# Data Partititoning Additions to Working Draft

#### D.3.1.3.1 Use of Data Partitioning

A two layer data partitioning scheme, and the use of two channels with different error performance, are used here is illustrate the error concealment capability of layering.

At the coder the value of the PBP pointer may be different for each slice such that the distribution of bits between the two channels (may be controlled (e.g. maintained constant). The distribution may be different for L.P. and P. trames.

Data partition 0 contains the address and control information and the low frequency DCT coefficients, while this partition t contains the high frequency DCT coefficients.

Two adependent channels are required. Channel errors are distributed so that data partition 1 receives most errors

It is assumed that at the decoder channel errors can be detected so that actions can be taken to prevent errored data from being displayed. For data partition 1 errored data is simply not displayed. For data partition 3 a suitably lined, quality channel, or decoder concealment actions are required.

# D.4.1 Bitrate allocation in Data Partitioning

Data partitioning allows splitting a bitstream for increased error resilience when two channels with different error performance are available. It is often required to constrain the bitrate of each partition. This can be achieved at the encoder by adaptively changing priority breakpoint at each slice.

The encoder can use two virtual buffers for HP and LP channels, and implement feedback rate control by picking a priority breakpoint that approximately meets the target rate for each channel. Difference between target and actual rates is used to revise the target for the next frame in a feedback loop.

It is desirable to vary the bitrate split from frame to frame for higher error resilience. Typically, I frames benefit from a higher HP/LP ratio, while B frames could be placed entirely in the LP partition.

### 1.4.2.4 Data partitioning extension

Data Partitoning—is a tool intended for use when multiple channels are available for transmission and or storage of a video bitstream, as may be the case in ATM networks, terrestrial broadcast, magnetic media etc. The bitstream is multiplexed onto these channels such that more critical parts of the bitstream (such as headers, motion vectors, DC coefficients) are transmitted in the channel with the better error performance, and less critical data (such as higher DCT coefficients) are transmitted in the channel with poor error performance. Thus, degradation to channel errors are minimized since the critical parts of a bitstream are better protected.

## 6.2.? Sequence\_data\_partitioning extension

This extension will be present between sequence and picture headers following the sequence\_scalable\_extension when the scalable\_mode field in this extension indicates the presence of data partitioning.

**partition\_identifier** — This is a 4-bit integer that identifies the partition number of the current bitstream. Data partitions are multiplexed at the systems layer, and partitition 0 is the base, or high priority layer, while partition 1 is the low priority layer. All other values are reserved.

### 6.?.? Slice layer in data partitioned bitstreams

**priority\_breakpoint** — This is a 7 bit integer that indicates the point in the syntax where the bitstream is supposed to be partitioned. The allowed values and their semantic interpretation is given in Table ?-?. Note that when MSB of priority\_breakpoint is set, the remaining 6 bits indicate the number of (run.level) DCT codewords present in the base partition.

sequence_data_partitioning extension() {	No. of bits	Mnemonio
extension_start_code	32	bslbf
extension_start_code_identifier	+	uimsbf
partition_identifier	4	uimsof
next_start_code()		
!		

slice() {	No. of bits	Mnemonic
slice_start_code	32	bslbf
if <data extension="" partitioning="" present="" was=""></data>		
priority_breakpoint	_	uimsbf
if ((vertical_size_value > 2800) && (! <iso 11172-2))<="" iec="" td=""><td></td><td></td></iso>		
slice_vertical_position_extension	3	uimsbf
quantizer_scale_code	5	uimsbi`
while ( nextbits() == '1' ) {		
extra_bit_slice	1	"1"
extra_information_slice	8	
}		
extra_bit_slice	1	"0"
do (		
macroblock()		
} while ( nextbits() != '000 0000 0000 0000 0000 0000')		
next_start_code()		
)		

Table ?-? Priority breakpoint values and associated semantics

priority_break point	Syntax elements included in the high priority stream				
0	All data at the sequence. GOP, picture and slice layers down to extra_bit_slice.				
1	All data included above, plus macroblock syntax elements up to and including macroblock address increment.				
2	All data included above, plus macroblock syntax elements up to but not including coded_block_pattern()				
3 63	Roserved.				
64	All syntax elements up to and including coded_block_pattern() or DC coefficient (det_dc_differential).				
65	All syntax elements above, plus 1 (run.level) DCT coefficient pair.				
64- <i>j</i>	All syntax elements above, plus $j$ (run.level) DCT coefficient pairs.				

Quant scale	DC Coeff	DCT Coeff 1	DCT Coeff 2	DCT Coeff 3	EOB	DC Coeff	DCT Coeff 1	EOB	
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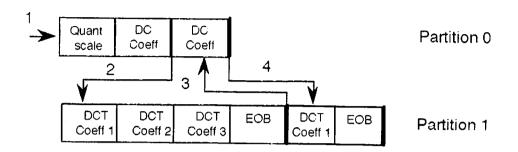


Figure ?-? A segment from a bitstream with two partitions, with priority\_breakpoint set to 64 (up to DC coefficient). The two partitions are shown, with arrows indicating how the decoder needs to switch between partitions.

#### 7.? Data Partitioning

Data partitioning is a technique that splits a video bitstream into two layers. A priority breakpoint indicates which syntax elements are placed in the base partition (which is also called HP, or high priority partition). The remainder of the bitstream is placed in the second partition (which is also called LP, or low priority partition). Sequence, GOP, and picture headers and extensions are redundantly copied in the LP partition to facilitate error recovery. Semantic interpretation of priority\_breakpoint is given in Table ?-?. The decoding process is modified in the following manner:

Set *current\_partition* to 0, and start decoding from bitstream with the identifier *current\_partition* If *current\_partition*=0, check to see if the current point in the bitstream is a priority breakpoint.

If yes, set current\_partition to 1. Next item will be decoded from LP partition.

Otherwise, continue decoding, Remove sequence, GOP, and picture headers from both partitions.

If current\_partition=1 check the priority breakpoint to see if the next item to be decoded is expected in the base partition.

If yes, set current\_partition to 0. Next item will be decoded from HP partition.

Otherwise, continue decoding,

An example is shown in Figure 2-2, where the priority breakpoint is set at 64 (up to DC coefficient).