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**OAM functions and mechanisms for Ethernet  
based networks**

Recommendation ITU-T G.8013/Y.1731



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# Recommendation ITU-T G.8013/Y.1731

## OAM functions and mechanisms for Ethernet based networks

### Summary

Recommendation ITU-T G.8013/Y.1731 provides mechanisms for user-plane OAM functionality in Ethernet networks according to the requirements and principles given in Recommendation ITU-T Y.1730. This Recommendation is designed specifically to support point-to-point connections and multipoint connectivity in the ETH layer as identified in Recommendation ITU-T G.8010/Y.1306.

The OAM mechanisms defined in this Recommendation offer capabilities to operate and maintain the network and service aspects of ETH layer.

### History

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## **Introduction**

ITU-T has prepared this Recommendation in cooperation with the IEEE Project 802.1ag (Connectivity Fault Management). Since the IEEE work is now complete, this Recommendation contains amendments to fully align the final results and include the appropriate normative references to IEEE documents. Moreover, further detailed work on the implementation details (i.e., the specification of the equipment functions) has been undertaken by ITU-T.



# Recommendation ITU-T G.8013/Y.1731

## OAM functions and mechanisms for Ethernet based networks

### 1 Scope

This Recommendation specifies mechanisms required to operate and maintain the network and service aspects of the ETH layer. It also specifies the Ethernet OAM frame formats and syntax and semantics of OAM frame fields. The OAM mechanisms as described in this Recommendation apply to both point-to-point ETH connections and multipoint ETH connectivity. The OAM mechanisms as described in this Recommendation are applicable to any environment independently of how the ETH layer is managed (e.g., using network management systems or operational support systems).

The architectural basis for this Recommendation is the Ethernet specification [ITU-T G.8010] which also accounts for [IEEE 802.1D], [IEEE 802.1Q], and [IEEE 802.3]. The OAM functions of the server layer networks used by the Ethernet network are not within the scope of this Recommendation. The OAM functions of the layers above the ETH layer are not within the scope of this Recommendation either.

### 2 References

The following ITU-T Recommendations and other references contain provisions, which, through reference in this text, constitute provisions of this Recommendation. At the time of publication, the editions indicated were valid. All Recommendations and other references are subject to revision; users of this Recommendation are therefore encouraged to investigate the possibility of applying the most recent edition of the Recommendations and other references listed below. A list of the currently valid ITU-T Recommendations is regularly published.

The reference to a document within this Recommendation does not give it, as a stand-alone document, the status of a Recommendation

- [ITU-T G.805] Recommendation ITU-T G.805 (2000), *Generic functional architecture of transport networks.*
- [ITU-T G.806] Recommendation ITU-T G.806 (2009), *Characteristics of transport equipment – Description methodology and generic functionality.*
- [ITU-T G.809] Recommendation ITU-T G.809 (2003), *Functional architecture of connectionless layer networks.*
- [ITU-T G.826] Recommendation ITU-T G.826 (2002), *End-to-end error performance parameters and objectives for international, constant bit-rate digital paths and connections.*
- [ITU-T G.7710] Recommendation ITU-T G.7710/Y.1701 (2007), *Common equipment management function requirements.*
- [ITU-T G.8001] Recommendation ITU-T G.8001/Y.1354 (2011), *Terms and definitions for Ethernet frames over transport.*
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- [IEEE 802.1Q] IEEE 802.1Q-2005, *IEEE Standard for Local and Metropolitan Area Networks: Virtual Bridged Local Area Networks*.  
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## 3 Definitions

### 3.1 Terms defined elsewhere

This Recommendation uses the following terms defined elsewhere:

- 3.1.1 **adaptation**: [ITU-T G.809].
- 3.1.2 **adapted information**: [ITU-T G.809].
- 3.1.3 **client/server relationship**: [ITU-T G.809].
- 3.1.4 **connection point**: [ITU-T G.805].
- 3.1.5 **connectionless trail**: [ITU-T G.809].
- 3.1.6 **defect**: [ITU-T G.806].
- 3.1.7 **failure**: [ITU-T G.806].
- 3.1.8 **flow**: [ITU-T G.809].
- 3.1.9 **flow domain**: [ITU-T G.809].
- 3.1.10 **flow domain flow**: [ITU-T G.809].
- 3.1.11 **flow point**: [ITU-T G.809].
- 3.1.12 **flow point pool**: [ITU-T G.809].
- 3.1.13 **flow point pool link**: [ITU-T G.809].
- 3.1.14 **flow termination**: [ITU-T G.809].
- 3.1.15 **flow termination sink**: [ITU-T G.809].
- 3.1.16 **flow termination source**: [ITU-T G.809].
- 3.1.17 **in-profile**: [ITU-T G.8001].
- 3.1.18 **in-service OAM**: [ITU-T G.8001].
- 3.1.19 **layer network**: [ITU-T G.809].
- 3.1.20 **link**: [ITU-T G.805].
- 3.1.21 **link connection**: [ITU-T G.805].
- 3.1.22 **link flow**: [ITU-T G.809].
- 3.1.23 **maintenance entity**: [ITU-T G.8001].
- 3.1.24 **maintenance entity group**: [ITU-T G.8001].
- 3.1.25 **MEG end point (MEP)**: [ITU-T G.8001].
- 3.1.26 **MEG intermediate point (MIP)**: [ITU-T G.8001].
- 3.1.27 **network**: [ITU-T G.809].
- 3.1.28 **network connection**: [ITU-T G.805].
- 3.1.29 **on-demand OAM**: [ITU-T G.8001].
- 3.1.30 **organizationally unique identifier**: [IEEE 802].
- 3.1.31 **out-of-service OAM**: [ITU-T G.8001].
- 3.1.32 **port**: [ITU-T G.809].
- 3.1.33 **proactive OAM**: [ITU-T G.8001].

- 3.1.34 reference point:** [ITU-T G.809].
- 3.1.35 server MEP:** [ITU-T G.8001].
- 3.1.36 termination connection point:** [ITU-T G.805].
- 3.1.37 termination flow point:** [ITU-T G.809].
- 3.1.38 traffic unit:** [ITU-T G.809].
- 3.1.39 trail:** [ITU-T G.805].
- 3.1.40 trail termination:** [ITU-T G.805].
- 3.1.41 transport:** [ITU-T G.809].
- 3.1.42 transport entity:** [ITU-T G.809].
- 3.1.43 transport processing function:** [ITU-T G.809].

#### **4 Abbreviations**

This Recommendation uses the following abbreviations:

1DM	One-way Delay Measurement
AIS	Alarm Indication Signal
APS	Automatic Protection Switching
CCM	Continuity Check Message
CoS	Class of Service
CSF	Client Signal Fail
CP	Connection Point
DA	Destination MAC Address
DEI	Drop Eligible Indicator
DMM	Delay Measurement Message
DMR	Delay Measurement Reply
ETH	Ethernet MAC layer network
ETH-AIS	Ethernet Alarm Indication Signal function
ETH-APS	Ethernet Automatic Protection Switching function
ETH-CC	Ethernet Continuity Check function
ETH-CSF	Ethernet Client Signal Fail function
ETH-DM	Ethernet Delay Measurement function
ETH-EXP	Ethernet Experimental OAM function
ETH-LCK	Ethernet Lock signal function
ETH-LB	Ethernet Loopback function
ETH-LM	Ethernet Loss Measurement function
ETH-LT	Ethernet Link Trace function
ETH-MCC	Ethernet Maintenance Communication Channel function
ETH-RDI	Ethernet Remote Defect Indication function

ETH-SLM	Ethernet Synthetic Loss Measurement function
ETH-Test	Ethernet Test function
ETH-TFP	Ethernet Termination Flow Point
ETH-VSP	Ethernet Vendor-Specific OAM function
ETY	Ethernet PHY layer network
EXM	Experimental OAM Message
EXR	Experimental OAM Reply
FLR	Frame Loss Ratio
FT	Flow Termination
ICC	ITU Carrier Code
LBM	Loopback Message
LBR	Loopback Reply
LCK	Locked
LMM	Loss Measurement Message
LMR	Loss Measurement Reply
LOC	Loss Of Continuity
LTM	Link Trace Message
LTR	Link Trace Reply
MAC	Media Access Control
MCC	Maintenance Communication Channel
ME	Maintenance Entity
MEG	ME Group
MEL	MEG Level
MEP	MEG End Point
MIP	MEG Intermediate Point
NMS	Network Management System
NNI	Network Node Interface
NT	Network Termination
OAM	Operation, Administration and Maintenance
OSS	Operations Support System
OTN	Optical Transport Network
OUI	Organizationally Unique Identifier
PCP	Priority Code Point
PDU	Protocol Data Unit
PE	Provider Edge
PHY	Ethernet Physical layer entity consisting of the PCS, the PMA, and, if present, the PMD sub layers

PRBS	Pseudo Random Bit Sequence
RDI	Remote Defect Indication
SA	Source MAC Address
SES	Severely Errored Seconds
SLA	Service Level Agreement
SLM	Synthetic Loss Message
SLR	Synthetic Loss Reply
SRV	Server
STP	Spanning Tree Protocol
TCI	Tag Control Information
TrCP	Traffic Conditioning Point
TLV	Type, Length, and Value
TST	Test PDU
TTL	Time To Live
UMC	Unique MEG ID Code
UNI	User Network Interface
UNI-C	Customer side of UNI
UNI-N	Network side of UNI
VLAN	Virtual LAN
VSM	Vendor-Specific OAM Message
VSR	Vendor-Specific OAM Reply

## 5 Conventions

The diagrammatic conventions for connection-oriented and connectionless layer networks described in this Recommendation are those of [ITU-T G.805], [ITU-T G.809] and [ITU-T G.8010].

For the purposes of this Recommendation, the following OAM terms and diagrammatic conventions are also defined.

### 5.1 ME group (MEG)

An ME group (MEG) includes different MEs that satisfy the following conditions:

- All MEs in a MEG exist in one same administrative boundary; and
- All MEs in a MEG have the same MEG level (see clause 5.3), and
- All MEs in a MEG belong to the same point-to-point ETH connection or multipoint ETH connectivity.

For a point-to-point ETH connection, a MEG contains a single ME. For a multipoint ETH connectivity containing  $n$  end-points, a MEG contains  $n*(n - 1)/2$  MEs.

### 5.2 Traffic conditioning point (TrCP)

A traffic conditioning point (TrCP) is an ETH flow point which is capable of an ETH traffic conditioning function, as specified in [ITU-T G.8010].

### 5.3 MEG level

In case MEGs are nested, the OAM flow of each MEG has to be clearly identifiable and separable from the OAM flows of the other MEGs. In cases where the OAM flows are not distinguishable by the ETH layer encapsulation itself, the MEG level in the OAM frame distinguishes between the OAM flows of nested MEGs.

Eight MEG levels are available to accommodate different network deployment scenarios.

When customer, provider, and operator data path flows are not distinguishable based on means of the ETH layer encapsulations, the eight MEG levels can be shared amongst them to distinguish between OAM frames belonging to nested MEGs of customers, providers and operators. The default MEG level assignment amongst customer, provider, and operator roles is:

- The customer role is assigned three MEG levels: 7, 6, and 5
- The provider role is assigned two MEG levels: 4 and 3
- The operator role is assigned three MEG levels: 2, 1, and 0

The default MEG level assignment can be changed via a mutual agreement among customer, provider, and/or operator roles.

Though eight MEG levels are available, not all MEG levels may be used. When not all eight MEG levels are used, there is no constraint on the continuity of MEG levels (e.g., MEG levels 7, 5, 2, and 0 may be used). The number of MEG levels used depends on the number of nested MEGs for which the OAM flows are not distinguishable based on means of the ETH layer encapsulation.

The specific assignment of MEG levels across different roles in specific deployments is outside the scope of this Recommendation. Refer to [ITU-T G.8010] for some examples.

### 5.4 OAM transparency

OAM transparency refers to the ability to allow transparent carrying of OAM frames belonging to higher-level MEGs across other lower-level MEGs when the MEGs are nested.

OAM frames belonging to an administrative domain originate and terminate in MEPs present at the boundary of that administrative domain. A MEP prevents OAM frames corresponding to a MEG in the administrative domain, from leaking outside that administrative domain. However, when a MEP is not present or is faulty, the associated OAM frames could leave the administrative domain.

Similarly, a MEP present at the boundary of an administrative domain protects the administrative domain from OAM frames belonging to other administrative domains. The MEP allows OAM frames from outside administrative domains belonging to higher-level MEGs to pass transparently; while it blocks OAM frames from outside administrative domains belonging to same or lower-level MEGs.

The customer role can use any of the eight MEG levels when the MEG levels are not shared with provider and operator roles, as mentioned in clause 5.3. However, if MEG Levels are shared with provider and operator roles, transparency of customer's OAM frames across provider's and/or operator's administrative domains will only be guaranteed for mutually agreed MEG levels, e.g., default MEG levels 7, 6 and 5. Similarly, transparency of a provider's OAM frames across an operator's administrative domain when MEG levels are shared will be guaranteed for mutually agreed MEG levels, e.g., default MEG levels 4 and 3, while the operator role can use default MEG levels 2, 1, and 0.

OAM frames can be prevented from leaking by implementing an OAM filtering process in the MEP atomic functions.

## 5.5 Representation of octets

In this Recommendation, octets are represented as defined in [IEEE 802.1D].

When consecutive octets are used to represent a binary number, the lower octet number has the most significant value. As an example, if Octet1 and Octet2 in Figure 5-1 represent a binary number, Octet1 has the most significant value.

The bits in an octet are numbered from 1 to 8, where bit 1 is the least significant bit (LSB) and bit 8 is the most significant bit (MSB).

		1								2								3								4							
		8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1		Octet1								Octet2								Octet3								Octet4							
5		Octet5								Octet6								Octet7								Octet8							
9		Octet9								Octet10								Octet11								Octet12							
:																																	

**Figure 5.5-1 – Example PDU format**

## 6 OAM relationships

### 6.1 MEs, MEPs, MIPs and TrCPs relationship

Appendix I provides different network scenarios to show how MEGs, MEPs and MIPs at different MEG levels can be deployed, and where TrCPs are likely to be placed.

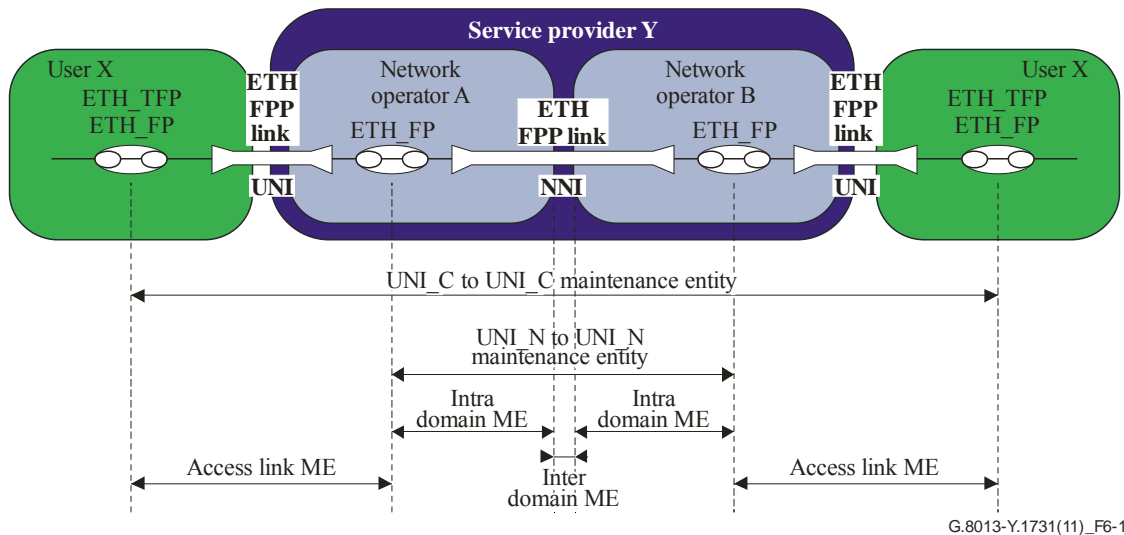
NOTE – Not all MEGs and corresponding MEPs and MIPs may be used or provided in the example network scenarios in Appendix I. For example, providers may not provide customer MIPs.

### 6.2 MEs, MEGs and MEG level relationship

The MEPs associated with an administrative domain operate at an assigned MEG level. Inter-domain MEPs, associated with MEGs between two administrative domains, can operate at a MEG level agreeable between the two administrative domains, such that associated inter-domain OAM flows are prevented from leaking into either administrative domain. The default MEG level for inter-domain OAM flows is 0.

MEs in Ethernet networks have been identified in Figure 7-1 of [ITU-T G.8010] (see Figure 6.2-1), Figure 7-2 of [ITU-T G.8010] and clause 9 of [ITU-T Y.1730]. MEs can nest but not overlap.





**Figure 6.2-1 – Example of MEs associated with a point-to-point connection administrative domain shown in Figure 7-1 of [ITU-T G.8010]**

Table 6-1 highlights possible MEG level assignments for MEGs within the context of customer, provider and operator administrative domains that share the MEG levels, as mapped to [ITU-T G.8010] and [ITU-T Y.1730].

**Table 6-1 – Example MEG level assignments for shared MEG levels**

ITU-T G.8010 MEG	ITU-T Y.1730 ME	MEG level(s)
UNI-C to UNI-C ME	UNI-UNI (Customer)	7, 6, or 5
UNI-N to UNI-N ME	UNI-UNI (Provider)	4, or 3
Intra-domain ME	Segment (PE-PE) intra-provider	4, or 3
Inter-domain ME	Segment (PE-PE) inter-provider (Provider – Provider)	0 (default)
Access link ME	ETY link OAM – UNI (Customer – Provider)	0 (default)
Inter-domain ME	ETY link OAM – NNI (Operator – Operator)	0 (default)

As mentioned in clause 5.3, MEG levels are shared when the OAM flows of nested MEGs of customer, provider, and operator cannot be distinguished based on ETH layer encapsulation. However, when OAM flows of nested MEGs of customer, provider, and operator can be distinguished based on ETH layer encapsulation, MEG levels are not shared except for inter-domain MEGs, (e.g., MEGs between customer and provider, MEGs between provider and operator, MEGs between operators, MEs between providers, etc.).

Table 6-2 highlights possible MEG level assignments for MEs within the context of customer, provider and operator administrative domains that do not share the MEG levels but require inter-domain MEs.

**Table 6-2 – Example MEG level assignments for independent MEG levels**

ITU-T G.8010 MEG	ITU-T Y.1730 ME	MEG level(s)
UNI-C to UNI-C ME	UNI-UNI (Customer)	7-1
UNI-N to UNI-N ME	UNI-UNI (Provider)	7-1
Intra-domain ME	Segment (PE-PE) intra-provider	7-1
Inter-domain ME	Segment (PE-PE) inter-provider (Provider – Provider)	0 (default)
Access Link ME	ETY link OAM – UNI (Customer – Provider)	0 (default)
Inter-domain ME	ETY link OAM – NNI (Operator – Operator)	0 (default)

Further, if inter-domain MEs are not required, each customer, provider, and operator can use all eight MEG levels. However, as already stated in clause 5.3, not all MEG levels may be used.

### 6.3 MEPs and MIPs configuration

MEPs and MIPs are configured via the management plane and/or control plane. The management plane configurations may be carried out through manual local administration of each device or via network management systems (NMS).

This configuration is outside the scope of this Recommendation.

## 7 OAM functions for fault management

OAM functions for fault management allow detection, verification, localization and notification of different defect conditions.

### 7.1 Ethernet continuity check (ETH-CC)

The Ethernet continuity check function (ETH-CC) is used for proactive OAM. It is used to detect loss of continuity (LOC) between any pair of MEPs in a MEG. ETH-CC also allows detection of unintended connectivity between two MEGs (mismerge), unintended connectivity within the MEG with an unexpected MEP (unexpected MEP), and other defect conditions (e.g., unexpected MEG level, unexpected period, etc.). ETH-CC is applicable for fault management, performance monitoring, or protection switching applications.

A MEP must always report reception of a frame with unexpected ETH-CC information. ETH-CC transmission may be enabled or disabled in a MEG. When ETH-CC transmission is enabled in a MEG, all MEPs are enabled to periodically transmit frames with ETH-CC information to all other MEPs in the MEG. The ETH-CC transmission period is the same for all MEPs in the MEG. When a MEP is enabled to generate frames with ETH-CC information, it also expects to receive frames with ETH-CC information from its peer MEPs in the MEG.

When ETH-CC transmission is disabled in a MEG, all MEPs are disabled to transmit frames with ETH-CC information.

The specific configuration information required by each MEP to support ETH-CC is the following:

- MEG ID – Identifies the MEG to which the MEP belongs.
- MEP ID – MEP's own identity in the MEG.
- List of peer MEP IDs – List of peer MEPs in the MEG. For a point-to-point MEG with a single ME, the list would consist of a single MEP ID for the peer.
- MEG level – MEG level at which the MEP exists.
- ETH-CC transmission period – This is application dependent. ETH-CC has three different applications (for each application, a default transmission period is specified):

- Fault management: default transmission period is 1 s (i.e., transmission rate of 1 frame/second)
- Performance monitoring: default transmission period is 100 ms (i.e., transmission rate of 10 frames/second)
- Protection switching: default transmission period is 3.33 ms (i.e., transmission rate of 300 frames/second).
- Priority – Identifies the priority of frame with ETH-CC information. By default, the frame with ETH-CC information is transmitted with the highest priority available to the data traffic. This information is configurable per operation.
- Drop eligibility – Frames with ETH-CC information are always marked as drop ineligible. This information is not necessarily configured.

A MIP is transparent to the ETH-CC information and therefore does not require any configuration information to support ETH-CC.

When a MEP does not receive ETH-CC information from a peer MEP, in the list of peer MEPs, within an interval of 3.5 times the ETH-CC transmission period, it detects loss of continuity to that peer MEP. The interval corresponds to a loss of three consecutive frames carrying ETH-CC information from the peer MEP. ETH-CC also allows detection of other defect conditions as described in clause 7.1.2.

The OAM PDU used for ETH-CC information is CCM, as described in clause 9.2. Frames which carry the CCM PDU are called CCM frames.

#### **7.1.1 CCM (with ETH-CC information) transmission**

When ETH-CC is enabled, a MEP periodically transmits CCM frames as often as the configured transmission period. The transmission period can be one of the following seven values:

- 3.33 ms: default transmission period for protection switching application (transmission rate of 300 frames/second)
- 10 ms: (transmission rate is 100 frames/second)
- 100 ms: default transmission period for performance monitoring application (transmission rate of 10 frames/second)
- 1 s: default transmission period for fault management application (transmission rate of 1 frame/second)
- 10 s: (transmission rate of 6 frames/minute)
- 1 min: (transmission rate of 1 frame/minute)
- 10 min: (transmission rate of 6 frames/hour)

NOTE – Even though seven different values are specified for the transmission period, the default values are recommended based on the application area for which ETH-CC is being used. When a transmission period other than the default value for an application area is used, the behaviour of the intended application is not guaranteed.

The period field in CCM is transmitted with a value for the transmission period configured at the transmitting MEP, so that a receiving MEP can detect Unexpected Period, if the transmission period is not the same across the transmitting and receiving MEPs.

### **7.1.2 CCM (with ETH-CC information) reception**

When a MEP receives a CCM frame, it examines it to ensure that its MEG ID matches the configured MEG ID in the receiving MEP, and that the MEP ID in the CCM frame is one from the configured list of peer MEP IDs. The information in the CCM frame is catalogued in the receiving MEP.

CCM frames allow the detection of different defect conditions, which include:

- If no CCM frames from a peer MEP are received within the interval equal to 3.5 times the receiving MEP's CCM transmission period, loss of continuity with peer MEP is detected.
- If a CCM frame with a MEG level lower than the receiving MEP's MEG level is received, Unexpected MEG Level is detected.
- If a CCM frame with the same MEG level but a MEG ID different than the receiving MEP's own MEG ID is received, Mismatch is detected.
- If a CCM frame with the same MEG level and a correct MEG ID but with an incorrect MEP ID, including the receiving MEP's own MEP ID, is received, Unexpected MEP is detected.
- If a CCM frame is received with a correct MEG level, a correct MEG ID, a correct MEP ID, but with a period field value different than the receiving MEP's own CCM transmission period, Unexpected Period is detected.

A receiving MEP must notify the equipment fault management process when it detects the above defect conditions.

## **7.2 Ethernet loopback (ETH-LB)**

The Ethernet loopback function (ETH-LB) is used to verify connectivity of a MEP with a MIP or peer MEP(s). There are two ETH-LB types:

- Unicast ETH-LB.
- Multicast ETH-LB.

### **7.2.1 Unicast ETH-LB**

Unicast ETH-LB is an on-demand OAM function that can be used for the following applications:

- To verify bidirectional connectivity of a MEP with a MIP or a peer MEP.
- To perform a bidirectional in-service or out-of-service diagnostics test between a pair of peer MEPs. This includes verifying bandwidth throughput, detecting bit errors, etc.

Frames with unicast ETH-LB information can be transmitted in several ways for different on-demand command types, e.g., single transmission, repetitive transmission, etc. The specific on-demand command types are outside the scope of this Recommendation.

When used to verify bidirectional connectivity, a MEP sends a unicast frame with ETH-LB request information and expects to receive a unicast frame with ETH-LB reply information from a MIP or peer MEP within a specified period of time. The MIP or peer MEP is identified by its MAC address. This MAC address is encoded in the DA of the unicast request frame. If the MEP does not receive the unicast frame with ETH-LB reply information within the specified period of time, loss of connectivity with the MIP or peer MEP can be inferred. Unicast ETH-LB can also be used to test the bidirectional connectivity with different frame sizes between a MEP and a MIP or peer MEP.

When used for performing bidirectional diagnostics tests, a MEP sends unicast frames with ETH-LB request information to a peer MEP. This ETH-LB request information includes test patterns. When out-of-service diagnostic tests are performed, data traffic is not delivered on either side of the diagnosed ME. Instead the MEPs are configured to send frames with ETH-LCK information, as described in clause 7.6, for the immediate client MEG level on either side of the ME.

NOTE 1 – Unicast ETH-LB can be used to perform only one of the two applications at any time. It must finish the pending on-demand command related to one application (either connectivity verification or diagnostic test) before it can act on a new on-demand command for the other application.

NOTE 2 – The maximum rate at which frames with Unicast ETH-LB information can be sent without adversely impacting the data traffic for in-service bidirectional connectivity verification or in-service bidirectional diagnostic tests is outside the scope of this Recommendation. It may be mutually agreed between the user of Unicast ETH-LB and the user of the service.

Specific configuration information required by a MEP to support unicast ETH-LB is the following:

- MEG level – MEG level at which the MEP exists.
- Unicast MAC address of remote MIP or MEP to which ETH-LB is intended. This information is configurable per operation.
- Data – Optional element whose length and contents are configurable at the MEP. The contents can be a test pattern and an optional checksum. Examples of test patterns include pseudo-random bit sequence (PRBS) ( $2^{31}-1$ ) as specified in clause 5.8 of [ITU-T O.150], all '0' pattern, etc. For bidirectional diagnostic test applications, configuration is required for a test signal generator and a test signal detector associated with the MEP.
- Priority – Identifies the priority of frames with unicast ETH-LB information.
- Drop eligibility – Identifies the eligibility of frames with unicast ETH-LB information to be discarded when congestion conditions are encountered.

NOTE 3 – Additional configuration information elements may be needed for repetitive transmission, e.g., repetition rate, total interval of repetition, etc. These additional configuration information elements are outside the scope of this Recommendation.

A remote MIP or MEP, upon receiving the unicast frame with ETH-LB request information which is addressed to the MIP or MEP, responds with a unicast frame with ETH-LB reply information.

Specific configuration information required by a MIP to support unicast ETH-LB is the following:

- MEG level – MEG level at which the MIP exists.

The OAM PDU used for unicast-LB request information is LBM, as described in clause 9.3. The OAM PDU used for unicast-LB reply information is LBR, as described in clause 9.4. Unicast frames carrying the LBM PDU are called unicast LBM frames. Unicast frames carrying the LBR PDU are called unicast LBR frames.

#### **7.2.1.1 Unicast LBM transmission**

Unicast LBM frames are transmitted by a MEP on an on-demand basis.

When used for bidirectional connectivity verification, a MEP transmits a unicast LBM frame addressed to the remote MIP or remote peer MEP with a specific transaction ID inserted in the Transaction ID/Sequence Number field. After unicast LBM frame transmission, a MEP expects to receive a unicast LBR frame within 5 seconds. The transmitted transaction ID is therefore retained by the MEP for at least 5 seconds after the unicast LBM frame is transmitted. A different transaction ID must be used for every unicast LBM frame, and no transaction ID from the same MEP may be repeated within one minute.

A MEP can optionally use a data TLV or test TLV. When configured for checking the successful transmission of different frame sizes, the MEP uses a data TLV. However, when used for diagnostic tests, a MEP transmits a unicast LBM frame addressed to the remote peer MEP with a test TLV. The test TLV is used to carry the test pattern generated by a test signal generator associated with the MEP. When the MEP is configured for an out-of-service diagnostic test, the MEP also generates LCK frames, as described in clause 7.6, at the client MEG level.

### **7.2.1.2 Unicast LBM reception and LBR transmission**

Whenever a valid unicast LBM frame is received by a MIP or MEP, an LBR frame is generated and transmitted to the requesting MEP. A unicast LBM frame with a valid MEG level and a destination MAC address equal to the MAC address of receiving MIP or MEP is considered to be a valid unicast LBM frame. Every field in the unicast LBM frame is copied to the LBR frame with the following exceptions:

- The source and destination MAC addresses are swapped.
- The OpCode field is changed from LBM to LBR.

Further, when a receiving MEP is configured for an out-of-service diagnostic test, it also generates LCK frames, as described in clause 7.6, at the client MEG level.

### **7.2.1.3 LBR reception**

When a MEP configured for connectivity verification receives an LBR frame addressed to it with the same MEG level as its own MEG level, and with an expected transaction ID and within 5 seconds after transmitting the unicast LBM frame, the LBR frame is valid. Otherwise the LBR frame addressed to it is invalid and is discarded.

When a MEP configured for a diagnostics test receives an LBR frame addressed to it with the same MEG level as its own MEG level, the LBR frame is valid. The test signal receiver associated with MEP may also validate the received sequence number against expected sequence numbers.

If a MIP receives an LBR frame addressed to it, such an LBR frame is invalid and the MIP should discard it.

### **7.2.2 Multicast ETH-LB**

The multicast ETH-LB function is used to verify bidirectional connectivity of a MEP with its peer MEPs. Multicast ETH-LB is an on-demand OAM function. When a multicast ETH-LB function is invoked on a MEP, the MEP returns to the initiator of multicast ETH-LB a list of its peer MEPs with whom the bidirectional connectivity is detected.

When multicast-LB is invoked on a MEP, a multicast frame with ETH-LB request information is sent from a MEP to other peer MEPs in the same MEG. The MEP expects to receive a unicast frame with ETH-LB reply information from its peer MEPs within a specified period of time. Upon reception of a multicast frame with ETH-LB request information, the receiving MEPs validate the multicast frame with ETH-LB request information and transmit a unicast frame with ETH-LB reply information after a randomized delay in the range of 0 to 1 second.

Specific configuration information required by each MEP to support multicast ETH-LB is the following:

- MEG level – MEG level at which the MEP exists.
- Priority – Identifies the priority of Multicast frames with ETH-LB request information. This information is configurable per operation.
- Drop eligibility – Multicast frames with ETH-LB request information are always marked as drop ineligible. This information is not necessarily configured.

A MIP is transparent to the multicast frames with ETH-LB request information and therefore does not require any information to support multicast ETH-LB.

The OAM PDU used for multicast ETH-LB request information is LBM, as described in clause 9.3. The OAM PDU used for ETH-LB reply is LBR, as described in clause 9.4. Multicast frames carrying the LBM PDU are called multicast LBM frames.

#### **7.2.2.1 Multicast LBM transmission**

Multicast LBM frames are transmitted by a MEP on an on-demand basis. After transmitting the multicast LBM frame with a specific transaction ID, the MEP expects to receive LBR frames within 5 seconds. The transmitted transaction ID is therefore retained for at least 5 seconds after the multicast LBM frame is transmitted. A different transaction ID must be used for every multicast LBM frame, and no transaction ID from the same MEP may be repeated within one minute.

#### **7.2.2.2 Multicast LBM reception and LBR transmission**

Whenever a valid multicast LBM frame is received by a MEP, an LBR frame is generated and transmitted to the requesting MEP after a randomized delay in the range of 0 to 1 second. The validity of the multicast LBM frame is determined based on the correct MEG level.

Every field in the multicast LBM frame is copied to the LBR frame with the following exceptions:

- The source MAC address in the LBR frame is the unicast MAC address of the replying MEP. The destination MAC address in the LBR frame is copied from the source MAC address of the multicast LBM frame which should be a unicast address.
- The OpCode field is changed from LBM to LBR.

#### **7.2.2.3 LBR reception**

When an LBR frame is received by a MEP with an expected transaction ID and within 5 seconds of transmitting the multicast LBM frame, the LBR frame is valid. If a MEP receives an LBR frame with a transaction ID that is not in the list of transmitted transaction IDs maintained by the MEP, the LBR frame is invalid and is discarded.

If a MIP receives an LBR frame addressed to it, such an LBR frame is invalid and the MIP should discard it.

### **7.3 Ethernet link trace (ETH-LT)**

The Ethernet link trace function (ETH-LT) is an on-demand OAM function that can be used for the following two purposes:

- Adjacent relation retrieval – The ETH-LT function can be used to retrieve adjacency relationship between a MEP and a remote MEP or MIP. The result of running ETH-LT function is a sequence of MIPs from the source MEP until the target MIP or MEP. Each MIP and/or MEP is identified by its MAC address.
- Fault localization – The ETH-LT function can be used for fault localization. When a fault (e.g., a link and/or a device failure) or a forwarding plane loop occurs, the sequence of MIPs and/or MEP will likely be different from the expected one. Difference in the sequences provides information about the fault location.

ETH-LT request information is initiated in a MEP on an on-demand basis. After transmitting a frame with ETH-LT request information, the MEP expects to receive frames with ETH-LT reply information within a specified period of time. Network elements containing MIPs or MEPs and receiving the frame with ETH-LT request information respond selectively with frames containing ETH-LT reply information.

A network element containing MIP or MEP responds with a frame with ETH-LT reply information upon receiving a valid frame with ETH-LT request information only if:

- the network element where the MIP or MEP resides is aware of the TargetMAC address in the ETH-LT request information and associates it to a single egress port, where the egress port is not the same as the port on which the frame with ETH-LT request information was received; or
- the TargetMAC address is the same as the MIP's or MEP's own MAC address.

A network element containing MIPs may also relay the frame with ETH-LT request information, as described in clause 7.3.2.

The specific configuration information required by a MEP to support ETH-LT is the following:

- MEG level – MEG level at which the MEP exists.
- Priority – Identifies the priority of the frames with ETH-LT request information. This information is configurable per operation.
- Drop eligibility – Frames with ETH-LT information are always marked as drop ineligible. This information is not necessarily configured.
- Target MAC address (usually of MIPs or MEPs of the MEG, but not limited to that) for which ETH-LT is intended. This information is configurable per operation.
- TTL – Allows the receiver to determine if frames with ETH-LT request information can be terminated. TTL is decremented every time frames with ETH-LT request information are relayed. Frames with ETH-LT request information with  $TTL \leq 1$  are not relayed.

The specific configuration information required by a MIP to support ETH-LT is the following:

- MEG level – MEG level at which the MIP exists.

The PDU used for ETH-LT request information is LTM, as described in clause 9.5. The PDU used for ETH-LT reply information is LTR, as described in clause 9.6. Frames carrying the LTM PDU are called LTM frames. Frames carrying the LTR PDU are called LTR frames.

NOTE 1 – As each network element, containing the MIPs or MEP, needs to be aware of the TargetMAC address in the received LTM frame and associates it to a single egress port, in order that the MIP or MEP can process the received LTM frames, a Unicast ETH-LB to the TargetMAC address could be performed by a MEP before transmitting the LTM frame. This would ensure that the network elements along the path to the TargetMAC address would have information about the route to the TargetMAC address if the TargetMAC address is reachable in the same MEG.

NOTE 2 – During a failure condition the information about the route to the TargetMAC address may age out after a certain time. The ETH-LT function has to be performed before the age out occurs in order to provide information about the route.

### 7.3.1 LTM transmission

An LTM frame is transmitted by a MEP on an on-demand basis. If the MEP resides at an ingress port, the LTM frame is forwarded towards the network element's own ETH-LT responder. However, if the MEP resides on an egress port, the LTM frame is transmitted out of that egress port. The LTM frame contains an LTM egress identifier TLV which identifies the network element initiating the LTM frame.

NOTE – ETH-LT responder is not defined in [ITU-T Y.1731] but just MEP and MIP of ingress and egress port are defined. And LTM egress identifier TLV is regarded as optional in [ITU-T Y.1731].

After transmitting the LTM frame with a specific transaction number, the MEP expects to receive LTR frames within 5 seconds. The transaction number of each LTM frame transmitted is therefore retained for at least 5 seconds after the LTM frame is transmitted. A different transaction number must be used for every LTM frame, and no transaction number from the same MEP may be repeated within one minute.



### 7.3.2 LTM reception and forwarding, and LTR transmission

If an LTM frame is received by a MEP or MIP, it forwards the LTM frame to the network element's ETH-LT responder, which performs the following validation:

- Only LTM frames with the same MEG level as the receiving MEP's or MIP's own MEG level are validated.
- Thereafter, the TTL field value of the LTM frame is checked. If the TTL field value is 0, the LTM frame is discarded (a TTL field value of 0 is an invalid value).
- Thereafter, the LTM frame is checked to see if LTM egress identifier TLV is present. The LTM frame is discarded if it does not contain LTM egress identifier TLV. It is noted that the LTM frame generated by [ITU-T Y.1731] may not contain the LTM egress identifier TLV. See Annex B for keeping the compatibility, i.e., LTM frame TLV may be processed at the MIP or MEP even if the LTM egress identifier TLV is absent.

If the LTM frame is valid, the ETH-LT responder does the following:

- It determines the destination address for the LTR frame from the OriginMAC address in the received LTM frame.
- If the network element is aware of the TargetMAC address in the LTM frame and associates it with a single egress port, where the egress port is not the same as the ingress port, or the LTM frame terminates at the MIP or MEP (when the TargetMAC address is the MIP's or MEP's own MAC address), an LTR frame is sent backwards to the originating MEP after a random time interval in the range of 0 to 1 second.
- Further, if the above condition applies and the LTM frame does not terminate at the MIP (when the TargetMAC address is not the same as the MIP's own address, if received by a MIP) and the TTL field in LTM frame is greater than 1, the LTM frame is forwarded towards the single egress port. All the fields of the relayed LTM frame are the same as the original LTM frame except for TTL which is decremented by 1, the source address which becomes the MIP's own MAC address, and LTM egress identifier TLV which identifies the network element relaying the modified LTM frame. It is noted that MIPs supporting [ITU-T Y.1731] may forward the LTM egress identifier TLV as it is. See Annex B for keeping the compatibility.
- Further, when the TargetMAC address is not the same as the MEP's own address, if received by a MEP, the LTM frames are always terminated at the MEP and the MEP does not send back the LTR frames.

The LTR frame contains LTR egress identifier TLV which identifies the source and destination of the LTM that triggered the transmission of this LTR. LTR egress identifier TLV contains the Last Egress Identifier field which identifies the network element that originated or forwarded the LTM frame for which this LTR frame is the response. This field takes the same value as the LTM egress identifier TLV of that LTM frame. The LTR egress identifier TLV also contains the Next Egress Identifier field which identifies the network element that transmitted this LTR frame, and can relay a modified LTM frame to the next hop. This field takes the same value as the LTM egress identifier TLV of the relayed modified LTM frame, if any. If no modified LTM frame is relayed, the FwdYes bit of the Flags field in the LTM frame is clear and the contents of Next Egress Identifier are undefined, and must be ignored by the LTR frame receiver.

Additionally, if the LTM frame was received by a MIP or MEP at an ingress port, the LTR frame includes a reply ingress TLV which describes the MIP or MEP at the ingress port.

Similarly, if the LTM frame was not received by a MEP at the ingress port, and if the egress port has a MIP or MEP, the LTR frame includes a reply egress TLV which describes the MIP or MEP at the egress port.

It is noted that both including Reply Ingress TLV and Reply Egress TLV are documented as optional in [ITU-T Y.1731 (06)] so that they may not be included in the LTR frame of that version. See Annex B for keeping the compatibility.

### **7.3.3 LTR reception**

When an LTR frame is received by a MEP with an expected transaction number and within 5 seconds of transmitting the LTM frame, the LTR frame is valid. If a MEP receives an LTR frame with a transaction number that is not in the list of transmitted transaction numbers maintained by the MEP, the LTR frame is invalid.

If a MIP receives an LTR frame addressed to it, such an LTR frame is invalid and the MIP should discard it.

### **7.4 Ethernet alarm indication signal (ETH-AIS)**

The Ethernet alarm indication signal function (ETH-AIS) is used to suppress alarms following detection of defect conditions at the server (sub) layer. Due to independent restoration capabilities provided within the Spanning Tree Protocol (STP) environments, ETH-AIS is not expected to be applied in the STP environments.

Transmission of frames with ETH-AIS information can be enabled or disabled on a MEP (or on a server MEP).

Frames with ETH-AIS information can be issued at the client MEG level by a MEP, including a server MEP, upon detecting defect conditions. For example, the defect conditions may include:

- Signal fail conditions in the case that ETH-CC is enabled.
- AIS condition or LCK condition in the case that ETH-CC is disabled.

NOTE – Since a Server MEP does not run ETH-CC, a server MEP can transmit frames with ETH-AIS information upon detection of any signal fail condition.

For multipoint ETH connectivity, a MEP cannot determine the specific server (sub) layer entity that has encountered defect conditions upon receiving a frame with ETH-AIS information. More importantly, it cannot determine the associated subset of its peer MEPs for which it should suppress alarms since the received ETH-AIS information does not contain that information. Therefore, upon reception of a frame with ETH-AIS information, the MEP will suppress alarms for all peer MEPs whether there is still connectivity or not.

However, for a point-to-point ETH connection, a MEP has only a single peer MEP. Therefore, there is no ambiguity regarding the peer MEP for which it should suppress alarms when it receives the ETH-AIS information.

Only a MEP, including a server MEP, is configured to issue frames with ETH-AIS information. Upon detecting a defect condition the MEP can immediately start transmitting periodic frames with ETH-AIS information at a configured client MEG level. A MEP continues to transmit periodic frames with ETH-AIS information until the defect condition is removed. Upon receiving a frame with ETH-AIS information a MEP detects AIS condition and suppresses loss of continuity alarms associated with all its peer MEPs. A MEP resumes loss of continuity alarm generation upon detecting loss of continuity defect conditions in the absence of AIS condition.

The specific configuration information required by a MEP to support ETH-AIS transmission is the following:

- Client MEG level – MEG level at which the most immediate client layer MIPs and MEPs exist.
- ETH-AIS transmission period – Determines transmission periodicity of frames with ETH-AIS information.

- Priority – Identifies the priority of frames with ETH-AIS information.
- Drop eligibility – Frames with ETH-AIS information are always marked as drop ineligible. This information is not necessarily configured.

Specific configuration information required by a MEP to support ETH-AIS reception is the following:

- Local MEG level – MEG level at which the MEP operates.

A MIP is transparent to frames with ETH-AIS information and therefore does not require any information to support ETH-AIS functionality.

The PDU used for ETH-AIS information is AIS, as described in clause 9.7. Frames carrying the AIS PDU are called AIS frames.

#### **7.4.1 AIS transmission**

A MEP, upon detecting a defect condition, can transmit AIS frames in a direction opposite to its peer MEP(s). The periodicity of AIS frames transmission is based on the AIS transmission period. An AIS transmission period of 1 second is recommended. The first AIS frame must always be transmitted immediately following the detection of a defect condition.

The client (sub) layer may consist of multiple MEGs that should be notified to suppress alarms resulting from defect conditions detected by the server (sub) layer MEP. The server (sub) layer MEP, upon detecting the signal fail condition, needs to send AIS frames to each of these client (sub) layer MEGs. In such cases, the first AIS frame for all client (sub) layer MEGs must be transmitted within 1 second of defect condition.

NOTE – To support ETH-AIS across current equipment, which may be stressed when issuing AIS frames every 1 second potentially across all 4094 VLANs, another AIS transmission period of 1 minute is also supported. An AIS frame communicates the used AIS transmission period via the Period field.

#### **7.4.2 AIS reception**

Upon receiving an AIS frame, a MEP examines it to ensure that its MEG level corresponds to its own MEG level. The Period field indicates the period at which the AIS frames can be expected. Upon receiving an AIS frame, the MEP detects AIS defect condition. Following the detection of AIS defect condition, if no AIS frames are received within an interval of 3.5 times the AIS transmission period indicated in the AIS frames received before, the MEP clears the AIS defect condition.

### **7.5 Ethernet remote defect indication (ETH-RDI)**

The Ethernet remote defect indication function (ETH-RDI) can be used by a MEP to communicate to its peer MEPs that a defect condition has been encountered. ETH-RDI is used only when ETH-CC transmission is enabled.

ETH-RDI has the following two applications:

- Single-ended fault management: The receiving MEP detects an RDI defect condition, which gets correlated with other defect conditions in this MEP and may become a fault cause. The absence of received ETH-RDI information in a single MEP indicates the absence of defects in the entire MEG.
- Contribution to far-end performance monitoring: It reflects that there was a defect condition in the far-end which is used as an input to the performance monitoring process.

A MEP that is in a defect condition transmits frames with ETH-RDI information. A MEP, upon receiving frames with ETH-RDI information, determines that its peer MEP has encountered a defect condition. However, for multipoint ETH connectivity, a MEP, upon receiving frames with ETH-RDI information, cannot determine the associated subset of its peer MEPs with which the

MEP transmitting RDI information encounters defect conditions, as the transmitting MEP itself does not always have that information.

The specific configuration information required by a MEP to support ETH-RDI function is the following:

- MEG level – MEG level at which the MEP exists.
- ETH-RDI transmission period – Application dependent and is configured to be the same as ETH-CC transmission period.
- Priority – Identifies the priority of frames with ETH-RDI information. The priority is the same as ETH-CC priority.
- Drop eligibility – Frames with ETH-RDI information are always marked as drop ineligible. This information is not necessarily configured.

A MIP is transparent to frames with ETH-RDI information and therefore does not require any configuration information to support ETH-RDI functionality.

The PDU used to carry ETH-RDI information is CCM, as described in clause 9.2.

#### **7.5.1 CCM with ETH-RDI transmission**

A MEP, upon detecting a defect condition with its peer MEP, sets the RDI field in the CCM frames for the duration of the defect condition. CCM frames, as described in clause 7.1.1, are transmitted periodically based on the CCM transmission period, when the MEP is enabled for CCM frames transmission. When the defect condition clears, the MEP clears the RDI field in the CCM frames in subsequent transmissions.

#### **7.5.2 CCM with ETH-RDI reception**

Upon receiving a CCM frame, a MEP examines it to ensure that its MEG level corresponds to its configured MEG level and detects RDI condition if the RDI field is set. For a point-to-point ETH connection, a MEP can clear the RDI condition when it receives the first CCM frame from its peer MEP with the RDI field cleared. For multipoint ETH connectivity, a MEP can clear the RDI condition when it receives the CCM frames from its entire list of peer MEP with the RDI field cleared.

### **7.6 Ethernet locked signal (ETH-LCK)**

The Ethernet locked signal function (ETH-LCK) is used to communicate the administrative locking of a server (sub) layer MEP and consequential interruption of data traffic forwarding towards the MEP expecting this traffic. It allows a MEP receiving frames with ETH-LCK information to differentiate between a defect condition and an administrative locking action at the server (sub) layer MEP. An example of an application that would require administrative locking of a MEP is the out-of-service ETH-Test, as described in clause 7.7.

A MEP continues to transmit periodic frames with ETH-LCK information at the configured client MEG level until the administrative/diagnostic condition is removed.

A MEP extracts frames with ETH-LCK information at its own MEG level and detects a LCK condition, which contributes to the signal fail condition of the MEP. The signal fail condition may result in the transmission of AIS frames to its client MEPs.

Specific configuration information required by a MEP to support ETH-LCK transmission is the following:

- Client MEG level – MEG level at which the most immediate client layer MIPs and MEPs exist.
- ETH-LCK transmission period – Determines transmission periodicity of frames with ETH-LCK information.

- Priority – Identifies the priority of frames with ETH-LCK information.
- Drop eligibility – Frames with ETH-LCK information are always marked as drop ineligible. This information is not necessarily configured.

The specific configuration information required by a MEP to support ETH-LCK reception is the following:

- Local MEG level – MEG level at which the MEP operates.

A MIP is transparent to the frames with ETH-LCK information and therefore does not require any information to support ETH-LCK functionality.

The PDU used for ETH-LCK information is LCK, as described in clause 9.8. Frames carrying the LCK PDU are called LCK frames.

### 7.6.1 LCK transmission

A (server) MEP, when administratively locked, transmits LCK frames to each of its client (sub-) layer MEGs as shown in Figure 7.6-1.



**Figure 7.6-1 – Example of ETH-LCK transmission**

The periodicity of LCK frames transmission is based on the LCK transmission period. The LCK transmission period is the same as the AIS transmission period. The first LCK frame must always be transmitted immediately following the administrative/diagnostic action.

The client (sub) layer may consist of multiple MEGs that should be notified to suppress alarms resulting from intentional maintenance/diagnostic related configuration at the server (sub) layer MEP. The server (sub) layer MEP, upon being administratively locked, needs to send LCK frames to each of its client (sub) layer MEGs. In such cases, the first LCK frame for all client (sub) layer MEGs must be transmitted within 1 second of defect condition.

### 7.6.2 LCK reception

Upon receiving an LCK frame, a MEP examines it to ensure that its MEG level corresponds to its configured MEG level. The Period field indicates the periodicity at which the LCK frames can be expected. Upon receiving an LCK frame, the MEP detects an LCK condition. Following detection of an LCK condition, if no LCK frames are received within an interval of 3.5 times the LCK transmission period indicated in the LCK frames received before, the MEP clears the LCK condition.

## 7.7 Ethernet test signal (ETH-Test)

The Ethernet test signal function (ETH-Test) is used to perform one-way on-demand in-service or out-of-service diagnostics tests. This includes verifying bandwidth throughput, frame loss, bit errors, etc.

When configured to perform such tests, a MEP inserts frames with ETH-Test information with specified throughput, frame size and transmission patterns.

When the out-of-service ETH-Test function is performed, client data traffic is disrupted in the diagnosed entity. The MEP configured for the out-of-service test transmits LCK frames, as described in clause 7.6, for the immediate client ETH (sub) layer.

When an in-service ETH-Test function is performed, data traffic is not disrupted and the frames with ETH-Test information are transmitted in such a manner that a limited part of the service bandwidth is utilized. This rate of transmission for frames with ETH-Test information is pre-determined for in-service ETH-Test function.

NOTE 1 – The maximum rate at which frames with ETH-Test information can be sent without adversely impacting the data traffic for an in-service ETH-Test is outside the scope of this Recommendation. It may be mutually agreed between the user of ETH-Test and the user of the service.

The specific configuration information required by a MEP to support ETH-Test is the following:

- MEG level – MEG level at which the MEP exists.
- Unicast MAC address of the peer MEP for which ETH-Test is intended. This information is configurable per operation.
- Data – Optional element whose length and contents are configurable at the MEP. The contents can be a test pattern and an optional checksum. Examples of test patterns include pseudo-random bit sequence (PRBS) ( $2^{31}-1$ ) as specified in clause 5.8 of [ITU-T O.150], all '0' pattern, etc. At the transmitting MEP, configuration is required for a test signal generator which is associated with the MEP. At a receiving MEP, configuration is required for a test signal detector which is associated with the MEP.
- Priority – Identifies the priority of frames with ETH-Test information. This information is configurable per operation.
- Drop eligibility – Identifies the eligibility of frames with ETH-Test information to be dropped when congestion conditions are encountered.

NOTE 2 – Additional configuration information elements may be needed, such as the transmission rate of ETH-Test information, the total interval of ETH-Test, etc. These additional configuration information elements are outside the scope of this Recommendation.

A MIP is transparent to the frames with ETH-Test information and therefore does not require any configuration information to support ETH-Test functionality.

A MEP inserts frames with ETH-Test information towards a targeted peer MEP. The receiving MEP detects these frames with ETH-Test information and makes the intended measurements.

The PDU used for ETH-Test information is TST, as described in clause 9.9. Frames carrying the TST PDU are called TST frames.

### **7.7.1 TST transmission**

A test signal generator associated with a MEP can transmit TST frames as often as the test signal generator configuration. Each TST frame is transmitted with a specific sequence number. A different sequence number must be used for every TST frame, and no sequence number from the same MEP may be repeated within one minute.

When a MEP is configured for an out-of-service test, the MEP also generates LCK frames for the immediate client MEG level.

### **7.7.2 TST reception**

When a MEP receives TST frames, it examines them to ensure that the MEG level corresponds to its own configured MEG level. If the receiving MEP is configured for ETH-TST function, the test signal detector associated with the MEP detects bit errors from the pseudo-random bit sequence of the received TST frames and reports such errors. Further, when the receiving MEP is configured for an out-of-service test, it also generates LCK frames for the client MEG level.

## **7.8 Ethernet automatic protection switching (ETH-APS)**

The Ethernet automatic protection switching function (ETH-APS) is used to control protection switching operations to enhance reliability. The specific details of protection switching operations are outside the scope of this Recommendation.

The OAM frame type used for ETH-APS is APS frame, as described in clause 9.10.

Applications of ETH-APS mechanisms are defined in Recommendations [ITU-T G.8031] and [ITU-T G.8032].

## **7.9 Ethernet maintenance communication channel (ETH-MCC)**

The Ethernet maintenance communication channel function (ETH-MCC) provides a maintenance communication channel between a pair of MEPs. ETH-MCC can be used to perform remote management. The specific use of ETH-MCC is outside the scope of this Recommendation.

A MEP can send a frame with ETH-MCC information to its peer MEP with remote maintenance request, remote maintenance reply, notification, etc.

Specific configuration information required by a MEP to support ETH-MCC is the following:

- MEG level – MEG level at which the MEP exists.
- Unicast MAC address of the remote MEP for which ETH-MCC is intended.
- OUI – Organizationally unique identifier used to identify the organization defining a specific format and meaning of ETH-MCC.
- Data – Additional information that may be needed and is dependent on the specific application of ETH-MCC. Application specific information is outside the scope of this Recommendation.
- Priority – Identifies the priority of frames with ETH-MCC information. This information is configurable per operation.
- Drop eligibility – Frames with ETH-MCC information are always marked as drop ineligible. This information is not necessarily configured.

A remote MEP, upon receiving a frame with ETH-MCC information and with a correct MEG level, passes the ETH-MCC information to the management agent which may additionally respond.

A MIP is transparent to the frames with ETH-MCC information and therefore does not require any configuration information to support ETH-MCC functionality.

The PDU used for ETH-MCC information is MCC, as described in clause 9.11. Frames carrying the MCC PDU are called MCC frames.

## **7.10 Ethernet experimental OAM (ETH-EXP)**

ETH-EXP is used for the experimental OAM functionality which can be used within an administrative domain on a temporary basis. Interoperability of the experimental OAM functionality is not expected across different administrative domains.

The specific application of ETH-EXP is outside the scope of this Recommendation.

EXM PDU, as described in clause 9.17, and EXR PDU, as described in clause 9.18, can be used for experimental OAM. Details of experimental OAM mechanisms are outside the scope of this Recommendation.

## **7.11 Ethernet vendor-specific OAM (ETH-VSP)**

ETH-VSP is used for vendor-specific OAM functionality which may be used by a vendor across its equipment. Interoperability of vendor-specific OAM functionality is not expected across different vendors' equipment.

The specific application of ETH-VSP is outside the scope of this Recommendation.

VSM PDU, as described in clause 9.19, and VSR PDU, as described in clause 9.20, can be used for vendor-specific OAM. Details of vendor-specific OAM mechanisms are outside the scope of this Recommendation.

## **7.12 Ethernet client signal fail (ETH-CSF)**

The Ethernet client signal fail function (ETH-CSF) is used by a MEP to propagate to a peer MEP the detection of a failure or defect event in an Ethernet client signal when the client itself does not support appropriate fault or defect detection or propagation mechanisms, such as ETH-CC or ETH-AIS. The ETH-CSF messages propagate in the direction from the Ethernet source-adaptation function detecting the failure or defect event to the Ethernet sink-adaptation function associated with the peer MEP.

ETH-CSF is only applicable to point-to-point Ethernet transport applications. In particular, the use of ETH-CSF with [IEEE 802.1Q], [IEEE 802.1ad] or other Ethernet Spanning Tree Protocol (STP)-based networking environments is strictly restricted to point-to-point segments of the Ethernet flow. The use of client signal fail indications to support client failure applications is described in Appendix VIII of [ITU-T G.806].

Specific configuration information required by a MEP to support ETH-CSF transmission is:

- Local MEG level – MEG level at which the source MEP operates.
- ETH-CSF transmission period – Determines transmission periodicity of frames with ETH-CSF information.
- Priority – Identifies the priority of frames with ETH-CSF information.
- Drop eligibility – Frames with ETH-CSF information are always marked as drop ineligible.

Specific configuration information required by a MEP to support ETH-CSF reception is:

- Local MEG level – MEG level at which the sink MEP operates.

A MIP is transparent to frames with ETH-CSF information and therefore does not require any information to support ETH-CSF functionality.

The ETH-CSF message indicates also the type of defect. Three CSF defect types are currently defined:

- Client Loss of Signal (C-LOS)
- Client Forward Defect Indication (C-FDI)
- Client Reverse Defect Indication (C-RDI)

The PDU used to convey ETH-CSF information is referred as CSF PDU, as described in clause 9.21. Frames carrying the ETH-CSF indications are also referred to as CSF frames.

### **7.12.1 CSF transmission**

Frames with ETH-CSF information can be issued by a MEP, upon notification of an Ethernet CSF event from its associated Ethernet client source adaptation function. Detection rules for Ethernet CSF events are Ethernet client and application specific.

Transmission of packets with CSF information can be enabled or disabled on a MEP.



Upon receiving an Ethernet CSF notification from its Ethernet client-specific source adaptation function the MEP can immediately start periodic transmission of frames with ETH-CSF information. A MEP continues periodic transmission of frames with ETH-CSF information until the Ethernet CSF indication is removed by the source adaptation function.

Clearing an Ethernet CSF condition by the Ethernet client-specific source adaptation function is Ethernet client and application specific. The clearance of the Ethernet CSF condition by the source adaptation function is communicated to the sink adaptation function associated with the peer MEP via:

- the non-sending ETH-CSF or
- the forwarding of a ETH-CSF PDU with Client Defect Clear Indication (C-DCI) information.

### **7.12.2 CSF reception**

Upon receiving a CSF frame with ETH-CSF information a MEP declares the beginning or end of an Ethernet remote CSF condition, depending on the received ETH-CSF information as described in [ITU-T G.8021], and propagates this Ethernet client defect condition towards the corresponding egress client port.

Clearing an Ethernet CSF condition by the Ethernet client-specific sink adaptation function is Ethernet client and application specific. The clearance of the Ethernet CSF condition by the sink adaptation function is communicated to the source adaptation function associated with the peer MEP when:

- no ETH-CSF frame is received in within an interval of N times the CSF transmission period ms (suggested value of N is 3.5), or
- an ETH-CSF PDU with Client Defect Clear Indication (C-DCI) information is received.

Note that consequent actions by the sink adaptation function associated with the MEP to propagate the received ETH-CSF information to the Ethernet client are by definition Ethernet client and application specific.

## **8 OAM functions for performance monitoring**

OAM functions for performance monitoring allow measurement of different performance parameters. The functions and measurement methods for point-to-point ETH connections and multipoint ETH connectivity are defined.

This Recommendation covers the following performance parameters which are based on [MEF 10.2].

- **Frame loss ratio**

Frame loss ratio is defined as a ratio, expressed as a percentage, of the number of frames not delivered divided by the total number of frames during time interval T, where the number of frames not delivered is the difference between the number of frames arriving at the ingress ETH flow point to be delivered to the egress ETH flow point and the number of frames delivered at the egress ETH flow point in a point-to-point ETH connection or multipoint ETH connectivity. The frame loss ratio may be measured using either service frames or synthetic frames, belonging to single CoS. The use of synthetic frames can also be applicable for multipoint ETH connectivity. The use of service frames is applicable only to point-to-point ETH connection where all the frames arriving at the ingress ETH flow point are to be delivered to the egress ETH flow point.

- **Frame delay**

Frame delay can be expressed as one-way delay for a frame, where one-way frame delay is defined as the time elapsed since the start of transmission of the first bit of the frame by a source node until the reception of the last bit of the same frame by the destination node. When two-way delay is measured, a loopback is performed at the frame's destination node and the frame is received at the original source node. In the round-trip case, there are four timestamps available which enable both one-way and two-way delay calculations. Ideally, the mean one-way frame delay should be accessible for a set of frames. Mean one-way frame delay is defined in [ITU-T Y.1563]. The service frames belong to the same CoS instance on a point-to-point ETH connection or multipoint ETH connectivity.

- **Frame delay variation**

Frame delay variation is a measure of the variations in the frame delay between a pair of service frames. The service frames belong to the same CoS instance on a point-to-point ETH connection or multipoint ETH connectivity.

- **Availability**

The Ethernet service definition is defined in [ITU-T Y.1563]. Although the mechanisms defined in this Recommendation can contribute to availability-related measurements, the details of measurement methods in this Recommendation are for further study.

Frame performance parameters are applicable to frames that conform to an agreed-upon priority level X and that are deemed by the network as not drop-eligible (i.e., so called "green" frames) of bandwidth profile conformance. Such "green" frames are also called in-profile (see [ITU-T G.8021]). Service frames are admitted at the ingress ETH flow point of a point-to-point ETH network, tandem or link connection and should be delivered to the egress ETH flow point.

In addition, another performance parameter is identified as per [b-IETF RFC 2544]:

- **Throughput**

Throughput is the average rate of successful traffic delivery over a communication channel. [b-IETF RFC 2544]-based procedure is typically used under test conditions, i.e., out-of-service test, where when there is no traffic for the tested ETH connection. The procedure for in-service testing is for further study.

## **8.1 Frame loss measurement (ETH-LM)**

ETH-LM is used to collect counter values applicable for ingress and egress service frames where the counters maintain a count of transmitted and received data frames between a pair of MEPs.

ETH-LM is performed by sending frames with ETH-LM information to a peer MEP and similarly receiving frames with ETH-LM information from the peer MEP. Each MEP performs frame loss measurements which contribute to unavailable time. Since a bidirectional service is defined as unavailable if either of the two directions is declared unavailable, ETH-LM must facilitate each MEP to perform near-end and far-end frame loss measurements.

For a MEP, near-end frame loss refers to frame loss associated with ingress data frames while far-end frame loss refers to frame loss associated with egress data frames. Both near-end and far-end frame loss measurements contribute to near-end severely errored seconds (near-end SES) and far-end severely errored seconds (far-end SES), respectively, which together contribute to unavailable time, in a manner similar to Recommendation [ITU-T G.826] and Recommendation [ITU-T G.7710].

A MEP maintains the following two local counters for each peer MEP and for each priority class being monitored in a point-to-point ME for which loss measurements are to be performed:

- TxFCI: counter for in-profile data frames transmitted towards the peer MEP.
- RxFCI: counter for in-profile data frames received from the peer MEP.

TxFCl and RxFCI counters do not count the OAM frames transmitted or received by the MEP at the MEP's MEG level in some conditions (see Notes). However, the counters do count OAM frames from the higher MEG levels that pass through the MEPs in a manner similar to the data frames.

NOTE 1 – For single-ended ETH-LM, both counters do not need to count on-demand OAM frames for ETH-LB, ETH-LT, ETH-LM, ETH-DM, and ETH-Test. Instead, proactive OAM frames for ETH-CC and ETH-APS need to be counted. For dual-ended ETH-LM, the counters do not need to count on-demand OAM frames for ETH-LB, ETH-LT, ETH-LM, ETH-DM, and ETH-Test and proactive OAM frames for ETH-CC. However, proactive OAM frames for ETH-APS need to be counted.

NOTE 2 – As OAM frames for ETH-AIS and ETH-LCK are only sent in the defect conditions where the result of loss measurements is invalid, it is unnecessary to count these frames.

The method of loss measurement involving pairs of consecutive frames with ETH-LM information, as shown in clauses 8.1.1.2 and 8.1.2.3, alleviates lack of synchronization across the initial counter values at the transmitting and receiving MEPs. Further, when a MEP detects a loss-of-continuity defect condition, it ignores loss measurements during the defect condition and assumes 100% losses.

NOTE 3 – The level of accuracy in the loss measurements is dependent on how frames with ETH-LM information are added to the data stream after the counter values are copied in the ETH-LM information. For example, if additional data frames get transmitted and/or received between the time of reading the counter values and adding the frame with ETH-LM information to the data stream, the counter values copied in ETH-LM information become inaccurate. However, a hardware-based implementation which is able to add frames with ETH-LM information to the data stream immediately after reading the counter values, provides enhanced accuracy.

NOTE 4 – Details on the processing of counters used for the transmitted and received data frames are described in [ITU-T G.8021].

NOTE 5 – In profile frames are so called 'green' frames where drop-eligibility is 'false'. Network operators or administrators can configure the encoding method to identify green frames. For example, green frames are those where DEI field is 'false', and yellow frames are those where such field is 'true'. PCP or PCP/DEI can be used for this identification.

Specific configuration information required by a MEP to support ETH-LM is the following:

- MEG level – MEG level at which the MEP exists.
- ETH-LM transmission period – The default transmission period is 100 ms (i.e., a transmission rate of 10 frames/second). The ETH-LM transmission period should be such that the frame and/or octet counters whose values are carried in ETH-LM information should not wrap around to the same value even if one or more ETH-LM frames are lost. This is primarily a concern for frame loss measurements at lower priority levels. Refer to clause II.2 for examples of frame counter wrapping periods.
- Priority – Identifies the priority of the frames with ETH-LM information. This information is configurable per operation.
- Drop eligibility – Frames with ETH-LM information are always marked as drop ineligible. This information is not necessarily configured.

A MIP is transparent to frames with ETH-LM information and therefore does not require any information to support ETH-LM functionality.

ETH-LM can be performed in two ways:

- Dual-ended ETH-LM.
- Single-ended ETH-LM.

### 8.1.1 Dual-ended ETH-LM

Dual-ended ETH-LM is used as proactive OAM for performance monitoring and is applicable to fault management. In this case, each MEP sends periodic dual-ended frames with ETH-LM information to its peer MEP in a point-to-point ME to facilitate frame loss measurements at the peer MEP. Each MEP terminates the dual-ended frames with ETH-LM information and makes the near-end and far-end loss measurements. This function is used for performance monitoring at the same priority level as used for ETH-CC.

The PDU used for dual-ended ETH-LM information is CCM, as described in clause 9.2.

#### 8.1.1.1 CCM with dual-ended ETH-LM transmission

When configured for proactive loss measurement, a MEP periodically transmits CCM frames with the following information elements:

- TxFCf: Value of the local counter TxFCI at the time of transmission of the CCM frame.
- RxFCb: Value of the local counter RxFCI at the time of reception of the last CCM frame from the peer MEP.
- TxFCb: Value of TxFCf in the last received CCM frame from the peer MEP.

The CCM PDU is transmitted with a period value equal to the CCM transmission period configured for performance monitoring application at the transmitting MEP. The receiving MEP detects an unexpected period defect condition if the CCM transmission period is not the same as the configured value.

#### 8.1.1.2 CCM with dual-ended ETH-LM frame reception

When configured for proactive loss measurement, a MEP, upon receiving a CCM frame, uses the following values to make near-end and far-end loss measurements:

- Received CCM frame's TxFCf, RxFCb, and TxFCb values and local counter RxFCI value at time this CCM frame was received. These values are represented as TxFCf[t<sub>c</sub>], RxFCb[t<sub>c</sub>], TxFCb[t<sub>c</sub>], and RxFCI[t<sub>c</sub>], where t<sub>c</sub> is the reception time of the current frame.
- Previous CCM frame's TxFCf, RxFCb, and TxFCb values and local counter RxFCI value at time the previous CCM frame was received. These values are represented as TxFCf[t<sub>p</sub>], RxFCb[t<sub>p</sub>], TxFCb[t<sub>p</sub>], and RxFCI[t<sub>p</sub>], where t<sub>p</sub> is the reception time of the previous frame.

$$\text{Frame Loss}_{\text{far-end}} = |\text{TxFCb}[t_c] - \text{TxFCb}[t_p]| - |\text{RxFCb}[t_c] - \text{RxFCb}[t_p]|$$

$$\text{Frame Loss}_{\text{near-end}} = |\text{TxFCf}[t_c] - \text{TxFCf}[t_p]| - |\text{RxFCI}[t_c] - \text{RxFCI}[t_p]|$$

If the Period field value in the received CCM frame is different than the MEP's own configured CCM transmission period, the MEP detects an unexpected period defect condition, in which case the frame loss measurements are not carried out.

### 8.1.2 Single-ended ETH-LM

Single-ended ETH-LM is used for on-demand OAM. In this case, a MEP sends frames with ETH-LM request information to its peer MEP and receives frames with ETH-LM reply information from its peer MEP to carry out loss measurements.

The PDU used for single-ended ETH-LM request is LMM, as described in clause 9.12. The PDU used for single-ended ETH-LM reply is LMR, as described in clause 9.13. Frames which carry the LMM PDU are called LMM frames. Frames which carry the LMR PDU are called LMR frames.

### 8.1.2.1 LMM transmission

For an on-demand loss measurement, a MEP periodically transmits LMM frames with the following information element:

- TxFCf: Value of the local counter TxFCI at the time of LMM frame transmission.

### 8.1.2.2 LMM reception and LMR transmission

Whenever a valid LMM frame is received by a MEP, an LMR frame is generated and transmitted to the requesting MEP. An LMM frame with a valid MEG level and a destination MAC address equal to the receiving MEP's MAC address is considered to be a valid LMM frame. An LMR frame contains the following values:

- TxFCf: Value of TxFCf copied from the LMM frame.
- RxFCf: Value of local counter RxFCI at the time of LMM frame reception.
- TxFCb: Value of local counter TxFCI at the time of LMR frame transmission.

### 8.1.2.3 LMR reception

Upon receiving an LMR frame, a MEP uses the following values to make near-end and far-end loss measurements:

- Received LMR frame's TxFCf, RxFCf, and TxFCb values and local counter RxFCI value at the time this LMR frame was received. These values are represented as TxFCf[t<sub>c</sub>], RxFCf[t<sub>c</sub>], TxFCb[t<sub>c</sub>], and RxFCI[t<sub>c</sub>], where t<sub>c</sub> is the reception time of the current reply frame.
- Previous LMR frame's TxFCf, RxFCf, and TxFCb values and local counter RxFCI value at the time the previous LMR frame was received. These values are represented as TxFCf[t<sub>p</sub>], RxFCf[t<sub>p</sub>], TxFCb[t<sub>p</sub>], and RxFCI[t<sub>p</sub>], where t<sub>p</sub> is the reception time of the previous reply frame.

$$\text{Frame Loss}_{\text{far-end}} = |\text{TxFCf}[t_c] - \text{TxFCf}[t_p]| - |\text{RxFCf}[t_c] - \text{RxFCf}[t_p]|$$

$$\text{Frame Loss}_{\text{near-end}} = |\text{TxFCb}[t_c] - \text{TxFCb}[t_p]| - |\text{RxFCI}[t_c] - \text{RxFCI}[t_p]|$$

## 8.2 Frame delay measurement (ETH-DM)

ETH-DM can be used for on-demand or proactive OAM to measure frame delay and frame delay variation. Frame delay and frame delay variation measurements are performed by sending periodic frames with ETH-DM information to the peer MEP and receiving frames with ETH-DM information from the peer MEP during proactive measurement session and/or the diagnostic interval. Each MEP may perform frame delay and frame delay variation measurement.

When a MEP is enabled to generate frames with ETH-DM information, it periodically sends frames with ETH-DM information to its peer MEP in the same ME. When a MEP is enabled to generate frames with ETH-DM information, it also expects to receive frames with ETH-DM information from its peer MEP in the same ME.

Specific configuration information required by a MEP to support ETH-DM is the following:

- MEG level – MEG level at which the MEP exists.
- Unicast MAC address of remote MEP to which ETH-DM is intended. Multicast MAC address is also allowed to use for multipoint-to-multipoint ETH connectivity. In case of a multipoint-to-multipoint ETH connectivity, a MEP may activate multiple monitoring toward different peer MEPs simultaneously. In this case, each MEP needs to manage the result of monitoring on a per-peer-MEP basis.

- DM application – Identifies the application, i.e., proactive versus on-demand delay measurement. This information is configurable per operation. A MEP may activate pro-active and on-demand monitoring simultaneously on the same CoS level and toward the same peer MEP. In this case, each MEP needs to manage the result of monitoring on a per peer-MEP basis.
- Data – Optional data element whose length is configurable at the MEP. The inclusion of the optional data element in the DM frame is to support configurable DM frame size.
- Priority – Identifies the priority of the frames with ETH-DM information. This information is configurable per operation. A MEP may activate the multiple monitoring on different CoS levels simultaneously. In this case, each MEP needs to manage the result of monitoring per CoS level.
- Drop eligibility – Frames with ETH-DM information are always marked as drop ineligible. This information is not necessarily configured.
- Test ID – Can optionally be used to distinguish each DM measurement if multiple measurements are simultaneously activated. It must be unique at least within the context of any DM measurement type (one-way/two-way and on-demand/pro-active) for the MEG and initiating MEP.

NOTE 1 – Additional configuration information elements may be needed, such as the transmission rate of ETH-DM information, the total interval of ETH-DM, etc. These additional configuration information elements are outside the scope of this Recommendation.

A MIP is transparent to the frames with ETH-DM information and therefore does not require any information to support ETH-DM functionality.

A MEP transmits frames with ETH-DM information with the following information element:

- TxTimeStampf: Timestamp at the transmission time of ETH-DM frame

The receiving MEP can compare this value with the RxTime<sub>f</sub>, the time at the reception of ETH-DM frame and calculate the one-way frame delay as:

$$\text{Frame Delay} = \text{RxTime}_f - \text{TxTimeStamp}_f$$

However, one-way frame delay measurement requires that the time and phase at the transmitting MEP and the receiving MEPs are synchronized. For the purposes of frame delay variation measurement, which is based on the difference between subsequent frame delay measurements, the requirement for the time and phase synchronizations can be relaxed since the out-of-phase period can be eliminated in the difference of subsequent frame delay measurements.

If it is not practical for the clocks to be synchronized, which is expected to be the most common scenario, the frame delay measurement can be made only for two-way measurements, where the MEP transmits a frame with ETH-DM request information with the TxTimeStamp<sub>f</sub>, and the receiving MEP responds with a frame with ETH-DM reply information with TxTimeStamp<sub>f</sub> copied from the ETH-DM request information. The MEP receiving the frame with ETH-DM reply information compares the TxTimeStamp<sub>f</sub> with the RxTime<sub>b</sub>, which is the time at the reception of frame with ETH-DM reply information and calculates the two-way frame delay as:

$$\text{Frame Delay} = \text{RxTime}_b - \text{TxTimeStamp}_f$$

The MEP can also make two-way frame delay variation measurements based on its ability to calculate the difference between two subsequent two-way frame delay measurements.

NOTE 2 – To allow a more precise two-way frame delay measurement, the MEP replying to frame with ETH-DM request information can also include two additional timestamps in the ETH-DM reply information: RxTimeStamp<sub>f</sub> (Timestamp at the time of receiving frame with ETH-DM request information), and TxTimeStamp<sub>b</sub> (Timestamp at the time of transmitting frame with ETH-DM reply information).

ETH-DM can be performed in two ways:

- One-way ETH-DM
- Two-way ETH-DM

### **8.2.1 One-way ETH-DM**

In this case, each MEP sends frame with one-way ETH-DM information to its peer MEP to facilitate one-way frame delay and/or one-way frame delay variation measurements at the peer MEP.

NOTE – If the clocks between the two MEPs are synchronized, one-way frame delay measurement can be carried out; otherwise, only one-way frame delay variation measurement can be performed.

The PDU used for one-way ETH-DM is 1DM, as described in clause 9.14. Frames which carry the 1DM PDU are called as 1DM frames. The same 1DM frame format can be used for proactive and on-demand one-way ETH-DM. The distinction of a proactive 1DM frame from an on-demand 1DM frame is by the value of a flag field in the 1DM frame.

#### **8.2.1.1 1DM transmission**

When configured for one-way delay measurements, a MEP periodically transmits 1DM frames with the TxTimeStampf value. A MEP can optionally use a Test ID TLV and/or a Data TLV. The MEP uses a Test ID TLV, containing a Test ID that is used to run multiple tests simultaneously, when configured. The MEP uses a Data TLV when configured for measuring delay and delay variation for different frame sizes.

#### **8.2.1.2 1DM reception**

When configured for one-way delay measurements, a MEP, upon receiving a valid 1DM frame, uses the following values to make one-way frame delay measurement. A 1DM frame with a valid MEG level and a destination MAC address equal to the receiving MEP's MAC address is considered to be a valid 1DM frame. These values serve as input to the one-way frame delay variation measurement:

- 1DM frame's TxTimeStampf value
- RxTimef, which is the time at reception of the 1DM frame

$$\text{Frame Delay}_{\text{one-way}} = \text{RxTimef} - \text{TxTimeStampf}$$

### **8.2.2 Two-way ETH-DM**

A MEP sends frames with ETH-DM request information to its peer MEP and receives frames with ETH-DM reply information from its peer MEP to carry out two-way frame delay and two-way frame delay variation measurements. If two optional timestamps of RxTimeStampf and TxTimeStampb are supported on its peer MEP, the results of one-way frame delay and one-way frame delay variation measurements can be also calculated by the same ETH-DM request/reply information.

NOTE – Regarding the one-way measurements, if the clocks between the two MEPs are synchronized, one-way frame delay measurement can be carried out. Otherwise, only one-way frame delay variation measurement can be performed.

The PDU used for ETH-DM request is DMM, as described in clause 9.15. The PDU used for ETH-DM reply is DMR, as described in clause 9.16. Frames which carry the DMM PDU are called as DMM frames. Frames which carry the DMR PDU are called as DMR frames. The same DMM and DMR frame formats can be used for proactive and on-demand two-way ETH-DM. The distinction of proactive DMM/DMR frames from on-demand DMM/DMR frames is by the value of a flag field in the DMM/DMR frames.

### 8.2.2.1 DMM transmission

When configured for two-way delay measurements, a MEP periodically transmits DMM frames with the TxTimeStampf value. A MEP can optionally use a test ID TLV and/or a data TLV. The MEP uses a test ID TLV, containing a test ID that is used to run multiple tests simultaneously, when configured. The MEP uses a data TLV when configured for measuring delay and delay variation for different frame sizes.

### 8.2.2.2 DMM reception and DMR transmission

Whenever a valid DMM frame is received by a MEP, a DMR frame is generated and transmitted to the requesting MEP. A DMM frame with a valid MEG level and a destination MAC address equal to the receiving MEP's MAC address is considered to be a valid DMM frame. Every field in the DMM frame is copied to the DMR frame with the following exceptions:

- The source MAC address is copied to the destination MAC address and the source MAC address is filled with the MEP MAC address.
- The OpCode field is changed from DMM to DMR.

NOTE – As an option, two additional timestamps may be used in the DMR frame to take into account the processing time at the remote MEP: RxTimeStampf (timestamp at the time of receiving the DMM frame) and TxTimeStampb (timestamp at the time of transmitting the DMR frame).

### 8.2.2.3 DMR reception

Upon receiving a DMR frame, a MEP uses the following values to calculate two-way frame delay. This value serves as input for two-way frame delay variation measurement:

- DMR frame's TxTimeStampf value
  - RxTimeeb – reception time of the DMR frame
- $$\text{Frame Delay}_{\text{two-way}} = \text{RxTimeeb} - \text{TxTimeStampf}$$

If the additional timestamps are carried in the DMR frame, which is determined by non-zero values of the RxTimeStampf and TxTimeStampb fields, the frame delay for one-way and two-way can be calculated to be:

$$\text{Frame Delay}_{\text{two-way}} = (\text{RxTimeeb} - \text{TxTimeStampf}) - (\text{TxTimeStampb} - \text{RxTimeStampf})$$

$$\text{Frame Delay}_{\text{one-way}_{\text{far}}} = \text{RxTimeStampf} - \text{TxTimeStampf}$$

$$\text{Frame Delay}_{\text{one-way}_{\text{near}}} = \text{RxTimeeb} - \text{TxTimeStampb}$$

## 8.3 Throughput measurement

[b-IETF RFC 2544] specifies measuring the throughput by sending frames at an increasing rate (up to the theoretical maximum), graphing the percentage of frames received, and reporting the rate at which frames start being dropped. In general this rate is dependent on the frame size.

The mechanisms specified in this Recommendation, e.g., unicast ETH-LB (e.g., LBM and LBR frames with the data field) and ETH-Test (e.g., TST frames with the data field) can be used for performing the throughput measurements. A MEP can insert TST frames or LBM frames with configured size, pattern, etc., at a rate to exercise the throughput and make one-way or two-way measurements.

## 8.4 Synthetic loss measurement (ETH-SLM)

Synthetic loss measurement is a mechanism to measure frame loss using synthetic frames, rather than data traffic. A number of synthetic frames are sent and received, and the number of those that are lost is hence calculated. This can be treated as a statistical sample, and used to approximate the frame loss ratio of data traffic.



ETH-SLM collects counters so as to maintain a count of transmitted and received synthetic frames between a set of MEPs.

ETH-SLM is used to perform on-demand or pro-active tests by sending a finite number of frames with ETH-SLM information to one or multiple peer MEPs and similarly receiving frames with ETH-SLM information from the peer MEPs. Each MEP then performs frame loss measurements which contribute to unavailable time. Since a bidirectional service is defined as unavailable if either of the two directions is declared unavailable, ETH-SLM must facilitate each MEP to perform near-end and far-end synthetic frame loss measurements.

A MEP maintains the following local counters for each test ID and for each peer MEP being monitored in a ME for which loss measurements are to be performed:

- TxFCI: number of synthetic frames transmitted towards the peer MEP, part of a given test ID. A source MEP increments this number for successive transmission of synthetic frames with ETH-SLM request information while a responder MEP increments it for successive transmission of synthetic frames with ETH-SLM reply information.
- RxFCI: number of synthetic frames received from the peer MEP and part of a given test ID. A source MEP increments this number for successive reception of synthetic frames with ETH-SLM reply information while a responder MEP increments it for successive reception of synthetic frames with ETH-SLM request information.

The method of loss measurement involving series of frames with increasing values of TxFCI with ETH-SLM information, as shown in the following subclauses.

NOTE 1 – No synchronization is required of test ID value between initiating and responding MEPs, as the test ID is configured at the initiating MEP, and the responding MEP uses the test ID it receives from the initiating MEP. The allocation and release of local counter resources for each test ID at the responding MEP is outside the scope of this Recommendation.

The specific configuration information required by a MEP to support ETH-SLM is the following:

- MEG level – MEG level at which the MEP exists.
- Data – Optional data element whose length is configurable at the MEP. The inclusion of the optional data element in the SLM frame is to support configurable SLM frame size.
- Destination MAC address – Identifies the target peer MEP.
- Test ID – Used to distinguish each SL measurement because multiple measurements can be simultaneously activated also on a given CoS and MEP pair. It must be unique at least within the context of any SL measurement for the MEG and initiating MEP.
- Priority – Identifies the priority of the frames with ETH-SLM information. This information is configurable per operation.
- Drop eligibility – Frames with ETH-SLM information are always marked as drop ineligible. This information is not necessarily configured.

A MIP is transparent to frames with ETH-SLM information and therefore does not require any information to support ETH-SLM functionality.

NOTE 2 – As ETH-SLM is a sampling technique, it is inevitably less accurate than counting the service frames. Furthermore, the accuracy depends on the number of SLM frames used or the period for transmitting SLM frames. The number of SLM frames or the period for SLM frames is outside the scope of this Recommendation, but some examples of accuracy are included for information in Appendix VI.

#### **8.4.1 Single-ended ETH-SLM**

Single-ended ETH-SLM is used for proactive or on-demand OAM. It carries out synthetic loss measurements applicable to both point-to-point ETH connection and multipoint ETH connectivity. It allows a MEP to initiate and report far-end and near-end loss measurements associated with one or a set of peer MEPs part of the same MEG.

The selection of on-demand or proactive is performed by the management function that initiates the test; however this is local information and does not need to be conveyed in the PDU.

For single-ended operation, a MEP sends frames with ETH-SLM request information to its peer MEP(s) and receives frames with ETH-SLM reply information from its peer MEP(s) to carry out synthetic loss measurements.

The PDU used for single-ended ETH-SLM request is SLM, as described in clause 9.22. The PDU used for single-ended ETH-SLM reply is SLR, as described in clause 9.23. Frames which carry the SLM PDU are called SLM frames. Frames which carry the SLR PDU are called SLR frames.

#### **8.4.1.1 SLM transmission**

A MEP periodically transmits SLM frames with the following information elements included:

- Test ID: Value containing a number configured by the MEP and then used to run multiple tests simultaneously.
- Source MEP ID: MEP's own identity in the MEG.
- TxFCf: Value of the local counter TxFCI at the time of SLM frame transmission.
- TxFCb: Set always to zero. Reserved for SLR transmission.

#### **8.4.1.2 SLM reception and SLR transmission**

Whenever a valid SLM frame is received by a MEP, an SLR frame is generated and transmitted to the requesting MEP. An SLM frame with a valid MEG level and a destination MAC address equal to the receiving MEP's MAC address is considered to be a valid SLM frame. Every field in the SLM frame is copied to the SLR frame with the following exceptions:

- The source MAC address is copied to the destination MAC address and the source MAC address is filled with the MEP MAC address.
- The OpCode is changed from SLM to SLR.
- Responder MEP ID: MEP's own identity in the MEG.
- TxFCb: Value of the local counter RxFCI at the time of SLR frame transmission.

Note that as an SLR frame is generated every time an SLM frame is received, RxFCI in the responder is equal to the number of SLM frames received and also equal to the number of SLR frames sent. In other words, in the responder, RxFCI = TxFCI.

#### **8.4.1.3 SLR reception**

After transmission of a SLM frame (with a given TxFCf value), a MEP will expect to receive a corresponding SLR frame (carrying same TxTCf value) within the timeout value from its peer MEP(s). SLR frames received after the timeout value (5 seconds) must be ignored.

With the information contained in SLR frames, a MEP determines frame loss for given measurement periods. The measurement period is a time interval during which the number of SLM frames transmitted is statistically adequate to make a measurement at a given accuracy. (See Appendix VI.) A MEP uses the following values to determine near-end and far-end frame loss in the measurement period:

- Last received SLR frame's TxFCf and TxFCb values and local counter RxFCI at the end of the measurement period. These values are represented as TxFCf[t<sub>c</sub>], TxFCb[t<sub>c</sub>] and RxFCI[t<sub>c</sub>], where t<sub>c</sub> is the end time of the measurement period.
- SLR frame's TxFCf and TxFCb values of the first received SLR frame after the test starts and local counter RxFCI at the beginning of the measurement period. These values are represented as TxFCf[t<sub>p</sub>], TxFCb[t<sub>p</sub>] and RxFCI[t<sub>p</sub>], where t<sub>p</sub> is the start time of the measurement period.

$$\text{Frame loss}_{\text{far-end}} = | \text{TxFCf}[t_c] - \text{TxFCf}[t_p] | - | \text{TxFCb}[t_c] - \text{TxFCb}[t_p] |$$

$$\text{Frame loss}_{\text{near-end}} = | \text{TxFCb}[t_c] - \text{TxFCb}[t_p] | - | \text{RxFCI}[t_c] - \text{RxFCI}[t_p] |$$

NOTE – If there are SLMs at the end of the measurement period for which no corresponding SLRs have been received within the timeout period (i.e., SLMs with sequence numbers after the sequence number of the last SLR received), it is not possible to determine whether they were lost in the near-end or far-end direction.

## 9 OAM PDU types

This clause describes the information elements and formats for different OAM PDU types used to meet the requirements of OAM functions described in clauses 7 and 8.

NOTE – When the values of OAM PDU fields are fixed, they are shown in parentheses in the OAM PDU formats in the following sub-clauses.

### 9.1 Common OAM information elements

Some information elements are common across the OAM PDUs that are identified in this Recommendation. These information elements are:

- MEG Level: MEG Level is a 3-bit field. It contains an integer value that identifies MEG level of OAM PDU. Value ranges from 0 to 7.
- Version: Version is a 5-bit field. It contains an integer value that identifies the OAM protocol version.
- OpCode: OpCode is a 1-octet field. It contains an OpCode that identifies an OAM PDU type. OpCode is used to identify the remaining content of an OAM PDU. The values of this information field are shown in Table 9-1.
- Flags: Flags is an 8-bit field. Use of the bits in this field is dependent on the OAM PDU type.
- TLV Offset: TLV Offset is a 1-octet field. It contains the offset to the first TLV in an OAM PDU relative to the TLV Offset field. The value of this field is associated with an OAM PDU type. When the TLV Offset is 0, it points to the first octet following the TLV Offset field.

Other information elements which are not present in OAM PDUs but are conveyed in frames carrying OAM PDUs include:

- Priority: Priority identifies the priority of a specific OAM frame.
- Drop eligibility: Drop eligibility identifies the drop eligibility of a specific OAM frame.

**Table 9-1 – OpCode values**

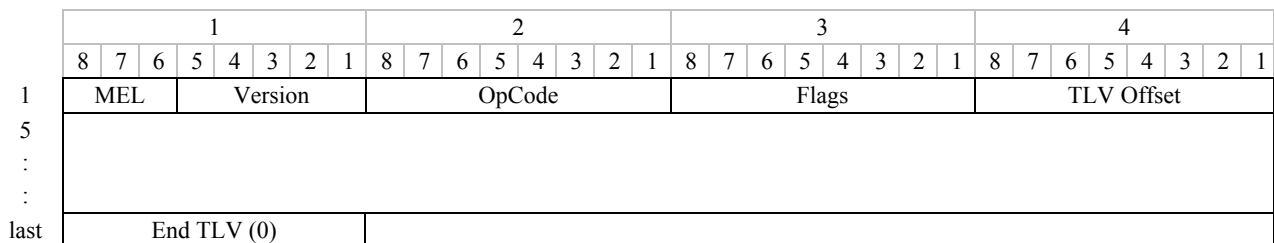
OpCode value	OAM PDU type	OpCode relevance for MEPs/MIPs
OpCodes common with IEEE 802.1		
1	CCM	MEPs
3	LBM	MEPs and MIPs (connectivity verification)
2	LBR	MEPs and MIPs (connectivity verification)
5	LTM	MEPs and MIPs
4	LTR	MEPs and MIPs
0, 6-31, 64-255	Reserved (Note 1)	

**Table 9-1 – OpCode values**

OpCode value	OAM PDU type	OpCode relevance for MEPs/MIPs
OpCodes specific to this Recommendation		
33	AIS	MEPs
35	LCK	MEPs
37	TST	MEPs
39	Linear APS	MEPs
40	Ring APS	MEPs
41	MCC	MEPs
43	LMM	MEPs
42	LMR	MEPs
45	IDM	MEPs
47	DMM	MEPs
46	DMR	MEPs
49	EXM	Outside the scope of this Recommendation
48	EXR	Outside the scope of this Recommendation
51	VSM	Outside the scope of this Recommendation
50	VSR	Outside the scope of this Recommendation
52	CSF	MEPs
55	SLM	MEPs
54	SLR	MEPs
32, 34, 36, 38, 44, 53, 56-63	Reserved (Note 2)	
NOTE 1 – Reserved for definition by IEEE 802.1.		
NOTE 2 – Reserved for future standardization by ITU-T.		

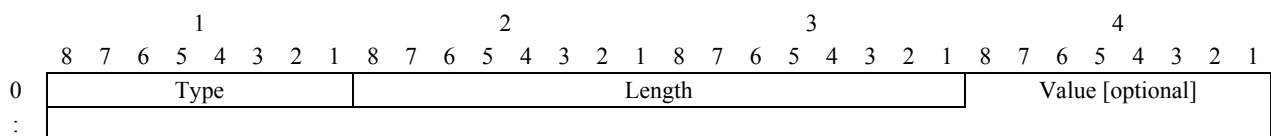
**9.1.1 Common OAM PDU format**

The common format used in all OAM PDUs is shown in Figure 9.1-1.



**Figure 9.1-1 – Common OAM PDU format**

The general format of TLVs is shown in Figure 9.1-2. Type values are specified in Table 9-2.



**Figure 9.1-2 – Generic TLV format**

NOTE – In an End TLV, Type = 0, and both Length and Value fields are not used.

**Table 9-2 – Type values**

Type value	TLV name
Types common with IEEE 802.1	
0	End TLV
3	Data TLV
5	Reply ingress TLV
6	Reply egress TLV
7	LTM egress identifier TLV
8	LTR egress identifier TLV
2, 4, 9-31, 64-255	Reserved (Note 1)
Types specific to this Recommendation	
32	Test TLV
36	Test ID TLV
33-35, 37-63	Reserved (Note 2)
NOTE 1 – Reserved for definition by IEEE 802.1.	
NOTE 2 – Reserved for future standardization by ITU-T.	

## 9.2 CCM PDU

CCM is used to support ETH-CC function, as described in clause 7.1, ETH-RDI function, as described in clause 7.5, and dual-ended ETH-LM function, as described in clause 8.1.1.

### 9.2.1 CCM information elements

Information elements carried in CCM to support ETH-CC are:

- **Period:** Period is a 3-bit information element carried in the three least significant bits of the Flags field. Period contains the value of the CCM transmission period configured at the CCM source. CCM Period values are specified in Table 9-3.
- **MEG ID:** MEG ID is a 48-octet field which contains the MEG ID of the MEG to which the MEP transmitting the CCM frame belongs. See Annex A.
- **MEP ID:** MEP ID is a 2-octet field where the 13 least significant bits are used to identify the MEP transmitting the CCM frame. MEP ID is unique within the MEG.

Information element carried in CCM to support ETH-RDI is:

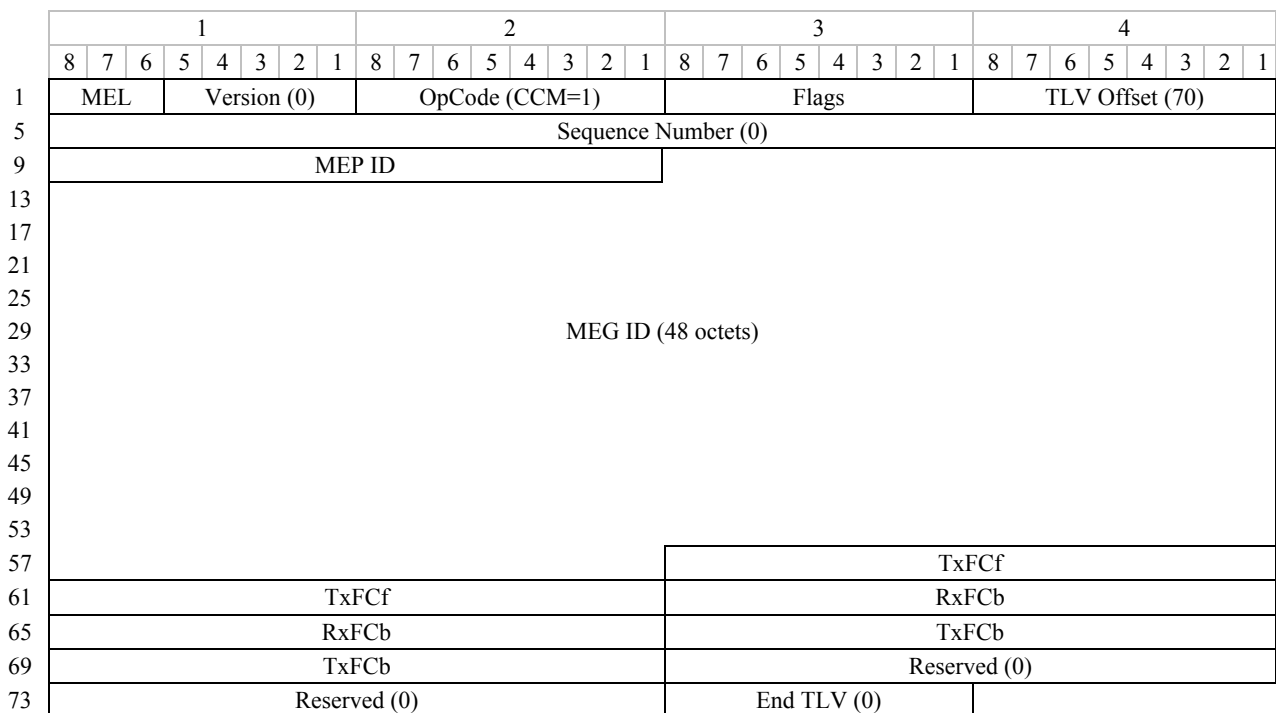
- **RDI:** RDI is a 1-bit information element carried in most significant bit of Flags field. When the RDI bit is 1, detection of a defect is indicated by the transmitting MEP. When the RDI bit is 0, no defect indication is communicated by the transmitting MEP.

The information elements carried in CCM to support dual-ended ETH-LM are:

- TxFCf: TxFCf is a 4-octet field which carries the value of the counter of in-profile data frames transmitted by the MEP towards its peer MEP, at the time of CCM frame transmission.
- RxFCb: RxFCb is a 4-octet field which carries the value of the counter of in-profile data frames received by the MEP from its peer MEP, at the time of receiving the last CCM frame from that peer MEP.
- TxFCb: TxFCb is a 4-octet field which carries the value of the TxFCf field in the last CCM frame received by the MEP from its peer MEP.

### 9.2.2 CCM PDU format

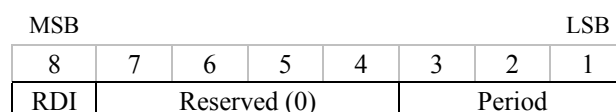
CCM PDU format used by a MEP to transmit CCM information is shown in Figure 9.2-1.



**Figure 9.2-1 – CCM PDU format**

The fields of the CCM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is CCM (1).
- Flags: Two information elements in the Flags field for CCM PDU: RDI, and Period as follows:



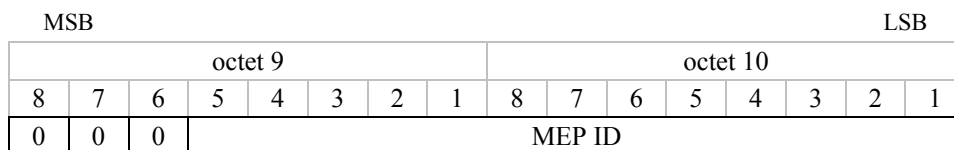
**Figure 9.2-2 – Flags format in CCM PDU**

- RDI: Bit 8 is set to 1 to indicate RDI, otherwise it is set to 0.
- Period: Bits 3 to 1 indicate the transmission period with the encoding shown in Table 9-3.

**Table 9-3 – CCM period values**

Flags[3:1]	Period value	Comments
000	Invalid value	Invalid value for CCM PDUs
001	3.33 ms	300 frames per second
010	10 ms	100 frames per second
011	100 ms	10 frames per second
100	1 s	1 frame per second
101	10 s	6 frames per minute
110	1 min	1 frame per minute
111	10 min	6 frame per hour

- TLV offset: Set to 70.
- Sequence Number: This field is set to all-ZEROes for this Recommendation.
- MEP ID: A 13-bit integer value identifying the transmitting MEP within the MEG. The three MSBs of the first octet are not used and set to ZERO:



**Figure 9.2-3 – MEP ID format in CCM PDU**

- MEG ID: 48-octet field. Refer to Annex A for the format used for the MEG ID field.
- TxFCf, TxFCb, RxFCb: 4-octet integer values with samples of the wrap-around frame counters, as specified in clause 9.2.1. These fields are set to all-ZEROes when not used.
- Reserved: Reserved fields are set to all ZEROes.
- End TLV: An all-ZEROes octet value.

### 9.3 LBM PDU

LBM is used to support ETH-LB request, as described in clause 7.2.

#### 9.3.1 LBM information elements

Information elements carried in LBM include:

- Transaction ID/Sequence Number: Transaction ID/Sequence Number is a 4-octet field that contains the transaction ID/sequence number for the LBM. The receiver is expected to copy the Transaction ID /Sequence Number in the LBR PDU, as described in clause 9.4.





The fields of the Data TLV format are as follows:

- Type: Identifies TLV type; value for this TLV type is Data (3).
- Length: Identifies size, in octets, of the Value field containing the Data Pattern. In a frame where the PDU is limited to 1492 octets, the maximum length value is 1480 (since 12 bytes are required for 8 octets of LBM PDU overhead, 3 octets of Data TLV overhead, and 1 octet of End TLV). Any other TLVs, if present in LBM, will furthermore detract from the maximum length value of 1480.
- Data Pattern: An n-octet (n = Length) arbitrary bit pattern. The receiver should ignore it.

1								2								3								4							
8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
Type (32)								Length								Pattern Type															
Test Pattern (NULL, PRBS)																															
CRC-32 (optional)																															

**Figure 9.3-4 – Test TLV format**

The fields of the Test TLV format are as follows:

- Type: Identifies TLV type; value for this TLV type is Test (32).
- Length: Identifies size, in octets, of the Value field containing the pattern type, test pattern and CRC-32. In a frame where the PDU is limited to 1492 octets, the maximum length value is 1480 octets (since 12 bytes are required for 8 octets of LBM PDU overhead, 3 octets of Test TLV overhead, and 1 octet of End TLV). Any other TLVs, if present in LBM, will furthermore detract from the maximum length value of 1480. (As one byte is used for pattern type, 1479 bytes are available for the test pattern.)
- Pattern Type: Identifies test pattern type; values are:
  - 0 'Null signal without CRC-32'
  - 1 'Null signal with CRC-32'
  - 2 'PRBS  $2^{31}-1$  without CRC-32'
  - 3 'PRBS  $2^{31}-1$  with CRC-32'
  - 4-255 Reserved for future standardization
- Test Pattern: An n-octet (n ≤ Length) test pattern: PRBS  $2^{31}-1$  or Null (all-zeroes) pattern.
- CRC-32: covers all fields (from Type to last octet before CRC-32)

## 9.4 LBR PDU

LBR is used to support ETH-LB reply, as described in clause 7.2.

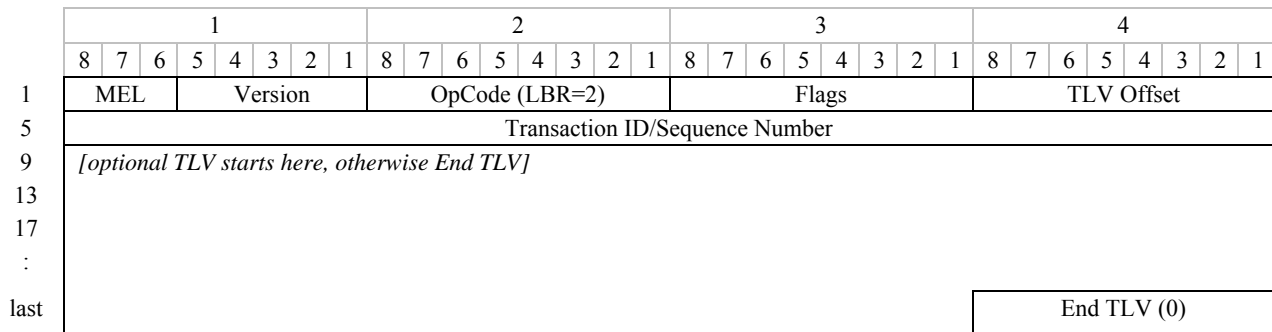
### 9.4.1 LBR information element

The information elements carried in LBR include:

- Transaction ID/Sequence Number: Transaction ID/Sequence Number is a 4-octet field that is copied from the Transaction ID/Sequence Number field in LBM.
- Data: Data is a field that is copied from the Data field in LBM.

## 9.4.2 LBR PDU format

The LBR PDU format used by a MEP to transmit LBR information is shown in Figure 9.4-1.



**Figure 9.4-1 – LBR PDU format**

The fields for the LBR PDU format are as follows:

- MEG Level: A 3-bit field the value of which is copied from the received LBM PDU.
- Version: A 5-bit field the value of which is copied from the LBM PDU.
- OpCode: Value for this PDU type is LBR (2).
- Flags: A 1-octet field the value of which is copied from the LBM PDU.
- TLV offset: A 1-octet field the value of which is copied from the LBM PDU.
- Transaction ID/Sequence Number: A 4-octet field the value of which is copied from the LBM PDU.
- Optional TLV: If present in LBM PDU, are copied from the LBM PDU.
- End TLV: A 1-octet field the value of which is copied from the LBM PDU.

## 9.5 LTM PDU

LTM is used to support ETH-LT request, as described in clause 7.3.

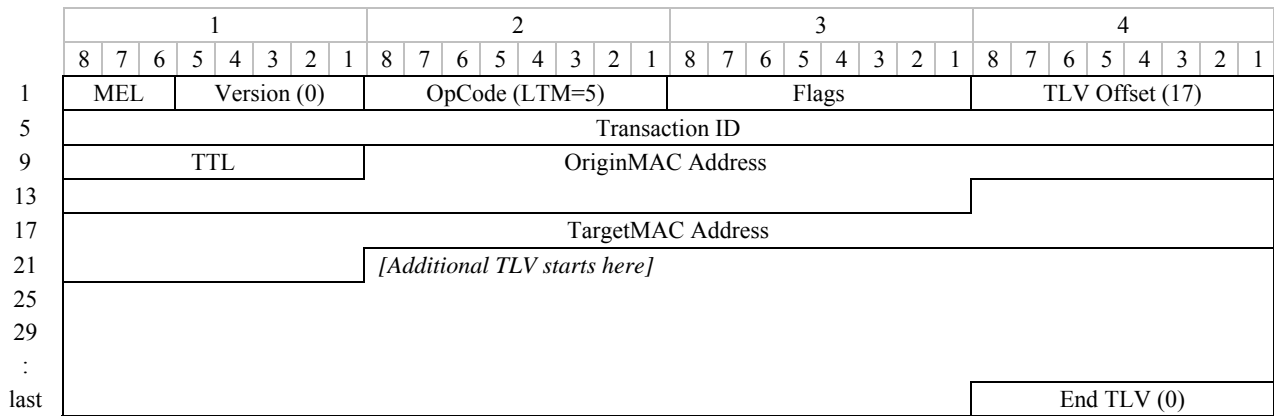
### 9.5.1 LTM information elements

The information elements carried in LTM include:

- Transaction ID: Transaction ID is a 4-octet field that contains the transaction number for the LTM. The receiver is expected to copy Transaction ID in the LTR PDU, as described in clause 9.6.
- TTL: TTL is a 1-octet field used to indicate whether a LTM should be terminated or not by the receiver. When a MIP receives LTM with TTL=1, the LTM is not relayed. A network element receiving LTM decrements the received TTL value by one and copies it into the TTL field of LTR PDU, as described in clause 9.6, as well as into the LTM that it forwards towards the next hop.
- TargetMAC: TargetMAC is a 6-octet field used to carry MAC address of the targeted end-point. An intermediate MIP copies this field into the LTM that it forwards towards the next hop.
- OriginMAC: OriginMAC is a 6-octet field used to carry the MAC address of the originating MEP. An intermediate MIP copies this field into the LTM that it forwards towards the next hop.

### 9.5.2 LTM PDU format

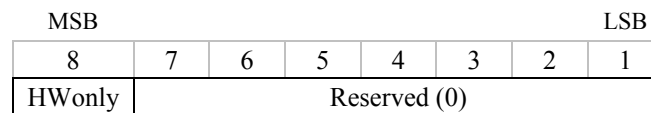
The LTM PDU format used by a MEP to transmit LTM information is shown in Figure 9.5-1.



**Figure 9.5-1 – LTM PDU format**

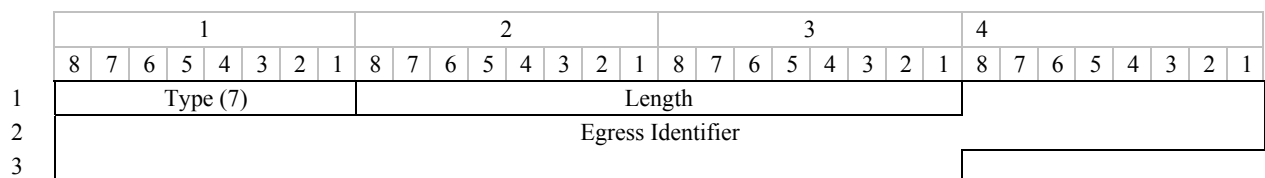
The fields of the LTM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: The value of this PDU type is LTM (5).
- Flags: The format is as shown in Figure 9.5-2.



**Figure 9.5-2 – Flags format in LTM PDU**

- HWonly: Bit 8 set to 1. Value 1 indicates that only MAC addresses learned in a bridge's active data forwarding tables is to be used to forward the LTM to the next hop. When forwarding a received LTM, HWonly is copied from incoming LTM value.
- TLV offset: Set to 17.
- Transaction ID: A 4-octet value containing the transaction ID for the LTM PDU.
- TTL: 1-octet field used to carry a TTL value as specified in clause 9.5.1.
- OriginMAC Address: A 6-octet OriginMAC as specified in clause 9.5.1.
- TargetMAC Address: A 6-octet TargetMAC as specified in clause 9.5.1.
- Additional TLV: LTM Egress Identifier TLV as specified in Figure 9.5-3.
- End TLV: All-ZEROes octet value.



**Figure 9.5-3 – LTM egress identifier TLV format**

The fields of the LTM egress identifier TLV format are as follows:

- Type: Identifies TLV type; value for this TLV type is LTM egress identifier (7).

- Length: Identifies size, in octets, of the Value field containing the egress identifier. This is set to 8.
- Egress Identifier: Identifies MEP initiating LTM frame or ETH-LT responder relaying modified LTM frame. Octets 4 and 5 are ZEROs while remaining six Octets 6-11 contain a 48-bit IEEE MAC address unique to network element where the MEP or ETH-LT responder resides.

## 9.6 LTR PDU

LTR is used to support ETH-LT reply, as described in clause 7.3.

### 9.6.1 LTR information elements

The information elements carried in LTR include:

- Transaction ID: Transaction ID is a 4-octet field that is copied from the Transaction ID field in LTM.
- TTL: TTL is a 1-octet field that contains the TTL field value decremented by 1 from the LTM for which LTR is being sent.

### 9.6.2 LTR PDU format

The LTR PDU format used by a MEP or MIP to transmit LTR information is shown in Figure 9.6-1.

		1				2				3				4											
		8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1		MEL			Version (0)				OpCode (LTR=4)				Flags				TLV Offset (6)								
5		Transaction ID																							
9		TTL				Relay Action				[TLVs starts here]															
17																									
21																									
:																									
last																						End TLV (0)			

**Figure 9.6-1 – LTR PDU format**

The fields of the LTR PDU format are as follows:

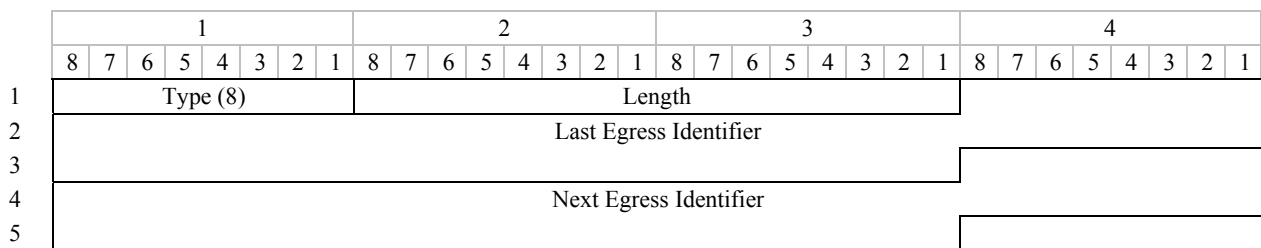
- MEG Level: A 3-bit field the value of which is copied from the received LTM PDU.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value of this PDU type is LTR (4).
- Flags: The format is as shown in Figure 9.6-2.

MSB				LSB				
8	7	6	5	4	3	2	1	
HwOnly	FwdYes	TerminalMEP	Reserved (0)					

**Figure 9.6-2 – Flags format in LTR PDU**

- HwOnly: Bit 8 (HwOnly) is copied from incoming LTM value.
- FwdYes: Bit 7 is set to 1 if modified LTM frame was relayed, or set to 0 if no LTM frame was relayed.
- TerminalMEP: Bit 6 is set to 1 if reply egress TLV (or reply ingress TLV, if the reply egress TLV is not present) is a MEP, or set to 0 otherwise.

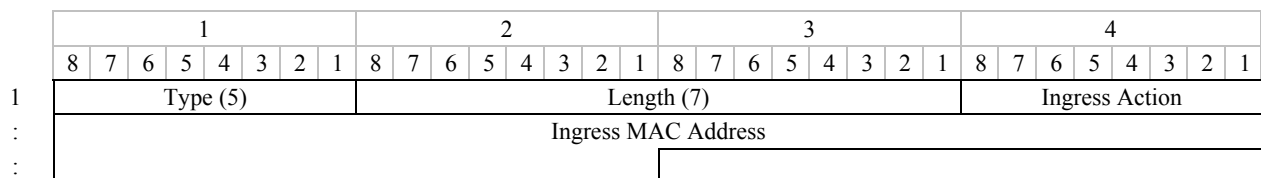
- TLV offset: Set to 6.
- Transaction ID: A 4-octet field the value of which is copied from the LTM PDU.
- TTL: A 1-octet field the value of which is copied from the LTM PDU after decrementing it by one.
- Relay Action: A 1-octet field that reports how the data frame targeted by the LTM would be passed through the MAC relay entity to the egress bridge port as described in 21.9.5 in [IEEE 802.1ag]. The value is defined in Table 21-27 of [IEEE 802.1ag].
- TLVs: LTR egress identifier TLV, reply ingress TLV and/or reply egress TLV as specified in Figures 9.6-3, 9.6.4 and 9.6-5 respectively.
- End TLV: All-ZEROes octet value.



**Figure 9.6-3 – LTR egress identifier TLV format**

The fields of the LTR egress identifier TLV format are as follows:

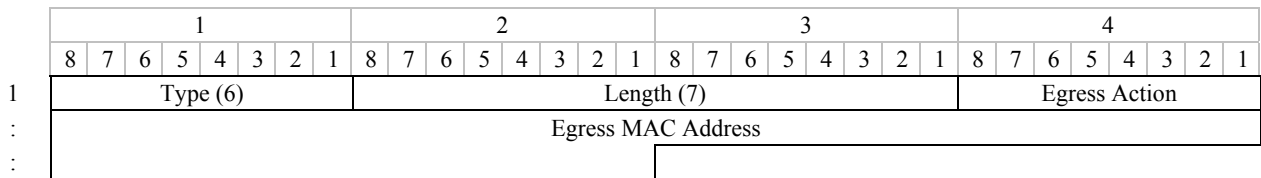
- Type: Identifies TLV type; value for this TLV type is LTR egress identifier (8).
- Length: Identifies size, in octets, of the Value field containing the last egress identifier and next egress identifier. This is set to 16.
- Last Egress Identifier: Identifies MEP that initiated, or ETH-LT responder that relayed the LTM frame to which this LTR frame is the response. This field is the same as the egress identifier in the LTM egress identifier TLV of the incoming LTM frame.
- Next Egress Identifier: Identifies ETH-LT Responder that transmitted this LTR frame, and which can relay a modified LTM frame to the next hop. If the FwdYes bit of Flags field is 0, the contents of this field are undefined, and ignored by the LTR frame receiver. When not undefined, Octets 12 and 13 are ZEROs while the remaining six octets 14-19 contain a 48-bit IEEE MAC address unique to network element where the ETH-LT responder resides.



**Figure 9.6-4 – Reply ingress TLV format**

The fields of the reply ingress TLV format are as follows:

- Type: Identifies TLV type; value for this TLV type is ingress reply (5).
- Length: Identifies size, in octets, of the Value field. This is set to 7.
- Ingress Action: A 1-octet field which is reserved for definition by IEEE 802.1.
- Ingress MAC Address: A 6-octet field which is reserved for definition by IEEE 802.1.



**Figure 9.6-5 – Reply egress TLV format**

The fields of the reply egress TLV format are as follows:

- Type: Identifies TLV type; value for this TLV type is egress reply (6).
- Length: Identifies size, in octets, of the Value field. This is set to 7.
- Egress Action: A 1-octet field which is reserved for definition by IEEE 802.1.
- Egress MAC Address: A 6-octet field which is reserved for definition by IEEE 802.1.

### 9.7 AIS PDU

The AIS PDU is used to support the ETH-AIS function, as described in clause 7.4.

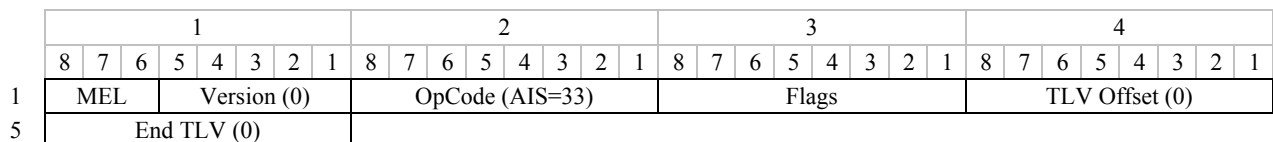
#### 9.7.1 AIS information elements

The information element carried in AIS is:

- Period: Period is a 3-bit information element carried in the three least significant bits of the Flags field. Period contains the value of AIS transmission periodicity. AIS period values are specified in Table 9-4.

#### 9.7.2 AIS PDU format

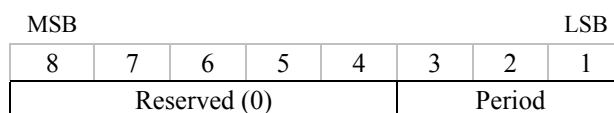
The AIS PDU format used by a MEP to transmit AIS information is shown in Figure 9.7-1.



**Figure 9.7-1 – AIS PDU format**

The fields of the AIS PDU format are as follows:

- MEG Level: A 3-bit field that is used to carry the MEG Level of the client MEG.
- Version: refer to sub-clause 9.1, value is always 0.
- OpCode: Value for this PDU type is AIS (33).
- Flags: One information element in the Flags field for the AIS PDU, Period, as follows:



**Figure 9.7-2 – Flags format in AIS PDU**

- Period: Bits 3 to 1 indicate transmission period with the encoding in Table 9-4.

**Table 9-4 – AIS/LCK period values**

Flags[3:1]	Period value	Comments
000-011	Invalid value	Invalid value for AIS/LCK PDUs
100	1s	1 frame per second
101	Invalid value	Invalid value for AIS/LCK PDUs
110	1 min	1 frame per minute
111	Invalid value	Invalid value for AIS/LCK PDUs

- TLV offset: Set to 0.
- End TLV: All-ZEROes octet value.

## 9.8 LCK frame

LCK PDU is used to support ETH-LCK function, as described in clause 7.6.

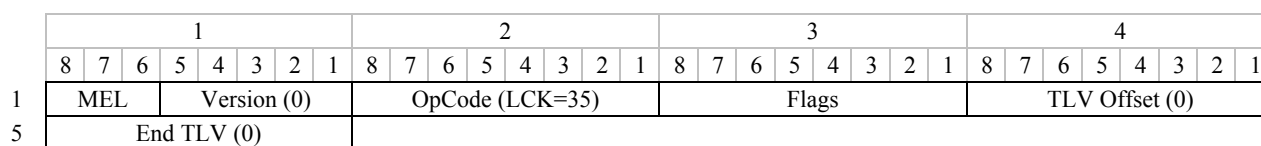
### 9.8.1 LCK information elements

The information element carried in LCK is:

- Period: Period is a 3-bit information element carried in the three least significant bits of the Flags field. Period contains the value of LCK transmission periodicity. LCK period values are specified in Table 9-4.

### 9.8.2 LCK PDU format

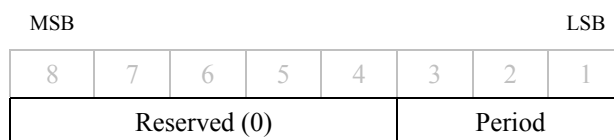
The LCK PDU format used by a MEP to transmit LCK information is shown in Figure 9.8-1.



**Figure 9.8-1 – LCK PDU format**

The fields of the LCK PDU format are as follows:

- MEG Level: A 3-bit field that is used to carry the MEG level of the client MEG.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is LCK (35).
- Flags: One information element in the Flags field for the LCK PDU, Period, as follows:



**Figure 9.8-2 – Flags format in LCK PDU**

- Period: Bits 3 to 1 indicate transmission period with the encoding in Table 9-4.
- TLV offset: Set to 0.
- End TLV: All-ZEROes octet value.

## 9.9 TST PDU

TST PDU is used to support the unidirectional ETH-Test function, as described in clause 7.7.

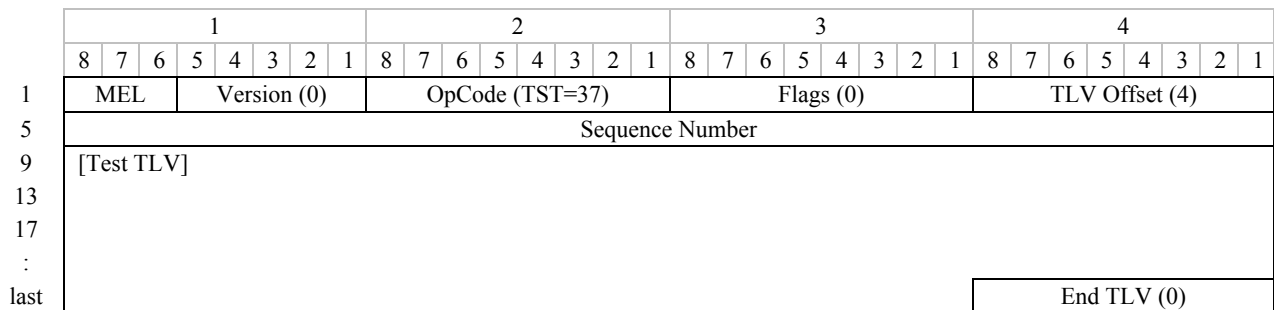
### 9.9.1 TST information elements

The information elements carried in TST are:

- Sequence Number: Sequence Number is a 4-octet field that contains the sequence number for the TST frames.
- Test: Test is an optional field whose length and contents are determined at the transmitting MEP. The contents of Test field indicate a test pattern and also carry an optional checksum. The test pattern can be a pseudo-random bit sequence (PRBS) ( $2^{31}-1$ ) as specified in clause 5.8 of [ITU-T O.150], an all '0' pattern, etc.

### 9.9.2 TST PDU format

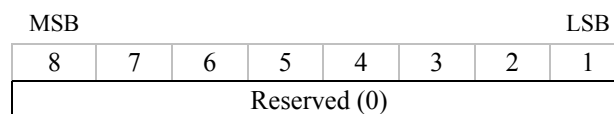
The TST PDU format used by a MEP to transmit TST information is shown in Figure 9.9-1.



**Figure 9.9-1 – TST PDU format**

The fields of the TST PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is TST (37).
- Flags: Set to all-ZEROes.



**Figure 9.9-2 – Flags format in TST PDU**

- TLV offset: Set to 4.
- Sequence Number: A 4-octet value containing the sequence number incremented for successive TST PDUs.
- Test TLV: Test TLV as specified in Figure 9.3-4.
- End TLV: All-ZEROes octet value.

## 9.10 APS PDU

APS is used to support ETH-APS function, as described in clause 7.8.

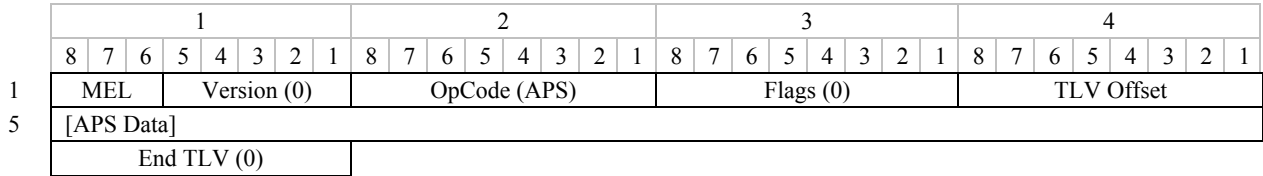


### 9.10.1 APS information elements

Information elements carried in APS are outside the scope of this Recommendation.

### 9.10.2 APS PDU format

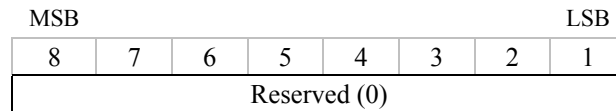
The APS PDU format used by a MEP to transmit APS information is shown in Figure 9.10-1.



**Figure 9.10-1 – APS PDU format**

The fields of the APS PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is (39) for linear APS and (40) for ring APS.
- Flags: Set to all-ZEROes.



**Figure 9.10-2 – Flags format in APS PDU**

- TLV offset: 1-octet field. Its specific value for APS is outside the scope of this Recommendation.
- APS Data: Format and length of this field are outside the scope of this Recommendation.
- End TLV: All-ZEROes octet value.

## 9.11 MCC PDU

MCC PDU is used to support ETH-MCC, as described in clause 7.9.

### 9.11.1 MCC Information Elements

Information elements carried in MCC include:

- OUI: OUI is a 3-octet field that contains the organizationally unique identifier of the organization defining the format of MCC Data and values SubOpCode.
- SubOpCode: SubOpCode is a 1-octet field that is used to interpret the remaining fields in the MCC PDU.
- MCC Data: Depending on the functionality indicated by the OUI and organizationally specific SubOpCode, MCC may carry one or more TLVs. MCC Data is outside the scope of this Recommendation.

### 9.11.2 MCC PDU format

The MCC PDU format used by a MEP to transmit MCC information is shown in Figure 9.11-1.

		1								2								3								4							
		8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1		MEL		Version (0)				OpCode (MCC=41)				Flags (0)				TLV Offset																	
5		OUI																SubOpCode															
9		<i>[optional MCC data; else End TLV]</i>																															
:																																	
:																																	
Last		End TLV (0)																															

**Figure 9.11-1 – MCC PDU format**

The fields of the MCC PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is MCC (41).
- Flags: Set to all-ZEROes.

MSB				LSB			
8	7	6	5	4	3	2	1
Reserved (0)							

**Figure 9.11-2 – Flags format in MCC PDU**

- TLV offset: 1-byte field. Its specific value for MCC is outside the scope of this Recommendation.
- OUI: 3-octet field the values of which are outside the scope of this Recommendation.
- SubOpCode: 1-octet field the values of which are outside the scope of this Recommendation.
- MCC Data: Format and length of this field are outside the scope of this Recommendation.
- End TLV: All-ZEROes octet value.

## 9.12 LMM PDU

LMM is used to support single-ended ETH-LM request, as described in clause 8.1.2.

### 9.12.1 LMM information elements

Information elements carried in LMM are:

- TxFCf: TxFCf is a 4-octet field which carries the value of counter responsible for counting in-profile data frames transmitted by the MEP towards its peer MEP, at the time of LMM frame transmission.

### 9.12.2 LMM PDU format

LMM PDU format used by a MEP to transmit LMM information is shown in Figure 9.12-1.

	1								2								3								4							
	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1	MEL				Version (0)				OpCode (LMM=43)				Flags (0)				TLV Offset (12)															
5	TxFCf																															
9	Reserved for RxFCf in LMR																															
13	Reserved for TxFCb in LMR																															
17	End TLV (0)																															

**Figure 9.12-1 – LMM PDU format**

The fields of the LMM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is LMM (43).
- Flags: Set to all-ZEROes.

MSB				LSB			
8	7	6	5	4	3	2	1
Reserved (0)							

**Figure 9.12-2 – Flags format in LMM PDU**

- TLV offset: Set to 12.
- TxFCf: 4-octet integer values with samples of the frame counters, as specified in clause 9.12.1.
- Reserved: Reserved fields are set to all ZEROes.
- End TLV: An all-ZEROes octet value.

### 9.13 LMR PDU

LMR PDU is used to support single-ended ETH-LM reply, as described in clause 8.1.2.

#### 9.13.1 LMR information elements

Information elements carried in LMR are:

- TxFCf: TxFCf is a 4-octet field which carries the value of the TxFCf field in the last LMM PDU received by the MEP from its peer MEP.
- TxFCb: TxFCb is a 4-octet field which carries the value of the counter of in-profile data frames transmitted by the MEP towards its peer MEP at the time of LMR frame transmission.
- RxFCf: RxFCf is a 4-octet field which carries the value of the counter of in-profile data frames received by the MEP from its peer MEP, at the time of receiving last LMM frame from that peer MEP.

#### 9.13.2 LMR PDU format

The LMR PDU format used by a MEP to transmit LMR information is shown in Figure 9.13-1.

	1								2								3								4							
	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1	MEL			Version					OpCode (LMR=42)					Flags					TLV Offset													
5	TxFCf																															
9	RxFCf																															
13	TxFCb																															
17	End TLV (0)																															

**Figure 9.13-1 – LMR PDU format**

The fields of the LMR PDU format are as follows:

- MEG Level: A 3-bit field the value of which is copied from the last received LMM PDU.
- Version: A 5-bit field the value of which is copied from the last received LMM PDU.
- OpCode: Value for this PDU type is LMR (42).
- Flags: A 1-octet field the value of which is copied from the last received LMM PDU.
- TLV offset: A 1-octet field the value of which is copied from the last received LMM PDU.
- TxFCf: 4-octet field the value of which is copied from the last received LMM PDU.
- RxFCf: 4-octet integer values with samples of the frame counters, as specified in clause 9.13.1.
- TxFCb: 4-octet integer values with samples of the frame counters, as specified in clause 9.13.1.
- End TLV: A 1-octet field the value of which is copied from the LMM PDU.

### 9.14 1DM PDU

The 1DM PDU is used to support one-way proactive and on-demand ETH-DM, as described in clause 8.2.1.

#### 9.14.1 1DM information element

The information element carried in 1DM is:

- TxTimeStampf: TxTimeStampf is an 8-octet field that contains the timestamp of 1DM transmission. The format of TxTimeStampf is equal to the TimeRepresentation format in [IEEE 1588].

#### 9.14.2 1DM PDU format

The 1DM PDU format used by a MEP to transmit 1DM information is shown in Figure 9.14-1.

	1								2								3								4							
	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1	MEL			Version (1)					OpCode (1DM=45)					Flags					TLV Offset (16)													
5	TxTimeStampf																															
9	Reserved for 1DM receiving equipment (0) (for RxTimeStampf)																															
13																																
17																																
21	<i>[optional TLV starts here; otherwise End TLV]</i>																															
25																																
29																																
:																																
last																									End TLV (0)							

**Figure 9.14-1 – 1DM PDU format**

The fields of the 1DM PDU format are as follows:

- MEG Level: refer to clause 9.1.
- Version: refer to clause 9.1, value for the 1DM PDU on this version is set to 1.
- OpCode: Value for this PDU type is 1DM (45).
- Flags: One information element in the Flags field, the LSB bit (Type), is used to indicate the type of the 1DM operation as follows:

MSB							LSB
8	7	6	5	4	3	2	1
Reserved (0)							Type

**Figure 9.14-2 – Flags format in 1DM PDU**

- Type: Bit 1 is set to 1 if it is the proactive operation, or set to 0 if it is the on-demand operation.
- TLV offset: Set to 16.
- TxTimeStampf: An 8-octet transmit time stamp field as described in clause 9.14.1.
- Reserved: Field for RxTimeStampf: An 8-octet reserved fields are set to all-ZEROes.
- OptionalTLV: If present, a Test ID TLV as specified in Figure 9.14-3 and/or a data TLV as specified in Figure 9.3-3, with configurable size, in octets. When test ID TLV is included in this area, it is recommended to put test ID TLV first (prior to data TLV). For the purpose of ETH-DM, the value part of data TLV is unspecified.
- End TLV: All-ZEROes octet value.

	1								2								3								4							
	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1	Type (36)								Length																Test ID							
5	Test ID																															

**Figure 9.14-3 – Test ID TLV format**

The fields of the Test ID TLV format are as follows:

- Type: Identifies TLV type; value for this TLV type is Test ID (36).
- Length: Identifies size. Must be 32.
- Test ID: Test ID is a 4-octet field set by the transmitting MEP when used to run multiple tests simultaneously between MEPs.

## 9.15 DMM PDU

DMM is used to support two-way proactive or on-demand ETH-DM request, as described in clause 8.2.2.

### 9.15.1 DMM information elements

The information elements carried in DMM are:

- TxTimeStampf: TxTimeStampf is an 8-octet field that contains the timestamp of DMM transmission. The format of TxTimeStampf is equal to the TimeRepresentation format in [IEEE 1588].

### 9.15.2 DMM PDU Format

DMM PDU format used by a MEP to transmit DMM information is shown in Figure 9.15-1.

	1								2								3								4							
	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1	MEL		Version (1)						OpCode (DMM=47)						Flags						TLV Offset (32)											
5	TxTimeStampf																															
9																																
13	Reserved for DMM receiving equipment (0)																															
17	<i>(for RxTimeStampf)</i>																															
21	Reserved for DMR (0)																															
25	<i>(for TxTimeStampb)</i>																															
29	Reserved for DMR receiving equipment (0)																															
33																																
37	<i>[optional TLV starts here; otherwise End TLV]</i>																															
41																																
45																																
:																																
last																															End TLV (0)	

**Figure 9.15-1 – DMM PDU format**

The fields of the DMM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value for the DMM PDU is set to 1.
- OpCode: Value for this PDU type is DMM (47).
- Flags: Set to all-ZEROes. One information elements in the Flags field, the LSB bit (Type), is used to indicate the type of the DMM operation as follows:

MSB							LSB
8	7	6	5	4	3	2	1
Reserved (0)							Type

**Figure 9.15-2 – Flags format in DMM PDU**

- Type: Bit 1 is set to 1 if it is the proactive operation, or set to 0 if it is the on-demand operation.
- TLV offset: Set to 32.
- TxTimeStampf: An 8-octet transmit time stamp field as described in clause 9.15.1.
- Reserved: A 24-octet reserved fields are set to all-ZEROes.
- Optional TLV: If present, a test ID TLV as specified in Figure 9.14-3 and/or a data TLV as specified in Figure 9.3-3, with configurable size, in octets. When Test ID TLV is included in this area, it is recommended to put test ID TLV first (prior to Data TLV). For the purpose of ETH-DM, the value part of data TLV is unspecified.
- End TLV: An all-ZEROes octet value.

### 9.16 DMR PDU

DMR is used to support two-way ETH-DM reply, as described in clause 8.2.2.

### 9.16.1 DMR information elements

The information elements carried in DMR are:

- TxTimeStampf: TxTimeStampf is an 8-octet field that contains the copy of TxTimeStampf field in received DMM.
- RxTimeStampf: RxTimeStampf is an optional 8-octet field that contains the timestamp of DMM reception. The format of RxTimeStampf is equal to the TimeRepresentation format in [IEEE 1588]. When not used, a value of all 0 is used.
- TxTimeStampb: TxTimeStampb is an optional 8-octet field that contains the timestamp of DMR transmission. The format of TxTimeStampb is equal to the TimeRepresentation format in [IEEE 1588]. When not used, a value of all 0 is used.

### 9.16.2 DMR PDU format

The DMR PDU format used by a MEP to transmit DMR information is shown in Figure 9.16-1.

		1								2								3								4							
		8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1		MEL				Version				OpCode (DMR=46)				Flags				TLV Offset															
5		TxTimeStampf																															
9		RxTimeStampf																															
13		TxTimeStampb																															
17		Reserved for DMR receiving equipment (0)																															
21		<i>(for RxTimeStampb)</i>																															
25		<i>[optional TLV starts here; otherwise End TLV]</i>																															
29																																	
33																																	
37																																	
41																																	
45																																	
:																																	
last																										End TLV (0)							

**Figure 9.16-1 – DMR PDU format**

The fields of the DMR PDU format are as follows:

- MEG Level: A 3-bit field the value of which is copied from the last received DMM PDU.
- Version: A 5-bit field the value of which is copied from the last received DMM PDU.
- OpCode: Value for this PDU type is DMR (46).
- Flags: A 1-octet field the value of which is copied from the last received DMM PDU.
- TLV offset: A 1-octet field the value of which is copied from the last received DMM PDU.
- TxTimeStampf: An 8-octet field the value of which is copied from last received DMM PDU.
- RxTimeStampf: An 8-octet transmit time stamp field as described in clause 9.16.1.
- TxTimeStampb: An 8-octet transmit time stamp field as described in clause 9.16.1.
- Reserved: Reserved fields are set to all ZEROes.
- Optional TLV: If present in DMM PDU, is copied from the DMM PDU. The order of the Optional TLVs is preserved.
- End TLV: A 1-octet field the value of which is copied from the DMM PDU.

## 9.17 EXM PDU

EXM is used as experimental OAM request PDU.

### 9.17.1 EXM PDU format

The information elements carried in EXM include:

- OUI: OUI is a 3-octet field that contains the organizationally unique identifier of the organization using the EXM.
- SubOpCode: SubOpCode is a 1-octet field that is used to interpret the remaining fields in the EXM frame.
- EXM Data: Depending on the functionality indicated by the OUI and organizationally specific SubOpCode, EXM may carry one or more TLVs. EXM Data is outside the scope of this Recommendation.

### 9.17.2 EXM PDU format

The EXM PDU format used by a MEP to transmit EXM information is shown in Figure 9.17-1.

1								2								3								4							
8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1		MEL		Version (0)				OpCode (EXM=49)				Flags				TLV Offset															
5		OUI														SubOpCode															
9		<i>[optional EXM data; else End TLV]</i>																													
:																															
:																															
:																End TLV (0)															

**Figure 9.17-1 – EXM PDU format**

The fields of the EXM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is EXM (49).
- Flags: Outside the scope of this Recommendation.
- TLV offset: 1-byte field. Its specific value for EXM is outside the scope of this Recommendation, but must conform to clause 9.1.
- OUI: 3-octet field the values of which are outside the scope of this Recommendation.
- SubOpCode: 1-octet field the values of which are outside the scope of this Recommendation.
- EXM Data: Format and length of this field are outside the scope of this Recommendation.
- End TLV: All-ZEROes octet value.

## 9.18 EXR PDU

EXR is used as experimental OAM reply PDU.

### 9.18.1 EXR information elements

The information elements carried in EXR include:

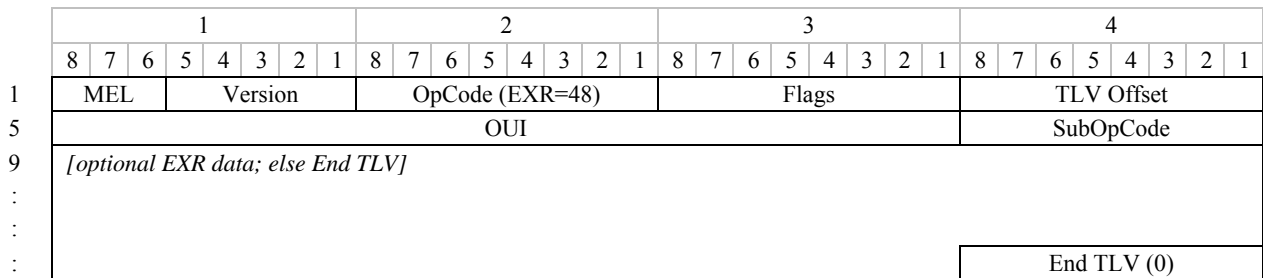
- OUI: OUI is a 3-octet field that contains the organizationally unique identifier of the organization using the EXR.



- SubOpCode: SubOpCode is a 1-octet field that is used to interpret the remaining fields in the EXR frame.
- EXR Data: Depending on the functionality indicated by the OUI and organizationally specific SubOpCode, EXR may carry one or more TLVs. EXR Data is outside the scope of this Recommendation.

### 9.18.2 EXR PDU format

EXR PDU format used by a MEP to transmit EXR information is shown in Figure 9.18-1.



**Figure 9.18-1 – EXR PDU format**

The fields of the EXR PDU format are as follows:

- MEG Level: A 3-bit field the value of which is copied from the last received EXM PDU.
- Version: A 5-bit field the value of which is copied from the last received EXM PDU.
- OpCode: Value for this PDU type is EXR (48).
- Flags: Outside the scope of this Recommendation.
- TLV offset: 1-byte field. Its specific value for EXR is outside the scope of this Recommendation, but must conform to clause 9.1.
- OUI: 3-octet field the value of which is copied from the last received EXM PDU.
- SubOpCode: 1-octet field the values of which are outside the scope of this Recommendation.
- EXR Data: Format and length of this field are outside the scope of this Recommendation.
- End TLV: All-ZEROes octet value.

### 9.19 VSM PDU

VSM is used as vendor-specific OAM request PDU.

#### 9.19.1 VSM PDU format

The information elements carried in VSM include:

- OUI: OUI is a 3-octet field that contains the organizationally unique identifier of the organization using the VSM.
- SubOpCode: SubOpCode is a 1-octet field that is used to interpret the remaining fields in the VSM frame.
- VSM Data: Depending on the functionality indicated by the OUI and organizationally specific SubOpCode, VSM may carry one or more TLVs. VSM Data is outside the scope of this Recommendation.

#### 9.19.2 VSM PDU format

VSM PDU format used by a MEP to transmit VSM information is shown in Figure 9.19-1.

1								2								3								4							
8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1	8	7	6	5	4	3	2	1
1	MEL		Version (0)					OpCode (VSM=51)					Flags					TLV Offset													
5	OUI															SubOpCode															
9	[optional VSM data; else End TLV]																														
:																															
:																															
:																												End TLV (0)			

**Figure 9.19-1 – VSM PDU format**

The fields of the VSM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is VSM (51).
- Flags: Outside the scope of this Recommendation.
- TLV offset: 1-byte field. Its specific value for VSM is outside the scope of this Recommendation, but must conform to clause 9.1.
- OUI: 3-octet field the values of which are outside the scope of this Recommendation.
- SubOpCode: 1-octet field the values of which are outside the scope of this Recommendation.
- VSM Data: Format and length of this field are outside the scope of this Recommendation.
- End TLV: All-ZEROes octet value.

## 9.20 VSR PDU

VSR is used as vendor-specific OAM reply PDU.

### 9.20.1 VSR information elements

The information elements carried in VSR include:

- OUI: OUI is a 3-octet field that contains the organizationally unique identifier of the organization using the VSR.
- SubOpCode: SubOpCode is a 1-octet field that is used to interpret the remaining fields in the VSR frame.
- VSR Data: Depending on the functionality indicated by the OUI and organizationally specific SubOpCode, VSR may carry one or more TLVs. VSR Data is outside the scope of this Recommendation.

### 9.20.2 VSR PDU format

VSR PDU format used by a MEP to transmit VSR information is shown in Figure 9.20-1.



**Table 9-5 – CSF type values**

Flags[6:4]	Type	Comments
000	LOS	Client loss of signal
001	FDI/AIS	Client forward defect indication
010	RDI	Client reverse defect indication
011	DCI	Client defect clear indication

- Period: Bits 3 to 1 indicate transmission period with the encoding Table 9-6.

**Table 9-6 – CSF period values**

Flags[3:1]	Period value	Comments
000	Invalid value	Invalid value for CSF PDUs
001	For further study	For further study
010	For further study	For further study
011	For further study	For further study
100	1s	1 frame per second
101	For further study	For further study
110	1 min	1 frame per minute
111	For further study	For further study

- TLV offset: Set to 0.
- End TLV: All-ZEROes octet value.

## 9.22 SLM PDU

SLM is used to support single-ended ETH-SLM requests, as described in clause 8.4.1.

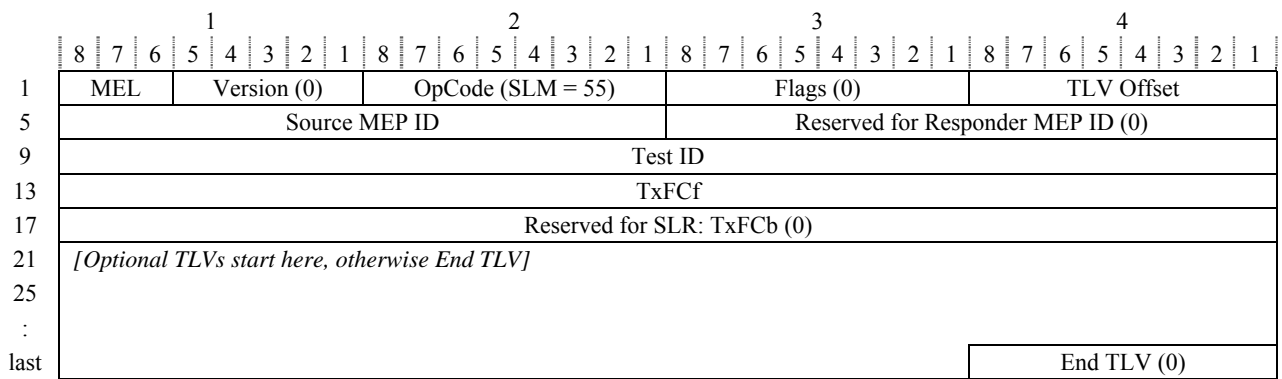
### 9.22.1 SLM information elements

Information elements carried in SLM include:

- Source MEP ID: Source MEP ID is a 2-octet field where the last 13 least significant bits are used to identify the MEP transmitting the SLM frame. MEP ID is unique within the MEG.
- Test ID: Test ID is a 4-octet field set by the transmitting MEP and used to identify a test when multiple tests run simultaneously between MEPs including on concurrent on-demand and proactive tests.
- TxFCf: TxFCf is a 4-octet field which carries the number of SLM frames transmitted by the MEP towards its peer MEP.

### 9.22.2 SLM PDU format

The SLM PDU format used by a MEP to transmit SLM information is shown in Figure 9.22-1.



**Figure 9.22-1 – SLM PDU format**

The fields of the SLM PDU format are as follows:

- MEG Level: Refer to clause 9.1.
- Version: Refer to clause 9.1, value is always 0.
- OpCode: Value for this PDU type is SLM (55).
- Flags: Set to all-ZEROes.
- TLV offset: Set to 16.
- Reserved: Reserved fields are set to all ZEROes.
- Source MEP ID: A 2-octet field used to identify the MEP transmitting the SLM frame, as specified in clause 9.22.1.
- Test ID: A 4-octet field used to identify an unique test among MEPs, as specified in clause 9.22.1.
- TxFCf: A 4-octet integer value representing the number of SLM frames transmitted, as specified in clause 9.22.1.
- Optional TLVs: A Data TLV (Figure 9.3-3) may be included in any SLM transmitted. For the purpose of ETH-SLM, the value part of Data TLV is unspecified.
- End TLV: An all-ZEROes octet value.

### 9.23 SLR PDU

SLR is used to support single-ended ETH-SLM reply, as described in clause 8.4.1.

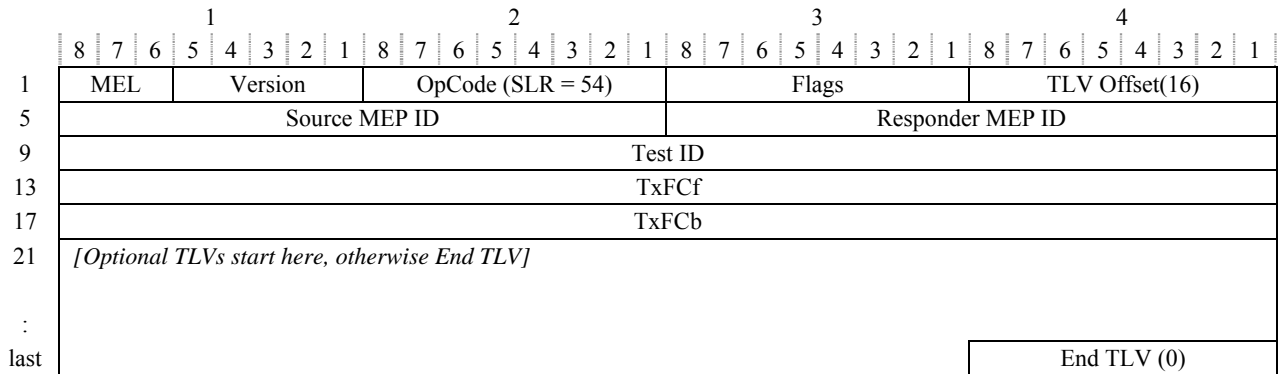
#### 9.23.1 SLR information elements

The information elements carried in SLR include:

- Source MEP ID: Source MEP ID is a 2-octet field that contains the copy of the Source MEP ID field in the received SLM.
- Responder MEP ID: Responder MEP ID is a 2-octet field where the last 13 least significant bits are used to identify the MEP transmitting the SLR frame. MEP ID is unique within the MEG.
- Test ID: Test ID is a 4-octet field that contains the copy of the Test ID field in the received SLM.
- TxFCf: TxFCf is a 4-octet field that contains the copy of the TxFCf field in the received SLM.
- TxFCb: TxFCb is a 4-octet field which carries the number of SLR frames transmitted by the MEP towards its peer MEP.

### 9.23.2 SLR PDU format

SLR PDU format used by a MEP to transmit SLR information is shown in Figure 9.23-1.



**Figure 9.23-1 – SLR PDU format**

The fields of the SLR PDU format are as follows:

- MEG Level: A 3-bit field the value of which is copied from the last received SLM PDU.
- Version: A 5-bit field the value of which is copied from the last received SLM PDU.
- OpCode: Value for this PDU type is SLR (54).
- Flags: A 1-octet field the value of which is copied from the SLM PDU.
- TLV offset: A 1-octet field the value of which is copied from the SLM PDU.
- Reserved: Reserved fields are set to all ZEROes.
- Source MEP ID: A 2-octet field the value of which is copied from the SLM PDU.
- Responder MEP ID: A 2-octet field used to identify the MEP transmitting the SLR frame, as specified in clause 9.22.1.
- Test ID: A 4-octet field the value of which is copied from the SLM PDU.
- TxFCf: A 4-octet field the value of which is copied from the SLM PDU.
- TxFCb: A 4-octet integer value representing the number of SLR frames transmitted, as specified in clause 9.22.1.
- Optional TLVs: If present in SLM PDU, are copied from the SLM PDU.
- End TLV: A 1-octet field the value of which is copied from the SLM PDU.

## 10 OAM frame addresses

OAM frames are identified by a unique EtherType the value of which is 0x8902. OAM frames processing and filtering at a MEP is based on the OAM EtherType and MEG Level fields for both unicast and multicast DA.

As indicated in clauses 7 and 8, the DA in an OAM frame could be multicast or unicast depending on the specific OAM functionality. The SA in an OAM frame is always unicast.

This clause provides further discussion on the choice of DA in specific OAM functions. Table 10-1 provides a summary of the DAs that are applicable for different OAM types.

### 10.1 Multicast destination addresses

Two types of multicast addresses are required depending on the type of OAM function:

- Multicast DA Class 1: OAM frames that are addressed to all MEPs in a MEG (e.g., CCM, Multicast LBM, AIS, etc.).
- Multicast DA Class 2: OAM frames that are addressed to all MIPs and MEPs associated with a MEG (e.g., LTM).

Normally, a single multicast DA Class 1 address and a single multicast DA Class 2 address would be sufficient. However, for a short-term deployment of Ethernet OAM across the current Ethernet equipment, a multicast DA could also implicitly carry the MEG level. This would require 8 distinct addresses for each of the Multicast DA Classes 1 and 2 for the 8 different MEG levels.

Specific values for 8 Multicast addresses for Class 1 and 8 Multicast addresses for Class 2 are 01-80-C2-00-00-3x and 01-80-C2-00-00-3y respectively; x represents the MEG level, with x being a value in the range 0-7, and y represents the MEG level, with y being a value in the range 8-F.

## **10.2 CCM**

CCM frames can be generated with a specific multicast Class 1 DA or unicast DA.

When a multicast DA is used, CCM frames allow discovery of MAC addresses associated with MEPs. The use of a multicast DA also allows for detection of misconnections among flow domain fragments. Detection of misconnections is described in clause 7.1.

When detection of the above conditions is important, a multicast DA must be used for CCM frames. When the above conditions are not expected or are not required to be detected and the data frames in different services instances are distinguished using the unicast DAs, CCM frames can be provisioned to use unicast DAs.

## **10.3 LBM**

LBM frames can be generated with unicast or multicast Class 1 DAs, as per the unicast ETH-LB or multicast ETH-LB functions respectively.

## **10.4 LBR**

LBR frames are always generated with unicast DAs.

## **10.5 LTM**

LTM frames are generated with a multicast Class 2 DA.

A multicast DA is used instead of unicast DA for LTM frames since in current bridges, the MIPs would not be able to intercept a frame with a unicast DA which was not their own address. Therefore the MIPs would not be able to reply and would simply forward the LTM frame with the unicast DA. The limitation is that current ports do not look at the EtherType before looking at the DA.

## **10.6 LTR**

LTR frames are always generated with unicast DAs.

## **10.7 AIS**

AIS frames can be generated with a multicast Class 1 DA, especially in a multipoint MEG.

A unicast DA may be used in provisioned environments for point-to-point connections. However, it requires that the unicast DA of the downstream MEP must have been configured on the MEP transmitting AIS.

## **10.8 LCK**

LCK frames can be generated with a multicast Class 1 DA, especially in a multipoint MEG.

A unicast DA may be used in provisioned environments for point-to-point connections. However, it requires that the unicast DA of the downstream MEP must have been configured on the MEP transmitting LCK.

## **10.9 TST**

TST frames are generated with Unicast DAs. TST frames may be generated with a multicast Class 1 DA if multipoint diagnostics are desired.

## **10.10 APS**

APS frames can be generated with a specific multicast Class 1 DA or Unicast DA.

## **10.11 MCC**

MCC frames are generated with Unicast DAs. For the case when a point-to-point VLAN is being used, a multicast Class 1 DA may be used.

## **10.12 LMM**

LMM frames are generated with unicast DAs. LMM frames may be generated with a multicast Class 1 DA if multipoint measurements are desired.

## **10.13 LMR**

LMR frames are always generated with unicast DAs.

## **10.14 1DM**

1DM frames are generated with unicast DAs. 1DM frames may be generated with a multicast Class 1 DA if multipoint measurements are desired.

## **10.15 DMM**

DMM frames are generated with unicast DAs. DMM frames may be generated with a multicast Class 1 DA if multipoint measurements are desired.

## **10.16 DMR**

DMR frames are always generated with unicast DAs.

## **10.17 EXM**

EXM frames DA is outside the scope of this Recommendation.

## **10.18 EXR**

EXR frames DA is outside the scope of this Recommendation.

## **10.19 VSM**

VSM frames DA is outside the scope of this Recommendation.

## **10.20 VSR**

VSR frames DA is outside the scope of this Recommendation.



### 10.21 CSF

CSF frames can be generated with a multicast Class 1 DA, especially in a multipoint MEG.

A unicast DA may be used in provisioned environments for point-to-point connections. However, it requires that the unicast DA of the downstream MEP must have been configured on the MEP transmitting CSF.

### 10.22 SLM

SLM frames are generated with unicast DAs. SLM frames may be generated with a multicast Class 1 DA if multipoint measurements are desired.

### 10.23 SLR

SLR frames are always generated with unicast DAs.

**Table 10-1 – OAM frame DA**

OAM type	DAs for frames with OAM PDU
CCM	Multicast Class 1 DA or unicast DA
LBM	Unicast DA or multicast Class 1 DA
LBR	Unicast DA
LTM	Multicast Class 2 DA
LTR	Unicast DA
AIS	Multicast Class 1 DA or unicast DA
LCK	Multicast Class 1 DA or unicast DA
TST	Unicast DA or multicast Class 1 DA
Linear APS	Multicast Class 1 DA or unicast DA
Ring APS	Multicast Class 1 DA or unicast DA
MCC	Unicast DA or multicast Class 1 DA
LMM	Unicast DA or multicast Class 1 DA
LMR	Unicast DA
IDM	Unicast DA or multicast Class 1 DA
DMM	Unicast DA or multicast Class 1 DA
DMR	Unicast DA
EXM, EXR, VSM, VSR	Outside the scope of this Recommendation
CSF	Multicast Class 1 DA or unicast DA
SLM	Unicast DA or multicast Class 1 DA
SLR	Unicast DA

## 11 OAM PDU validation and versioning

This clause describes rules for validation and versioning of OAM PDUs, which are designed to ensure that implementations of this Recommendation will interoperate with implementations of future versions of this Recommendation. In addition, these rules allow implementations to provide proprietary, non-standard, extensions to the protocol in a way which does not jeopardize interoperability with future versions of this Recommendation or restrict the ability of future versions of this Recommendation to extend the Recommendation functionality.

NOTE 1 – The change to the LTM format between the 2006 and 2008 versions of this Recommendation did not change the version number; however future revisions to this Recommendation must align with these rules.

NOTE 2 – The rules described here only apply to how PDUs with different versions are interpreted. Further details regarding how the PDUs are subsequently processed, where applicable, may be found in the atomic function definitions in [ITU-T G.8021] and [ITU-T G.8032].

NOTE 3 – These rules do not apply to parts of PDUs that are not specified in ITU-T Recommendations; for example, the data fields of VSM, VSR, EXM and EXR PDUs.

### **11.1 OAM PDU transmission**

OAM PDU transmission is required to meet the following requirements:

- The fixed header fields shall be transmitted exactly as specified in this Recommendation.
- All bits defined as "reserved" in this Recommendation shall be transmitted as 0.
- Additional fields shall not be added to the fixed header specified in this Recommendation.
- Code points reserved in this Recommendation or IEEE 802.1 shall not be transmitted in any OAM PDU; for example, reserved values of the OpCode field (Table 9-1), the TLV Type field (Table 9-2), or the MEG ID format field (Table A.1).
- Additional fields shall not be added to any TLV specified in this Recommendation.

### **11.2 OAM PDU validation in reception**

Received OAM PDUs are subject to a number of validation tests, and are discarded without further processing if they fail these tests. This clause does not provide an exhaustive list of such tests, it covers only those aspects that are most important to future compatibility. In addition to the tests specified here, it may be assumed that if an OAM PDU with a particular OpCode does not conform to the corresponding description in clause 9, it fails the tests. The initial validation test is to ensure that the OAM PDU is sufficiently long to contain the MEG level and version fields. OAM PDUs that fail this test are discarded.

The OAM PDU is subsequently processed in accordance with the numerically lower of 1) the Version field in the OAM PDU; and 2) the highest version number known to the receiving implementation. That is, a version 1 implementation receiving a version 0 OAM PDU processes it according to version 0, and it processes a version 1 OAM PDU according to version 1. It is noted that the imposition on future versions of this Recommendation that all earlier version implementations can process received OAM PDUs correctly, that is, that OAM PDUs specified by later versions of this Recommendation must remain valid when processed according to version 0.

The following validation tests are used, according to the version selected as described above:

- The fixed header length, as determined by the TLV Offset field, is not shorter than the length specified by the selected version.
- The OAM PDU is sufficiently long to contain a fixed header of the length specified by the selected version.

If the OAM PDU contains a TLV that needs to be processed, the following validation tests are used, according to the version selected as described above:

- The OAM PDU is sufficient long to contain a TLV Value field whose length is specified by the TLV Length field.
- A TLV Length field does not indicate a length that is shorter than the minimum length for that TLV as specified by the selected version.

The following criteria shall not be used to validate a received OAM PDU:

- The fixed header can be longer than the length specified by the selected version.

- Bits can be set in the reserved bits of the Flags field.
- A TLV can have a Type field not specified by the selected version of the standard.
- A TLV's Length field can be larger than the value (if any) specified in the selected version of the standard.
- Either the TLV Offset field, or the Length field of the last TLV in the OAM PDU, can indicate a position for the first (next) TLV that coincides with the end of the OAM PDU. That is, the end TLV can be missing from the OAM PDU.
- TLVs may occur in any order in the OAM PDU, unless the descriptions in clause 9 specify otherwise.

NOTE – The selection of the version to be used for processing a received OAM PDU does not impact the version copying requirement if an OAM PDU reply needs to be generated. This means that a version 0 implementation receiving a version 1 OAM PDU request interprets it according to version 0, but replies depending on the replying rules, unless this rule has version dependency. In this case, the reception of a version 1 OAM PDU reply cannot be used as an indication that the OAM PDU request has been processed according to version 1.

### **11.3 OAM PDU reception after validation**

Received OAM PDUs that pass the validation tests described above must be processed in accordance with the following rules, and in accordance with the same version selected for the validation tests (that is, the numerically lower of the Version field in the OAM PDU and the highest version number known to the receiving implementation).

- Only those fields in the fixed header portion of the OAM PDU that are defined in the selected version are processed; any extra octets in the fixed header, if it is longer than the length specified by the selected version, are ignored.
- Any TLV with a Type field not specified by the selected version is ignored, except that if the OAM PDU is forwarded or retransmitted (with or without modification), or if a new OAM PDU is sent in response to the received OAM PDU, the TLV is copied without modification into the forwarded or retransmitted PDU or into the response PDU.
- Any part of the OAM PDU following the End TLV is ignored (the lack of an End TLV is not an error).
- If any TLV's Length field is larger than the value (if any) specified by the selected version, then any octets following those specified by the selected version are ignored.
- All bits undefined in this Recommendation, e.g., reserved bits in the Flags field, are ignored.

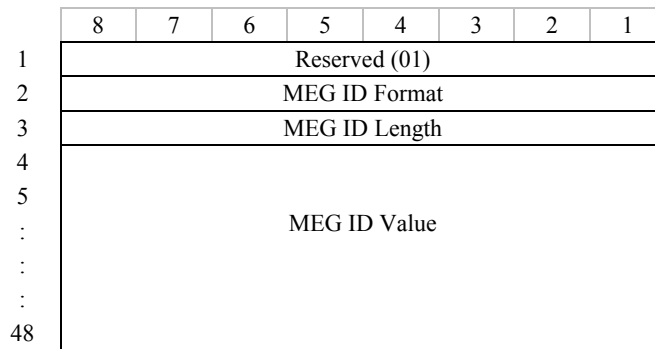
## Annex A

### MEG ID format

(This annex forms an integral part of this Recommendation.)

The features of maintenance entity group identifiers (MEG IDs) are:

- Each MEG ID must be globally unique.
- Where it may be expected that the MEG may be required for path set-up across an inter-operator boundary, the MEG ID must be available to other network operators.
- The MEG ID should not change while the MEG remains in existence.
- The MEG ID should be able to identify the network operator which is responsible for the MEG.
- The generic format of MEG IDs specific to this Recommendation is shown in Figure A.1.



**Figure A.1 – Generic MEG ID format**

The MEG ID format type is identified by the MEG ID Format field. Specific values of MEG ID format type are defined in Table A.1.

**Table A.1 – MEG ID format type**

MEG ID format type value	TLV name
00, 5-31, 64-255	Reserved (Note 1)
1-4	See below (Note 2)
Types specific to this Recommendation	
32	ICC-based format
33-63	Reserved (Note 3)
Note 1 – Reserved for definition by IEEE 802.1.	
Note 2 – Use values as defined in Table 21-20 of [IEEE 802.1ag].	
Note 3 – Reserved for future standardization by ITU-T.	

Figure A.2 shows the format that uses the ITU carrier code (ICC). ICC is a code assigned to a network operator/service provider, maintained by the ITU Telecommunication Standardization Bureau (TSB) as per [ITU-T M.1400].

	8	7	6	5	4	3	2	1
1	Reserved (01)							
2	MEG ID Format (32)							
3	MEG ID Length (13)							
4	0	MEG ID Value[1]						
5	0	MEG ID Value[2]						
	:	:	:	:	:	:	:	:
	:	:	:	:	:	:	:	:
15	0	MEG ID Value[12]						
16	0	MEG ID Value[13]						
19	Unused (= all-ZEROes)							
20								
47								
48								

**Figure A.2 – ICC-based MEG ID format**

The MEG ID value identified in Figure A.2 consists of 13 characters coded according to [ITU-T T.50] (International Reference Alphabet – 7-bit coded character set for information exchange).

It consists of two subfields: the ITU carrier code (ICC) followed by a unique MEG ID code (UMC). The ITU carrier code consists of 1-6 left-justified characters, alphabetic, or leading alphabetic with trailing numeric. The UMC code immediately follows the ICC and shall consist of 7-12 characters, with trailing NULLs, completing the 13-character MEG ID value. The UMC shall be a matter for the organization to which the ICC has been assigned, provided that uniqueness is guaranteed.

## Annex B

### Ethernet link trace (ETH-LT) of [ITU-T Y.1731] interoperability considerations

(This annex forms an integral part of this Recommendation.)

This annex describes the interworking of Ethernet MEPs and MIPs, supporting different types of Ethernet link trace (ETH-LT) (i.e., ETH-LT as defined in [ITU-T Y.1731] and that specified in this Recommendation) and identifies the basic requirements to support interworking under the ME where two types of MEPs or MIPs exist.

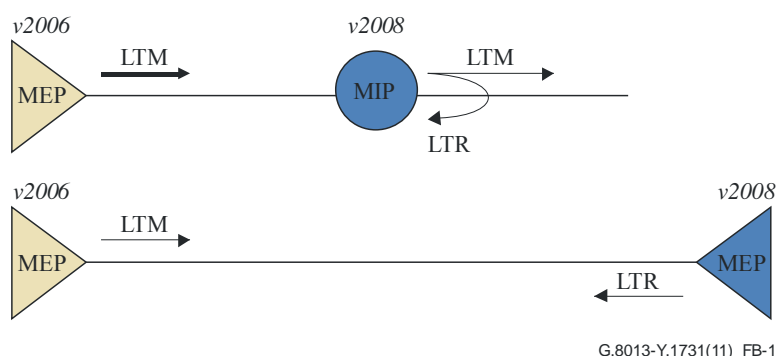
#### B.1 Ethernet link trace (ETH-LT) as defined in [ITU-T Y.1731]

The ETH-LT defined in [ITU-T Y.1731] differs from the one defined in this Recommendation in the following:

- LTM transmission and its PDU, as given in clauses 7.3.1 and 9.5 of [ITU-T Y.1731] do not define the LTM egress identifier TLV and its format, whereas they are defined as mandatory in this Recommendation.
- LTR transmission and its PDU, as given in clauses 7.3.2 and 9.6 of [ITU-T Y.1731] do not define the LTR egress identifier TLV and its format, whereas they are defined as mandatory in this Recommendation. Also, Reply Ingress TLV and Reply Egress TLV were optional in [ITU-T Y.1731], whereas they are defined as mandatory in this Recommendation.
- FwdYes and TerminalMEP were defined in bit 7 and bit 6 of Flags of LTR in this Recommendation, i.e., they were not defined in [ITU-T Y.1731].
- At a MIP, ETH-LT responder was not defined, and both ingress and egress ports could be set as MIP in a v2006 equipment, whereas in this Recommendation ETH-LT responder is defined so that there can only be one MIP per equipment.

#### B.2 Interworking with [ITU-T Y.1731]

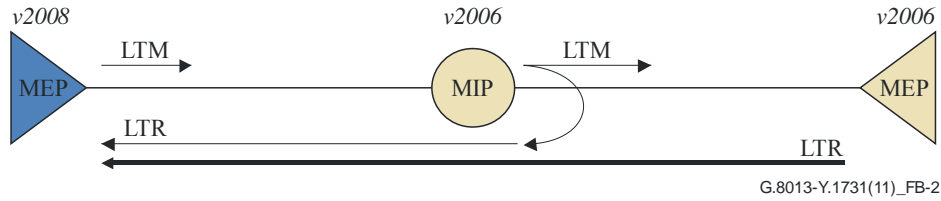
In the case of an ME consisting of a v2006 MEP that transmits ETH-LTM and some v2008 MIPs, or of an ME consisting of a v2006 MEP that transmits ETH-LTM and a v2008 MEP that receives ETH-LTM and transmits ETH-LTR, the v2008 MIP or v2008 MEP may discard ETH-LTM from the v2006 MEP due to the absence of LTM egress identifier TLV. In this case, for keeping the interoperability, the v2008 MIP may forward ETH-LTM and transmit ETH-LTR by recognizing that the ETH-LTM does not have the TLV and behaving as v2006 MIP. Similarly, the v2008 MEP may transmit ETH-LTR by recognizing that the ETH-LTM does not have the TLV and behaving as v2006 MEP. See Figure B.1.



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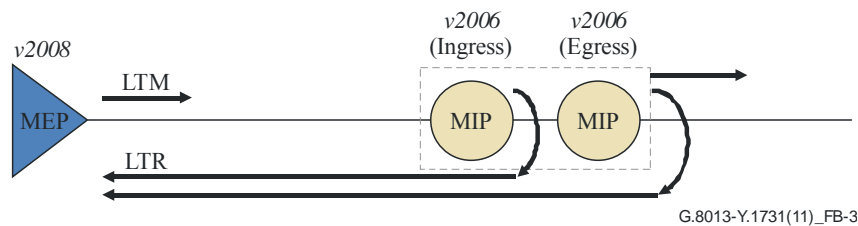
Figure B.1 – Interoperability case 1

In the case of an ME consisting of a v2008 MEP that transmits ETH-LTM and some v2006 MIPs and/or of a v2008 MEP that receives ETH-LTM and transmits ETH-LTR, the v2008 MEP receives ETH-LTR without LTR egress identifier TLV and without reply ingress TLV or reply egress TLV generated by v2006 MIPs and/or MEP. The absence of these TLVs in ETH-LTR is considered invalid in the v2008 version. In order to keep interoperability, the v2008 version may be configured to identify this ETH-LTR as valid (see Figure B.2)



**Figure B.2 – Interoperability case 2**

In the case of an ME consisting of a v2008 MEP that transmits ETH-LTM and some v2006 MIPs located in both the ingress and egress ports of an equipment, the equipment may transmit two ETH-LTRs to the v2008 MEP. In receiving the ETH-LTRs at the v2008 MEP, the behaviour is same as in the case mentioned above (see Figure B.3). It is noted that this behaviour is compatible with the LTR analysis according to Annex J.5 of [IEEE 802.1 ag], as long as each of the MIPs that decrement the LTM's TTL field also return an LTR.



**Figure B.3 – Interoperability case 2**

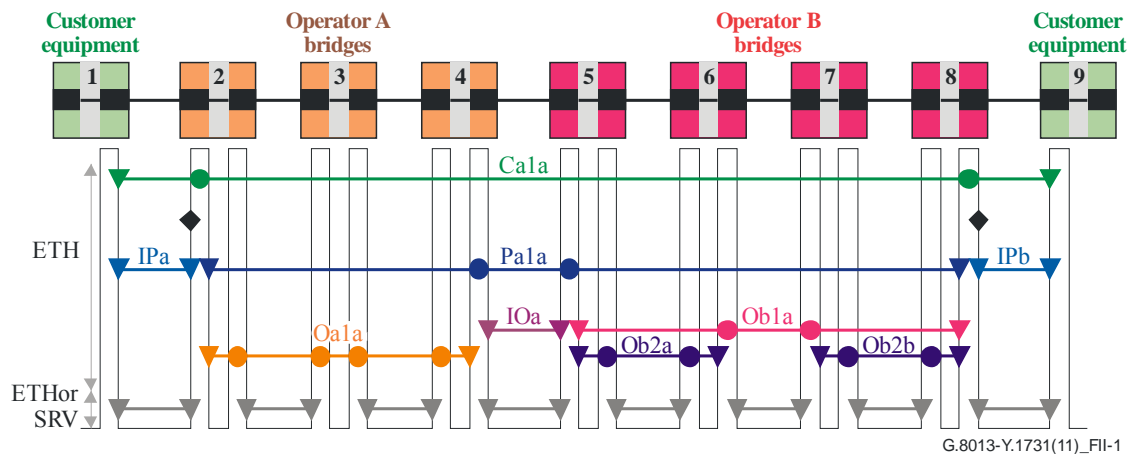
# Appendix I

## Ethernet Network Scenarios

(This appendix does not form an integral part of this Recommendation.)

### I.1 Shared MEG Levels Example

Figure I.1 provides an example scenario with the default assignment of MEG levels, where the customer, provider, and operator roles share the MEG levels. In the figure, triangles represent MEPs, circles represent MIPs, and diamonds represent TrCPs.



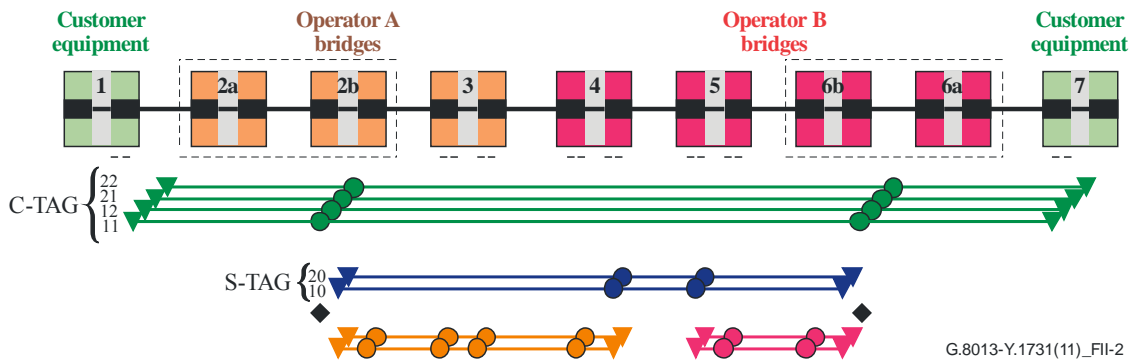
**Figure I.1 – Example MEG level assignment for shared MEG levels**

- Customer ME (Ca1a) can be assigned a customer MEG level 5. This allows for more customer MEs to be created at higher MEG levels, i.e., 6 and 7, if these customer MEs at additional customer MEG levels are needed.
- Provider ME (Pa1a) can be assigned a provider MEG level 4. This allows for more provider MEs to be created at a lower MEG level, i.e., 3, if additional MEs at a lower provider MEG level are needed.
- End-to-end operator MEs (Oa1a and Ob1a) can be assigned an operator MEG level 2. This allows for more operator MEs to be created at lower MEG levels, i.e., 1 and 0, if these operator MEs at additional operator MEG levels are needed in each operator network.
- Segment operator MEs in Operator B's network (Ob2a and Ob2b) can be now assigned a lower MEG level, for example 1, if Operator B needs such MEs.
- MEs between the customer and provider (IPa and IPb) can be assigned a MEG level 0. This allows provider to filter such OAM frames at UNI\_N since the provider is required to provide transparency only to customer MEG levels 7, 6, and 5.
- Inter-operator ME (IOa) can be assigned a MEG Level 0. This allows the operator to filter such OAM frames since the operator is required to provide transparency only to customer and provider MEG levels.

### I.2 Independent MEG levels example

Figure I.2 provides an example scenario where the customer and service provider do not share the MEG levels. However, the service provider and operator share the MEG levels. In the figure, triangles represent MEPs, circles represent MIPs, and diamonds represent TrCPs.





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**Figure I.2 – Example MEG level assignment for independent MEG levels**

- In the above example, four customer VLANs (11, 12, 21 and 22) and the corresponding customer MEGs (C-TAG, in the figure) are completely independent of the two service provider VLANs (20 and 10) and corresponding service provider MEGs (S-TAG in the figure).
- As a consequence, the customer and service provider can independently use the all eight MEG levels.
- The service provider and operator however share the MEG level space, in a manner similar to the one of Figure I.1. In this case, the eight MEG levels can be agreed mutually between the service provider and the operator.
- In the above example, the customer must send OAM frames as VLAN-tagged or priority-tagged frames to utilize all eight MEG levels independently. However, if the customer uses untagged OAM frames, the MEG levels may not be independent anymore and the customer and provider MEG levels need to be mutually agreed between the customer and the service provider.

## Appendix II

### Frame loss measurement

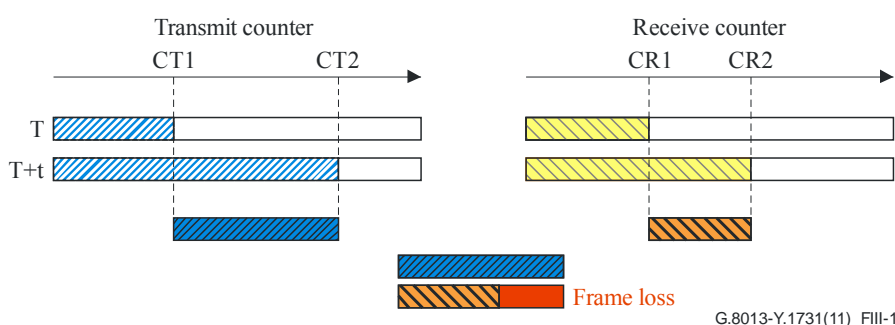
(This appendix does not form an integral part of this Recommendation.)

The four cases below should be taken into account for frame loss calculation.

- a) No wrap around for either the transmit or the receive counter.
- b) Only the transmit counter wraps around.
- c) Only the receive counter wraps around.
- d) Both transmit and receive counters wrap around.

For each case, the frame loss can be calculated as follows:

- a) No wrapping around for both transmit and receive counters.

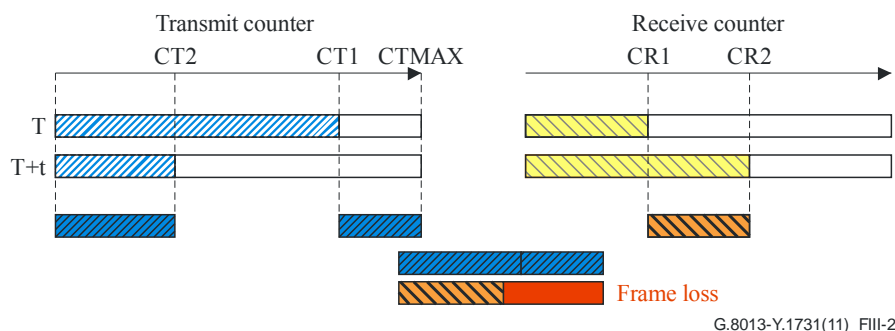


**Figure II.1 – No wrap around**

For this case, the frame loss can be calculated by the simple calculation:

$$\text{Frame Loss} = (\text{CT2} - \text{CT1}) - (\text{CR2} - \text{CR1})$$

- b) Only transmit counter wraps around.

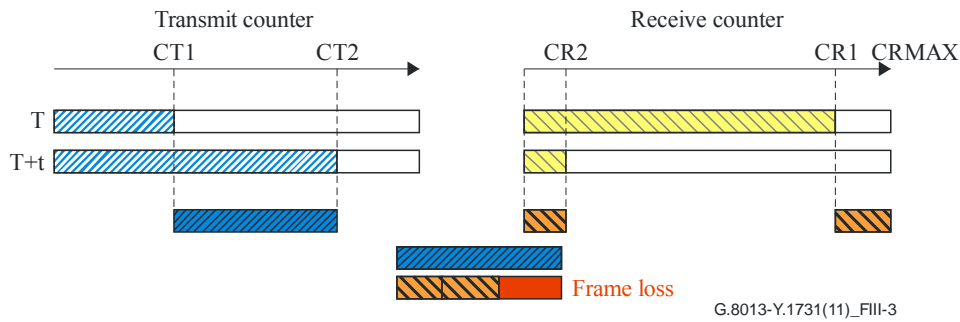


**Figure II.2 – Transmit counter wraps around**

In this case, frame loss can be calculated by the following equation, as described in the previous clause

$$\begin{aligned} \text{Frame Loss} &= ((\text{CTMAX} - \text{CT1}) + \text{CT2} + 1) - (\text{CR2} - \text{CR1}) \\ &= (\text{CT2} - \text{CT1}) - (\text{CR2} - \text{CR1}) + (\text{CTMAX} + 1) \end{aligned}$$

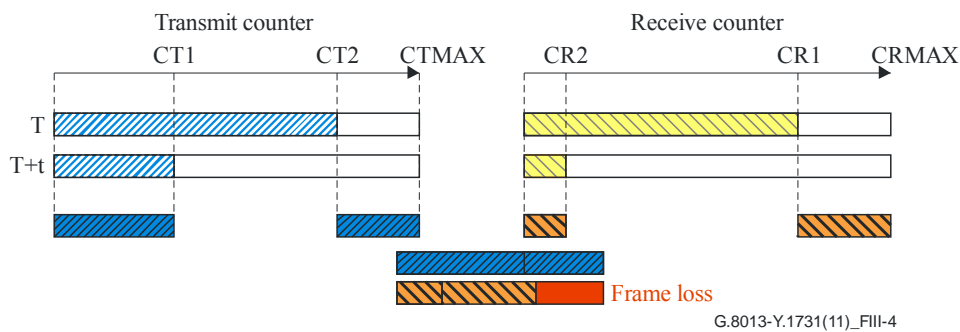
- c) Only the receive counter wraps around.



**Figure II.3 – Receive counter wraps around**

$$\begin{aligned} \text{Frame Loss} &= (CT2 - CT1) - ((CRMAX - CR1) + CR2 + 1) \\ &= (CT2 - CT1) - (CR2 - CR1) - (CRMAX + 1) \end{aligned}$$

d) Both transmit and receive counters wrap around.



**Figure II.4 – Both counters wrap around**

$$\begin{aligned} \text{Frame Loss} &= ((CTMAX - CT1) + CT2 + 1) - ((CRMAX - CR1) + CR2 + 1) \\ &= (CT2 - CT1) - (CR2 - CR1) + (CTMAX + 1) - (CRMAX + 1) \end{aligned}$$

## II.1 Simplified calculation for frame loss

If the calculation is processed in an unsigned value schema, the calculation formula for the frame loss can be greatly simplified by the following characteristics:

$$N + (MAX + 1) \equiv N \pmod{MAX + 1}$$

$$N - (MAX + 1) \equiv N \pmod{MAX + 1}$$

Therefore the formulas for frame loss (described in the clauses 8.1.1 and 8.1.2) can be transformed as follows.

- a) Frame Loss =  $\underline{(CT2 - CT1) - (CR2 - CR1)}$
- b) Frame Loss =  $(CT2 - CT1) - (CR2 - CR1) + CTMAX + 1$   
 =  $((CT2 + (CTMAX + 1)) - CT1) - (CR2 - CR1)$   
 =  $\underline{(CT2 - CT1) - (CR2 - CR1)}$
- c) Frame Loss =  $(CT2 - CT1) - (CR2 - CR1) - (CRMAX + 1)$   
 =  $(CT2 - CT1) - ((CR2 + CRMAX + 1) - CR1)$   
 =  $\underline{(CT2 - CT1) - (CR2 - CR1)}$
- d) Frame Loss =  $(CT2 - CT1) - (CR2 - CR1) + (CTMAX + 1) - (CRMAX + 1)$   
 =  $((CT2 + (CTMAX + 1)) - CT1) - ((CR2 + (CRMAX + 1)) - CR1)$   
 =  $\underline{(CT2 - CT1) - (CR2 - CR1)}$

As described above, the frame loss can be calculated by the single calculation formula for any case if it is calculated in unsigned value schema.

## II.2 Frame counter wrapping periodicity

This clause provides a view of wrapping periodicity of 4-octet frame counters for different interface rates and different frame sizes. The interfaces rates considered are 1 Gbit/s, 10 Gbit/s, and 100 Gbit/s. Frames sizes considered are 64-octet (minimum Ethernet frame size) and 1522-octet (maximum Ethernet frame size)

**Table II.1 – Frame counter wrapping period**

Interface rate	Frame size	4-octet frame counter wrapping period
1 Gbit/s	64-octet	$(2^{32})/((10^9)/((64+12)*8)) = 2611$ seconds
1 Gbit/s	1522-octet	$(2^{32})/((10^9)/((1522+12)*8)) = 52707$ seconds
10 Gbit/s	64-octet	$(2^{32})/(((10*(10^9))/((64+12)*8)) = 261$ seconds
10 Gbit/s	1522-octet	$(2^{32})/(((10*(10^9))/((1522+12)*8)) = 5270$ seconds
100 Gbit/s	64-octet	$(2^{32})/(((100*(10^9))/((64+12)*8)) = 26$ seconds
100 Gbit/s	1522-octet	$(2^{32})/(((100*(10^9))/((1522+12)*8)) = 527$ seconds

## **Appendix III**

### **Network OAM interworking**

(This appendix does not form an integral part of this Recommendation.)

The requirements for interworking between layered networks are the following:

- Upon detection of a defect condition in the server layer, the adaptation function between the server and client layer should be able to insert AIS in the client layer.
- The format of AIS inserted is specific to the client layer.

As an example, when the client layer is Ethernet, a server MEP is used.

## Appendix IV

### Mismerge detection limitation

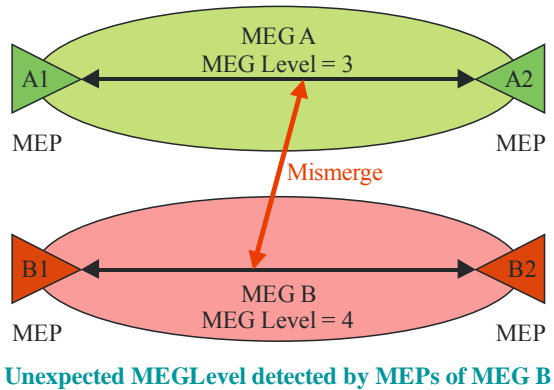
(This appendix does not form an integral part of this Recommendation.)

MEPs consider only CCM frames with their own or lower MEG level for defect detection. CCM frames with higher MEG levels are passed through in order to provide OAM transparency as defined in clause 5.4. This behaviour leads to a limitation in the Mismerge detection as shown in Figure IV.1 below.

In case of a mismerge between MEGs with different MEG levels, the MEPs of the MEG with the lower MEG level will not detect any defect as the CCM frames coming from the MEG with the higher MEG level are passed through transparently by the MEPs. The MEPs of the MEG with the higher MEG level will detect Unexpected MEGLevel.

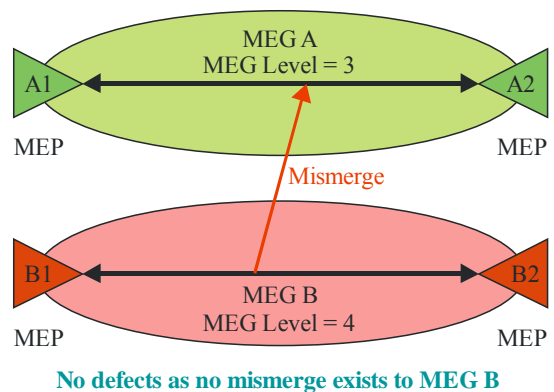
In case of a uni-directional mismerge from the MEG with the higher MEG level to the MEG with the lower MEG level, no defect will be detected.

**No defects detected by MEPS of MEG A as only MEG Levels 3 and lower are considered!**



a) Bi-directional mismerge

**No defects detected by MEPs of MEG A, as only MEG levels 3 and lower are considered!**



b) Uni-directional mismerge

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Figure IV.1 – Mismerge detection limitation

## Appendix V

### Terminology alignment with [IEEE 802.1ag]

(This appendix does not form an integral part of this Recommendation.)

The relationship of the terminology used in this Recommendation and [IEEE 802.1ag] is captured below.

**Table V.1 – Terminology mapping**

<b>ITU-T G.8013/Y.1731 term</b>	<b>IEEE 802.1ag term</b>	<b>Comments</b>
MEG	MA	
MEG ID	MAID (Domain Name + Short MA Name)	Unlike in [IEEE 802.1ag], the MEG ID does not imply a split between domain name and a short MEG name in ITU-T Y.1731.
MEG level	MA Level	

## Appendix VI

### Examples showing accuracy for ETH-SLM measurement

(This appendix does not form an integral part of this Recommendation.)

Synthetic loss measurement is a sampling technique for measuring frame loss; therefore, the measured FLR will be distributed around the actual loss value according to a binomial distribution. The mean measured FLR will always be equal to the actual FLR, while the standard deviation depends on the number of samples. The standard deviation can therefore be used to illustrate the accuracy of the measured FLR result. Table VI.1 shows the standard deviation for various real loss values and number of samples (i.e., SLM frames sent). When ETH-SLM is used, the number of samples should be chosen such that the standard deviation is low, when compared to any FLR threshold that is being used to trigger an action. This ensures the chance of false positives is low.

**Table VI.1 – Standard deviation for various real loss values and number of samples**

Actual FLR	Number of samples	Transmission interval	Std. Dev. (FLR % points)
50%	10	100 ms	15.81%
50%	100	10 ms	5.00%
50%	1000	1 ms	1.58%
10%	10	100 ms	9.49%
10%	100	10 ms	3.00%
10%	1000	1 ms	0.95%
1%	10	100 ms	3.15%
1%	100	10 ms	0.99%
1%	1000	1 ms	0.31%
0.1%	10	100 ms	1.00%
0.1%	100	10 ms	0.31%
0.1%	1000	1 ms	0.1%

Note that if the number of samples is increased by a factor of  $n$ , the standard deviation is reduced by a factor of  $\sqrt{n}$ .



## Bibliography

- [b-IETF RFC 2544] IETF RFC 2544 (1999), *Benchmarking Methodology for Network Interconnect Devices*.  
<<http://www.ietf.org/rfc/rfc2544.txt>>



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