# ITU-T 

TELECOMMUNICATION

SERIES Z: LANGUAGES AND GENERAL SOFTWARE ASPECTS FOR TELECOMMUNICATION SYSTEMS
Formal description techniques (FDT) - Application of formal description techniques

## Notations and guidelines for the definition of ITU-T languages

Recommendation ITU-T Z. 111

## ITU-T Z-SERIES RECOMMENDATIONS

## LANGUAGES AND GENERAL SOFTWARE ASPECTS FOR TELECOMMUNICATION SYSTEMS

| FORMAL DESCRIPTION TECHNIQUES (FDT) |
| :--- |
| Specification and Description Language (SDL) |
| Application of formal description techniques |
| Message Sequence Chart (MSC) |
| User Requirements Notation (URN) |
| Testing and Test Control Notation (TTCN) |
| PROGRAMMING LANGUAGES |
| CHILL: The ITU-T high level language |
| MAN-MACHINE LANGUAGE |
| General principles |
| Basic syntax and dialogue procedures |
| Extended MML for visual display terminals |
| Specification of the man-machine interface |
| Data-oriented human-machine interfaces |
| Human-machine interfaces for the management of telecommunications networks |
| QUALITY |
| Quality of telecommunication software |
| Quality aspects of protocol-related Recommendations |
| METHODS |
| Methods for validation and testing |
| MIDDLEWARE |
| Processing environment architectures |
| Z |

[^0]
## Recommendation ITU-T Z. 111

## Notations and guidelines for the definition of ITU-T languages

## Summary

Recommendation ITU-T Z. 111 provides meta-grammars for ITU-T Recommendations that define ITU-T languages in the ITU-T X. 680 series and the ITU-T Z series of Recommendations on languages for specification, implementation, modelling and testing. This allows the description of these meta-grammars that define the abstract or concrete grammar (syntax, constraints and semantics) of languages without having to repeat the meta-grammar (such as lexical naming rules, or the description of Backus-Naur Form syntax) as a preamble or annex to each language definition.

This Recommendation draws common elements from the meta-grammars of various languages, covering issues such as common lexical rules, the use of a universal character set, and syntax and constraint description for languages at both the abstract and concrete level.

## History

| Edition | Recommendation | Approval | Study Group | Unique ID* |
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## Keywords

ITU-T languages, meta-grammars, abstract grammar, concrete grammar, syntax, constraints, semantics.

[^1]
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The World Telecommunication Standardization Assembly (WTSA), which meets every four years, establishes the topics for study by the ITU-T study groups which, in turn, produce Recommendations on these topics.
The approval of ITU-T Recommendations is covered by the procedure laid down in WTSA Resolution 1.
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## Table of Contents

Page
1 Scope ..... 1
1.1 Objective ..... 1
1.2 Application ..... 1
2 References ..... 1
3 Definitions ..... 2
4 Abbreviations and acronyms ..... 2
5 Conventions ..... 2
5.1 Grammars ..... 2
5.2 Basic definitions ..... 2
5.3 Presentation style ..... 2
5.4 Metalanguages ..... 5
Annex A - Use of graphical meta-symbols ..... 20
Appendix I - Mapping from the metamodel presentation to the textual presentation: High-level description ..... 21
Bibliography ..... 23

## Introduction

Scope-objective: Notations and guidelines for the definition of formal languages are provided.
Coverage: The main features of this Recommendation are notations for defining the abstract and concrete syntaxes of languages, and a common structure for defining languages.
Applications: This Recommendation should be applied to new and possibly to revised formal language Recommendations.

Status/Stability: The notation and guidelines given are stable and have been used on some of the existing language Recommendations, such as Recommendations the ITU-T Z. 100 series and ITU-T Z.151. There is scope for further guidelines, extensions to the notations given, algorithms for conversions between the textual and metamodel presentations of abstract grammars, mappings to XML-based concrete syntaxes for interchange formats, and the possibility of further formal notations, such as a formal notation for specifying constraints.
Associated work: This Recommendation is generally associated with the study of formal languages for telecommunications applications in ITU-T, in particular the languages defined by the X. 680 series, Z .100 series, Z .120 series, Z .150 series and $\mathrm{Z} .160 / \mathrm{Z} .170$ series of ITU-T Recommendations.

This Recommendation is also related to the unified modelling language (UML) and meta object facility (MOF) work of the object management group (OMG).

Background: Before the introduction of this Recommendation, the language Recommendations defined by ITU-T used different ways of describing the syntax and semantics of languages. This hinders the understanding, verifiability, and maintainability of these languages, and prevents their simple harmonization. Some language Recommendations have used grammars in Backus-Naur Form (BNF) to define concrete syntaxes and, in some cases, abstract syntaxes. More recent language definitions have used MOF metamodels to capture several aspects of languages. In both approaches, there has been a separation of abstract grammar from the concrete notation with a defined relationship between them. Each approach has its benefits and drawbacks. MOF metamodels are appealing due to their graphical nature (where associations and inherited concepts are explicit), whereas BNF grammars are easily analyzable by tools and there is considerable experience in their use.

This Recommendation gives guidelines for the definition of MOF-based metamodels describing languages in a way that is compatible with the approaches based on BNF grammars.

Traditional BNF-based grammars usually lack some of the capabilities of the MOF approach such as inheritance. MOF and UML-based meta-metamodels contain many features that make metamodels unnecessarily complex, difficult to understand semantically, and difficult to map to BNF grammars. This Recommendation focuses on a subset of modelling features that is expressive enough to describe language metamodels and that is isomorphic between the two approaches.

## Recommendation ITU-T Z. 111

## Notations and guidelines for the definition of ITU-T languages

## 1 Scope

This Recommendation provides notation and guidelines for the definition of formal languages defined in new and optionally in revised language Recommendations that define formal languages or description techniques (such as those listed in clause 3.1.2 of the Recommendation "Criteria for use of formal description techniques by ITU-T" [b-ITU-T Z.110]).

### 1.1 Objective

The objective is to provide a basis for a common structure and meta-grammar notations to be used in formal language Recommendations, so that it is not necessary for each language Recommendation to describe its structure and the meta-grammars used: instead each language Recommendation can contain a reference to this Recommendation. By using a common structure and meta-grammar notations, it is also easier to integrate the different formal languages, therefore making it easier to use the languages together both with and without the support of tools. The use of a common structure and meta-grammar notations makes it easier to build tools that combine the ITU-T languages with each other and other notations. The use of a common structure and meta-grammar notations in different formal language Recommendations also makes it easier to understand several formal language Recommendations, because the structure and meta-grammar only has to be learnt once.

### 1.2 Application

When a new Recommendation is being drafted for a formal language, the application of this Recommendation should be applied. If it is decided not to apply this Recommendation, the reasons should be stated in the formal language Recommendation.
When a Recommendation for an existing formal language is being revised, the application of this Recommendation should be considered taking into account costs and benefits. If this Recommendation is not applied or is applied partially, the revised formal language Recommendation should at least contain a statement explaining that the structure and meta-language notations of the formal language Recommendation predated the approval of this Recommendation.

The structure and meta-grammar notations given in this Recommendation are also likely to be useful for the definition of other formal languages, not just formal languages defined in ITU-T Recommendations. If it is later decided to consider making a Recommendation for such a formal language, the prior application of this Recommendation will be a benefit.

## 2 References

The following ITU-T Recommendations and other references contain provisions which, through reference in this text, constitute provisions of this Recommendation. At the time of publication, the editions indicated were valid. All Recommendations and other references are subject to revision; users of this Recommendation are therefore encouraged to investigate the possibility of applying the most recent edition of the Recommendations and other references listed below. A list of the currently valid ITU-T Recommendations is regularly published. The reference to a document within this Recommendation does not give it, as a stand-alone document, the status of a Recommendation.
$\begin{array}{ll}\text { [ITU-T T.55] } & \begin{array}{l}\text { Recommendation ITU-T T. } 55 \text { (2008), Use of the universal multiple-octet coded } \\ \text { character set (UCS). } \\ \text { shttp://www.itu.int/rec/T-REC-T. } 55>\end{array}\end{array}$
[OMG UML] OMG Unified Modeling Language (OMG UML): Superstructure, version 2.5, formal/15-03-01 of the Object Management Group (OMG). <http://www.omg.org/spec/UML/2.5/PDF >

## 3 Definitions

No specific definition is provided in this Recommendation.

## 4 Abbreviations and acronyms

This Recommendation uses the following abbreviations:
BNF Backus-Naur Form
MOF Meta Object Facility
OCL Object Constraint Language
OMG Object Management Group
SDL Specification and Description Language
UML Unified Modeling Language
XML eXtensible Markup Language

## 5 Conventions

If a language Recommendation conforms to this Recommendation, it conforms to the conventions defined in this clause and therefore these do not need to be repeated in the language Recommendation.

### 5.1 Grammars

A description only conforms to a Recommendation if it conforms to both the Concrete grammar and Abstract grammar of the corresponding language definition: that is, the description must be both recognizable as the language defined in the Recommendation and have the same meaning as defined by the Semantics in the Recommendation. If further concrete grammars are defined (in additional clauses, annexes or Recommendations), each of the concrete grammars has a definition of its own syntax and of its relationship to the abstract grammar (that is, how to transform into the abstract syntax). Using this approach, there is only one definition of the semantics of a language: the semantics of each of the concrete grammars is identified via its relationship to the abstract grammar. This approach also ensures that any further grammars are equivalent.

For some constructs of the concrete grammar, there may be no directly equivalent abstract syntax. In these cases, a Model is given for the transformation from concrete syntax into the concrete syntax of other constructs that (directly or indirectly via further models) have an abstract syntax. Items that have no mapping to the abstract syntax (such as comments) do not have any formal meaning.

### 5.2 Basic definitions

Some general concepts and conventions are used throughout a language Recommendation; their definitions are given in the following subclauses.

### 5.3 Presentation style

The following presentation style is used to separate the different language issues under each topic.

### 5.3.1 Division of text

A language Recommendation is organized by language features described by an optional introduction, which by convention shall be informative rather than normative (see clause 5.3.2), followed by titled enumeration items for:
a) Abstract grammar - Described by abstract syntax (either a textual grammar or a graphical metamodel) and static conditions (that is, static constraints) for a model in the language to be well-formed for an abstract model where artifacts of the concrete syntax are ignored.
b) Concrete grammar - Described by the graphical syntax, static conditions and rules for the graphical syntax to be well-formed (including drawing rules), the relationship of this syntax with the abstract syntax.
c) Semantics - Gives meaning to a construct, provides the properties it has, the way in which it is interpreted and any dynamic conditions that have to be fulfilled for the construct to behave well in the sense of the language defined.
d) Model - Gives the mapping for notations that do not have a direct abstract syntax and modelled in terms of other concrete syntax constructs. A notation that is modelled by other constructs is known as a shorthand, and is considered to be derived syntax for the transformed form.

To avoid misunderstandings, a language Recommendation that follows this Recommendation and uses informative introductions shall include a clause that states:

- Where a section has introductory text followed by a titled enumeration items (such as Abstract grammar, Concrete grammar, Semantics, Model) the introductory text before the first titled enumeration item in a clause is for information only, and is treated as if written as a NOTE (that is, informative not normative), except if it is explicitly stated as normative.


### 5.3.2 Titled enumeration items

Where a topic in the language definition has an introduction followed by a titled enumeration item, the introduction is considered to be an informal part of the Recommendation presented only to aid understanding and not to make that Recommendation complete.
If there is no text for a titled enumeration item (Abstract grammar, Concrete grammar, Semantics, Model), the whole item is omitted.
It is permitted for a Recommendation to define additional titled enumeration items, such as an Example section.
The remainder of this clause describes the other special formalisms used in each titled enumeration item and the titles used. It can also be considered as an example of the typographical layout of firstlevel titled enumeration items defined above where this text is part of an introductory section.

## a) Abstract grammar

The notations for defining the abstract grammar are defined in clause 5.4.1. In the rest of this clause, the two notations are taken as equivalent. A rule in the textual abstract syntax (see clause 5.4.1.1) is equivalent to a model element in a graphical metamodel (see clause 5.4.1.2). The terms 'model element' and 'graphical metamodel' can be substituted for 'rule' and 'abstract syntax' respectively.
If the titled enumeration item Abstract grammar is omitted, then there is no additional abstract syntax for the topic being introduced, and the concrete syntax shall map onto the abstract syntax defined by another numbered text clause.

The rules in the abstract syntax may be referred to from any of the titled enumeration items by use of the rule name in italics.

The rules in the formal notation may be accompanied by paragraphs that define conditions which shall be satisfied by a well-formed piece (text and/or graphics) in the language being defined and which can be checked without dynamic interpretation of that piece. The static conditions at this point refer only to the abstract syntax. The static conditions shall be expressed in natural language, preferably supplemented by expression of the constraints in a more formal constraint language. Where it is possible to capture a constraint in the abstract syntax rather than by a static condition, this is preferable provided it does not make the syntax too complex. Static conditions, which are only relevant for the concrete syntax, are defined with the concrete syntax. Together with the abstract syntax, the static conditions for the abstract syntax define the abstract grammar of the language.

## b) Concrete grammar

The concrete syntax shall be specified in the extended Backus-Naur Form of syntax description defined in clause 5.4.2, except where a graphical metamodel has been used for the abstract grammar, in which case it is permitted to extend the abstract metamodel to include concrete attributes. The use of natural language to define concrete syntax or definition by examples should be avoided.

The concrete syntax is accompanied by paragraphs defining the static conditions which must be satisfied in a well-formed definition and which can be checked without interpretation of a definition. Static conditions (if any) for the abstract grammar also apply. Where it is possible to capture a constraint in the concrete syntax rather than by a static condition, this is preferable provided it does not make the syntax too complex. The concrete grammar should not repeat constraints present in the abstract grammar.

In many cases there is a simple relationship between the concrete and abstract grammar, because the concrete syntax rule is simply represented by a single rule (or metamodel element) in the abstract grammar. When the same name is used in the abstract grammar and concrete syntax in order to signify that they represent the same concept, the text "<x> in the concrete syntax represents $X$ in the abstract grammar" is implied in the language description and therefore does not need to be stated explicitly. In this context, spaces and hyphens are treated as equivalent, case is ignored but underlined semantic sub-categories (see clause 5.4.2) are significant, so that <integer name> represents Integer-name in the abstract syntax.
Concrete syntax that is not a shorthand form is strict concrete syntax. The relationship from concrete syntax to abstract syntax is defined only for the strict concrete syntax. Strict syntax is always defined in the Concrete grammar. The syntax for a shorthand is allowed to be defined in the Model.
The relationship between concrete syntax and abstract syntax is omitted if the topic being defined is a shorthand form that is modelled by other constructs of the defined language (see Model below).

When the name of a non-terminal ends in the concrete grammar with the word "diagram" and there is a name in the abstract grammar that differs only by ending in the word definition, then the two rules represent the same concept. For example, <system diagram> in the concrete grammar and System-definition in the abstract grammar correspond.

When the name of a non-terminal ends in the concrete grammar with the word "area" and there is a name in the abstract grammar that differs only by having the word area deleted, then the two rules represent the same concept. For example, <state partition area> in the concrete grammar and State-partition in the abstract grammar correspond.

## c) Semantics

Properties are relations between different concepts in the defined language. Properties are used in the rules for a model to be well-formed.

An example (from the specification of SDL-2010 in [b-ITU-T Z.101] clause 9 Agents, Semantics) of a property is the complete valid input signal set of an agent. This property is used in the static condition "Signals mentioned in <output body>s of the state machine of an agent type shall be in the
complete valid input signal set of the agent type or in the <signal list> of a gate in the direction from the agent type" (see [b-ITU-T Z.101] clause 11.13.4, Output, Concrete grammar).
Properties are static if they can be determined without interpretation of definitions in the defined language and are dynamic if an interpretation of the same is required to determine the property.

The interpretation is described in an operational manner. Whenever there is a list in the abstract syntax, the list is interpreted in the order given. That is, the Recommendation describes how objects of the semantic domain are created from a definition and how these objects are interpreted within an "abstract machine". Lists are denoted in the abstract syntax by the suffixes "*" and " + " (see clause 5.4.1).
Dynamic conditions are conditions to be satisfied during interpretation that cannot be checked without interpretation. Dynamic conditions may lead to errors.
NOTE - Behaviour of the object of the semantic domain is produced by "interpreting" the definition. The word "interpret" is explicitly chosen (rather than an alternative such as "executed") to include both mental interpretation by a human and the interpretation of the definition by a computer.
d) Model

Some constructs are considered to be "derived concrete syntax" (or a shorthand notation) for other equivalent concrete syntax constructs. For example, in SDL-2010 omitting an input for a signal is derived concrete syntax for an input for that signal followed by a null transition back to the same state.

The concrete syntax for a shorthand notation is allowed to be placed either in the Concrete grammar or in the Model.

The properties of a shorthand notation are derived from the way it is modelled in terms of (or transformed to) the primitive concepts. In order to ensure easy and unambiguous use of the shorthand notations, and to reduce side effects when several shorthand notations are combined, these concepts are transformed in a specified order.

The result of the transformation of a fragment of derived concrete syntax is usually either another fragment of derived concrete syntax, or a fragment of concrete syntax. The result of the transformation may also be empty. In the latter case, the original is removed from the specification.

Transformations can be inter-dependent and therefore the order in which various transformations are applied determines the validity and meaning of a definition.

### 5.4 Metalanguages

For the definition of properties and syntaxes of ITU-T languages, different metalanguages are used according to the particular needs.
In the following, an introduction of the metalanguages used is given.

### 5.4.1 Metalanguage for the abstract grammar

The abstract grammar of the language defines the relationships between elements of the language, without being concerned with issues such as punctuation marks that are needed to separate or terminate concrete syntax elements. There are two notations defined for the abstract grammar: a textual abstract syntax and a graphical metamodel notation.

One notation may not cover exactly everything found in the other notation, but the intention is that concrete grammar elements that map to the same abstract syntax elements have the same semantics.

### 5.4.1.1 Textual presentation

The following informally describes the textual presentation of the abstract syntax of ITU-T languages.

A definition rule in the abstract syntax can be regarded as a named composition defining a list of sub-components. The name starts with a letter (by convention uppercase) and is followed by any number of (by convention lowercase) letters and single hyphens ending in a (by convention lowercase) letter. When a name is used on the right-hand-side of a rule, it is allowed to be immediately followed by suffixes as described below.

To avoid confusion with the suffix "-set", a name shall not end in "-set". To avoid confusion with Quotation, a name should not be in all uppercase letters. The case of names is significant, so that My-item is distinct from My-Item, but the use of two names that differ only in the case of letters should be avoided, so that tools can be used to correct errors.

By convention names are in italics, which allows the rules and component names of the abstract syntax to be more easily recognized in the context of the other text and the name to be distinguished when used in plain text. The suffix "-set" by convention has the word "set" in bold.

A rule starts with the name of the rule followed by either "::" or " $=$ ". The rule continues on the same line and subsequent lines until followed by another rule or any printing character not allowed within a rule (such as a comma "," or a semi-colon ";"), or the end of the text containing the rules (for example, the end of a file, end of a text box, a heading, or a paragraph style not allowed for rules). Non-printing characters in rules are separators and are otherwise ignored, but by convention names of rules are always placed at the start of a new line, and continuation lines for a rule start with a nonprinting character (usually a space or tab character). The characters that are allowed within a rule are ":" (part of "::"), "=", "*", "+", "[", "]", "|", "\{", "\}", "." (as part of ".." or "..."), "-", digits (0, 1, 2, 3, $4,5,6,7,8,9$ ) and (uppercase and lowercase) letters (A, B, C, D, E, F, G, H, I, J, K, L, M, N, O, P, Q, R, S, T, U, V, W, X, Y, Z, a, b, c, d, e, f, g, h, i, j, k, l, m, n, o, p, q, r, s, t, u, v, w, x, y, z).
The symbol "::" can be read as "is defined as" and is used to define a domain in a definition rule.
For example:
Channel-path :: Originating-gate
Destination-gate
Signal-identifier-set
which defines the domain for the composition named Channel-path. This consists of three sub-components, which in turn are compositions or elementary domains (see below). It is allowed for a sub-component to contain the domain being defined, but (to avoid an infinitely recursive definition) there shall be at least one alternative of the domain definition that does not include such a sub-component.

It is allowed that a domain has no sub-components, and this is denoted by a " $\}$ " as the right-hand-side of a definition rule. It is the presence (or absence) of an instance of such a domain that is still significant, and an instance of such a domain is distinct from an instance of any other domain even if the other domain does not have any components.
The symbol " $=$ " can be read as "is equivalent to" and is used to define an equivalence rule where the name on the left-hand-side is equivalent to the syntactic expression on the right-hand-side. Wherever the equivalence name defined by the left-hand-side occurs, syntactically it can be replaced by the expression on the right-hand-side, grouped if the right-hand-side has multiple components. The reason for using " $=$ " is to avoid repeating a common grouping especially in the case of alternatives, or to give a more specific name to a domain in a particular context to aid description of the language.

## The definition

Agent-identifier $=$ Identifier
expresses that an Agent-identifier is equivalent to an Identifier and therefore cannot (syntactically in the abstract syntax) be distinguished from other Identifier items defined in the same way such as a Signal-identifier where:

Signal-identifier $=$ Identifier
The distinction between items defined as equivalent is not syntactic. For example, each Identifier instance has a value, which allows one instance to be distinguished from another and whether the value is the Identifier for a signal (Signal-identifier) or agent (Agent-identifier) in the model being defined. The name Agent-identifier is preferable wherever the Identifier has to identify an agent. The constraint (on the Identifier to identify an agent in this case) may be implied by general conventions or could be given by natural language and optionally by formal expression (in a language such as OCL) An alternative to naming equivalences is to name components of a domain (as described below).

The definition:
Agent-qualifier :: Agent-name
expresses that an Agent-qualifier is a domain that has a component that is an Agent-name and Agent-qualifier is distinct from the domain for Agent-name. For example, if:

| Agent-name | $=\quad$ Name |
| :--- | :--- |
| Interface-name | $=\quad$ Name |

then an Agent-qualifier can be distinguished syntactically from an Agent-name and other items that are in the same domain as Agent-name such as an Interface-name.

The definition:

## Nextstate-node $=$ Dash-nextstate | Named-nextstate

expresses that a Nextstate-node is equivalent to the alternatives on the right-hand-side. Therefore syntactically,

Terminator

$$
\begin{array}{ll}
:: \quad \text { Nextstate-node } \\
& \text { | Stop-node } \\
& \text {... }
\end{array}
$$

has the same meaning as:
Terminator $:: \quad\{$ Dash-nextstate $\mid$ Named-nextstate $\}$
|Stop-node
An abstract syntax item might also be of some elementary (non-composite) domain. Predefined elementary domains are:
a) Natural

- This is the elementary domain of non-negative integers, denoted by Nat.
- Example:
- Number-of-instances :: Nat [Nat]
- Number-of-instances denotes a composite domain containing one mandatory natural (Nat) value and one optional natural ([Nat]) value denoting respectively the initial number and the optional maximum number of instances.
b) Boolean
- Boolean denotes the domain of the two Boolean values TRUE and FALSE.
c) Quotation
- These are represented as any bold face sequence of uppercase letters and digits. There can be a single quotation or a number of grouped quotations separated by " $\mid$ ", and in both cases this defines an enumerated domain with as many values as there are quotations. For example,


## Agent-kind $=\quad$ SYSTEM $\mid$ BLOCK $\mid$ PROCESS

- The quotation defines an enumerated domain with three values, represented by SYSTEM, BLOCK and PROCESS. Usually the names used in the abstract syntax correspond to the names in the concrete syntax, in this case the keywords SYSTEM, BLOCK and PROCESS. In this example, the name Agent-kind is made equivalent to the enumerated domain. If the quotation group was a component of a domain with other components (see example Visibility-name below), the enumeration domain has an implicit name only and cannot be reused in other contexts, and another use of the same quotations defines another implicit domain.
- A single quotation corresponds to a domain with just one value, and will usually appear in an option. In such cases, it is usually preferable to replace the quotation with a Boolean.

Example:
Channel-definition :: Channel-name
[NODELAY]
Channel-path-set

- A channel definition has the property that it does not delay distinguished by an optional quotation NODELAY. This could be replaced by:
$\left.\begin{array}{lll}\text { Channel-definition }:: & \begin{array}{l}\text { Channel-name } \\ \\ \text { Is-delaying }\end{array} \\ \text { Channel-path-set }\end{array}\right\}$
- $\quad$ Similarly if there are only two alternatives, this can usually be replaced by a Boolean. For example:

Category-definition :: Category-name

## \{ ABSTRACT | CONCRETE \} <br> Category-details

- could be replaced by:

Category-definition $:: \quad$| Category-name |
| :--- |
|  |
|  |
| Is-concrete |

Category-details

$I s$-concrete $=\quad$| Boolean |
| :--- |

d) Token

Token denotes the domain of tokens. This domain can be considered to consist of a potentially infinite set of distinct atomic objects for which no representation is required.
Example:
Name :: Token
A name consists of an atomic object such that any Name can be distinguished from any other name.
e) Unspecified items

- An unspecified item denotes domains which might have some representation, but for which the representation is of no concern in the language Recommendation.
Example:
Informal-text :: ...
- Informal-text contains an object that does not have a formally defined content and structure and therefore is not formally interpreted. The meaning is therefore not formally defined, and if interpretation of the model leads to interpretation of the informal text, then the further behaviour of the system is not defined.

The following operators (constructors) of BNF (see clause 5.4.2) have the same use in the abstract syntax:

- suffix "[n..m]" for a list of between n and m items, where n is a non-negative integer, m is a positive integer $>=\mathrm{n}$;
- suffix "[n]" for a list of exactly $n$ items, where $n$ is a positive integer - equivalent to the suffix [n..n];
- suffix "[n..*]" for a list of at least n items, where n is a positive integer;
- suffix "*" for a possibly empty list, which is equivalent to the suffix [0..*];
- suffix "+" for a non-empty list, which is equivalent to the suffix [1..*];
- separator "|" for an alternative;
- brackets "[" "]" for meaning the enclosed item or items are optional.

Parentheses "\{" and "\}" are used for grouping of items that are logically related.
All the operators (including "-set" and "+set" optionally followed by the number of items - see below) are applied to either a domain name or to a grouping (within "\{" and "\}").
A domain is allowed to be specified to have a structure derived from another domain, in which case we say the former (child) domain inherits its structure from the latter (or parent) domain. A child domain has all the sub-components defined by the parent domain, and it is allowed to add subcomponents which follow the inherited components. A child item object is allowed in any place of the tree of objects where parent item objects are allowed.
For example, the definition:
Agent-identifier (Agent-name) :: Qualifier
defines an Agent-identifier to have the same structure as Agent-name but has another component that is a Qualifier. As far as the structure is concerned, the above definition is the same as defining:

## Agent-identifier :: Agent-name Qualifier

However, the former inheritance definition means that an Agent-identifier instance is allowed wherever an Agent-name instance is allowed, so rather than specify the alternatives \{ Agent-identifier |Agent-name \}, a single Agent-name is sufficient.
Finally, the abstract syntax uses two other postfix operators "-set" yielding a possibly empty set (unordered collection of distinct objects) and the postfix operator "+set" yields a set which shall not be empty. These can also have the suffixes " $[\mathrm{n}]$ " for a set of exactly n items, suffix "[n..m]" for a set of between n and m items, suffix "[ $\left.\mathrm{n} . .{ }^{*}\right]$ " for a set of at least n items.
Example:
Agent-graph :: Agent-start-node State-node-set
An Agent-graph consists of an Agent-start-node and a possibly empty set of State-nodes.
Where a sub-component is of an elementary domain, or a quotation, a default value can be specified for this sub-component. For example,
Number-of-instances :: Nat:=1 [Nat]
specifies that the first Nat component will have the value 1, if no value is specifically given in the concrete notation.
Visibility-name :: Name \{ PUBLIC | PROTECTED | PRIVATE := PUBLIC \}
specifies that the second component will have the Quotation value PUBLIC, if no value is specifically given in the concrete notation.

Suffixes that specify multiplicity for a list or set specify constraints on the grammar if the lower bound is greater than zero or if an upper bound is given (that is, it is not "*"). Often there are other constraints that need to be specified: for example, in Number-of-instances if there is a second Nat, this represents the maximum number of instances and has to be greater than or equal to the initial number of instances given by the first Nat. A natural language description of such constraints helps explain the reason for the limitation, but writing natural language so that the formulation of the constraint is clear and unambiguous can be difficult, especially if the text is to be translated into different natural languages. For that reason, it is desirable that constraints on the grammar are written in a formal notation (such as OCL). To be able to clearly refer to the domain components of a domain, each domain component can optionally be given a name starting with a lowercase letter preceding the component domain by a colon, for example:

```
Number-of-instances :: initialNumber: Nat:=1 [maximumNumber: Nat]
```

In this case, it is now possible to write Number-of-instances.maximumNumber >= Number-ofinstances.initialNumber. An alternative to named domain components is to introduce an equivalence rules for each component. For example:

| Number-of-instances | $:: \quad$ Initial-number [Maximum-number] |
| :--- | :--- | :--- |
| Initial-number | $=\quad$ Nat $:=1$ |
| Maximum-number | $=\quad$ Nat |

The main differences are the relative conciseness of the named component notation, and that a component name is local to a domain whereas an equivalence name can be used in several domains.

The following gives a more formal abstract description of the textual abstract grammar defined in the notation itself as far as possible and supplementary to the text above:

Grammar
Rule
Definition-rule
Equivalence-rule
Expression
] Domain
Alternatives
Composition
Option
List
Set
Domain

Non-terminal-domain
Elementary-domain
Elementary-domain-kind
Default-elementary-value
Enumeration-domain
:: Rule-set
:: Definition-rule | Equivalence-rule
:: Name [ Non-terminal-domain ] [Expression]
:: Name Expression
:: Alternatives $\mid$ Composition $\mid$ Option |List $\mid$ Set $\mid$ [ Component-name
:: Expression[2..*]
:: Expression[1..*]
:: Expression
:: Expression minimum: Nat: $=0$ [maximum: Nat $]$
:: Expression minimum: Nat:=0 [maximum: Nat ]
:: Non-terminal-domain
| Elementary-domain
| Enumeration-domain
:: Name
:: Elementary-domain-kind [ Default-elementary-value ]
$=$ NAT $\mid$ TOKEN $\mid$ BOOLEAN $\mid$ UNSPECIFIED
= Token
:: Quotation-set [ Default-enumeration-value ]

| Default-enumeration-value | $::$ Quotation |
| :--- | :--- |
| Component-name | $=$ Token |
| Name | $=$ Token |

There should normally be one Rule from which every other Rule can be reached.
The same character string always maps to the same Token value, and two different character strings map to different Token values. The character strings Token, Nat, Boolean and Informal-text shall not be used for a Name, because these are reserved for each of the respective domains.

A Name is a Token that is represented by a character string starting with a letter, followed by any number of alphanumeric characters and single hyphens and terminating in an alphanumeric. Each Name of a Definition-rule or Equivalence-rule shall have a Token value distinct from the Token value of the Name of any other Definition-rule or Equivalence-rule.

A Non-terminal-domain has a Token that is represented by name starting with a letter, followed by any number of alphanumeric characters and single hyphens and terminating in an alphanumeric. The Name of a Non-terminal-domain shall have the same Token value as the Name of a Definition-rule or Equivalence-rule.
NAT, TOKEN, BOOLEAN and UNSPECIFIED are represented by Token, Nat, Boolean and Informal-text respectively.

A Default-elementary-value is a Token that is represented in the case of Nat (NAT) by a non-negative integer (a string of digits) and in the case of Boolean (BOOLEAN) by one of the strings TRUE or FALSE. If the Elementary-domain-kind of an Elementary-domain is Token or Informal-text, then assignment is not allowed and the Default-enumeration-value shall be omitted. If the Elementary-domain-kind of an Elementary-domain is Nat, then Default-enumeration-value shall be a Token for an integer. If the Elementary-domain-kind of an Elementary-domain is Boolean, then Default-enumeration-value shall be a Token for a Boolean value.
Quotation is represented by bold face sequence of uppercase letters and digits, called a quotation in this paragraph. The same quotation always maps to the same Quotation value and two different quotations map to different Quotation values. The quotation for a Quotation shall not be the same as the character string for a Token. An Enumeration-domain is represented either by a single quotation or by a list of quotations separated by "|" characters. When a "|" character is between two quotations, it is a separator for an Enumeration-domain rather than for Alternatives. The Default-enumerationvalue (if present) of an Enumeration-domain shall be one of the Quotation values in the Quotation-set of the Enumeration-domain.

Alternatives or Composition within an Expression that is Alternatives or Composition or List or Set are grouped by "\{" and "\}" brackets in the concrete notation.

The minimum Nat for a List or Set is the minimum number of elements, and the maximum Nat is the maximum number. If a maximum exists, it must be greater than or equal to the minimum. If maximum is omitted, there is no maximum number of elements. The notation for a list or set is explained above.

### 5.4.1.2 Metamodel presentation

The following describes an alternative to the textual presentation introduced in clause 5.4.1.1. This metamodel representation of the abstract syntax of ITU-T languages is explained using a meta-metamodel (Figure 1).

Z.111(16)_F01

Figure 1 - Meta-metamodel for metamodeling language elements
The meta-metamodel elements described below are represented using UML in Figure 1, except where otherwise noted. Note that, in order to better align metamodels based on ITU-T Z. 111 with MOF and UML, Token elements are shown with the "String" keyword in metamodels, including the metametamodel in Figure 1.
In the following, a Named element is a meta-class that contains a name attribute of the Token primitive type (shown visually as String). Some Token values represent Nat values (that is non-negative Integers), and two distinct Token values represent the Boolean true and false values. These Token values for Nat and Boolean shall only be used in the value of a DefaultValue and shall not be used for the Token value of any name. When deriving the Token from a character string in the concrete notation, the Token value depends on the character string and the case of letters in the character string are taken into account. The same character string in the concrete notation for names shall always produce the same Token, so that constraints on Token values are constraints on character string in the concrete notation for the name attributes. Two name attributes are the same if they have the same Token value, otherwise they are distinct.
i) Class is a Named element that contains one or more Attribute items and can participate in binary Association relations via one or both of the AssociationEnd attributes of each Association. A Class can inherit the Attribute and AssociationEnd (and hence Association) items from another Class (single inheritance only).

- Each Class shall have a name with a Token value that is distinct from the Token value for the name of any other Class (including predefined primitives), or any Enumeration.
- The inheritedFrom attribute (if present) of Class references the parent Class of a Class. An inherited Class of a Class is either the parent Class or any inherited Class of the parent Class, transitively. A Class shall not have itself as an inherited Class. The attribute list of a Class
includes the attribute list of the parent Class, so that the same Attribute shall be referenced by each element of the attribute list of the parent Class and the corresponding (by order) element of the attribute list of the Class. A Class that inherits from a parent Class includes each end reference to an AssociationEnd of parent Class and is the endClass of the AssociationEnd. A Class that inherits from a parent Class is referenced as an endType by each AssociationEnd that references the parent Class as an endType. Multiple classification is not allowed, that is, an instance of a Class C 1 cannot be an instance of another Class C2, if C 1 is not inherited (directly or transitively) from C2.
- $\quad$ The name (if present) of each AssociationEnd that has an end in a given Class shall be distinct from the Token value for the name (if present) of any AssociationEnd that has an end in the Class or any inherited Class.
ii) Attribute is a Named element. It is a meta-class for an ordered element of a Class that has a DataType and an optional DefaultValue.
- $\quad$ The name of each attribute of a given Class shall be distinct from the name of any other attribute of the Class or any inherited Class and shall also be distinct from the name (if present) of any AssociationEnd that has an end in the Class or any inherited Class.
- $\quad$ The multiplicity of an Attribute is given by multiplicityMin and multiplicityMax, which have default values of 1 if they are omitted in concrete notation. However, if in the concrete notation the maximum is given by "*", the maximum is empty and there is no upperbound on the number of Attribute instances. If multiplicityMax is present, it shall be greater than 0 and also greater than or equal to multiplicityMin.
- If the Attribute has a DataType that is a PrimitiveType Token or PrimitiveType Unspecified, the DefaultValue shall be absent. If a DefaultValue is present, its Token value shall represent a value of the DataType of the Attribute (that is, it shall be true or false for a Boolean, a non-negative integer for a Nat, or a quotation Literal for an Enumeration).
iii) DataType is a meta-class that is either a PrimitiveType or an EnumeratedType.
iv) PrimitiveType is a DataType where the type is predefined (one of Token, Nat, Boolean, Unspecified) primitive identified by the Primitive value of the name attribute of the PrimitiveType. Unspecified is a primitive that might have some representation, but for which the representation is of no concern in the language Recommendation.
v) EnumeratedType is a DataType where the type is an Enumeration identified by its name attribute
vi) Enumeration is a Named element that represents a (quotation) DataType whose values are represented by a set of literal elements.
- $\quad$ Each Enumeration shall have a name with a Token value that is distinct from the Token value for the name of any other Enumeration or any Class (including predefined primitives).
vii) Literal is a Named element contained in an Enumeration that is one of the values of an Enumeration. The name of each literal of an Enumeration shall be distinct for any other literal of the Enumeration, but it is permitted for the name to be the same as the literal of another Enumeration.
viii) DefaultValue is a meta-class for the optional component of an Attribute that identifies a default value of the appropriate type (see Attribute).
ix) Association is a meta-class for the relation between Class meta-classes. It has two AssociationEnd attributes (source and target) for the logical connection of the Association with the related Class meta-classes.
- An Association shall have at least one AssociationEnd that is navigable (isNavigable $=$ true $)$. An Association is represented by a line in the concrete notation. An Association navigable in only one direction is represented by a line with an arrow on the navigable end.
x) AssociationEnd is a meta-class for the end of an Association. It has Boolean attributes that determine if the AssociationEnd is composite (isComposite), navigable (isNavigable) or ordered (isOrdered). If the AssociationEnd is navigable (isNavigable $=$ true), the name shall not be empty, and the name (which corresponds to the role name of the AssociationEnd) shall be distinct from the name of any other AssociationEnd that is an end of the Class that is the source of the AssociationEnd.
- An AssociationEnd shall not be both the source and target of an Association. In the following, an AssociationEnd that is the source of an Association is called a source AssociationEnd, and an AssociationEnd that is the target of an Association is called a target AssociationEnd.
- $\quad$ The Class of the endType of a source AssociationEnd is derived and is the same as the Class of the endClass of the target AssociationEnd of the Association that is the source of the source AssociationEnd. The Class of the endType of a target AssociationEnd is also derived and is the same as the Class of the endClass of the source AssociationEnd of the Association that is the target of the target AssociationEnd. More formally:
- source AssociationEnd.endType = AssociationEnd.associationlink.target.endClass; and
- target AssociationEnd.endType = AssociationEnd.associationlink.source.endClass.
- A source AssociationEnd is allowed to be composite (isComposite $=$ true), and in this case, the Class with the source AssociationEnd as an end has as a component an item (or a collection of items - see multiplicity) with the endType of the source AssociationEnd. A Class with a non-composite source AssociationEnd is related (by the Association) to an item or collection with the endType of the source AssociationEnd. In the concrete notation, an AssociationEnd that is composite is represented by filled diamond at the AssociationEnd that is the source.
- The AssociationEnd that is the target of an Association shall not be composite (isComposite $=$ false).
- An AssociationEnd has a multiplicity defined by the values of the multiplicityMin and multiplicityMax, which are natural numbers. Although default values are given for the values, it is not allowed to omit the values in the concrete notation. If, in the concrete notation, the maximum is given by "*", the maximum is empty and there is no upperbound on the number of instances. If multiplicityMax is present, it shall be greater than 0 and also greater than or equal to multiplicityMin. If multiplicityMax is equal to multiplicityMin, then it is sufficient in the metamodel representation to show only one number (e.g., 1 instead of $1 . .1$ ).
- If multiplicityMax is larger than 1 or omitted, the AssociationEnd denotes a collection of items. By default, this collection is ordered (isOrdered $=$ true), so the collection is a list. If the collection is not ordered (isOrdered $=$ false), the collection will normally be a set (but is allowed to be a bag). In a set, each item in the collection has a different value. In a bag, there may be any number of items with the same value.

Several modelling elements commonly used in class diagrams should be avoided when describing metamodels for abstract syntaxes. These include:

- Packages: A structuring mechanism that does not add semantic value at the abstract level.
- $\quad$ Visibility of attributes and association ends: This has no impact on the abstract model as it is hidden.
- Absence of association multiplicities: To avoid ambiguities in the understanding of default multiplicities (according to standards or people), multiplicities at association ends are required to be explicit.
- Multiple inheritance: May require name resolution and hence should be avoided.
- Operations: Not usually needed in the abstract model. Operations on model elements (in a language such as OCL) are useful to describe constraints formally, but constraints should also be expressed in natural language as well as formally so the intent can be captured.
- Abstract classes: Not needed because the mapping between a concrete syntax and an abstract syntax can prevent the instantiation of such classes.
- Interfaces: Not needed since there are no operations on instances.
- Do not include layout elements (graphical information such as colour, positions, shapes, sizes).


### 5.4.2 Metalanguage for the concrete grammar

The syntax of the concrete grammar is defined by the extended Backus-Naur Form (BNF) defined in more detail in this clause. If necessary, syntax is supplemented by further definition in natural language and preferably also in a formal language such as OCL. The concrete syntax is mapped to the abstract syntax (either implicitly as described in clause 5.3.2 b) Concrete Grammar above or explicitly) and the constraints and semantics of the abstract grammar apply.

The grammar should usually be separated into two levels:

1) lexical rules that define the syntax of lexical units such as names, numbers, character strings, composite symbols (such as "=>") and any reserved keyword names of the language;
2) non-lexical rules that use the lexical rules as terminal symbols.

If the grammar is simple, it is acceptable to include the lexical rules with the non-lexical rules. The definition of graphical symbols (such as a frame box or a flow line) are treated like keywords and defined in the lexical rules, or as part of the non-lexical grammar.

In the BNF for lexical rules, the terminals are <space> and the printed characters are specified as terminal symbols by the defined language. The printed characters are not limited to the Latin alphabet; the universal multiple-octet coded character set defined in [ITU-T T.55] is allowed to be used. The actual printed character set used and non-printing characters such as <space> are defined in the formal language Recommendation. Some examples of lexical rules are:

```
<asterisk> ::= *
<bit string> ::=
```

$$
\text { <apostrophe> }\{0 \mid 1\}^{*} \text { <apostrophe> }\{\mathrm{B} \mid \mathrm{b}\}
$$

NOTE - It is assumed that characters are ordered and presented left to right then top to bottom as in Latin scripts. If other orderings (such as right to left, or top to bottom, mixed ordering) are significant, the relationship between the order characters is analysed and how they are presented has to be defined.
In the Backus-Naur Form for non-lexical rules, a terminal symbol is one of the lexical units defined to be terminal. In non-lexical rules, a terminal can be represented by one of the following:
a) a keyword (such as state);
b) the character for the lexical unit if it consists of a single character (such as " $=$ " );
c) the lexical unit name (such as <quoted operation name> or <bit string>);
d) the name of a <composite special> lexical unit (such as <implies sign> which is defined as "=>").

To avoid confusion with BNF grammar, the lexical unit names <asterisk>, <plus sign>, <vertical line>, <left square bracket>, <right square bracket>, <left curly bracket> and <right curly bracket> are always used rather than the equivalent characters (except in the lexical rule that defines the unit - see the example of <asterisk> above). Note that non-lexical terminals that have multiple instances (typically the lexical units for <name> and <character string>) are permitted to have semantics stressed as defined below.

The angle brackets and enclosed word(s) are either a non-terminal symbol or one of the lexical units. Syntactic categories are the non-terminals indicated by one or more words enclosed between angle brackets. For each non-terminal symbol, a production rule is given in the concrete grammar. For example,
<block reference> ::=
block <block name> referenced <end>
A production rule for a non-terminal symbol consists of the non-terminal symbol at the left-hand side of the symbol "::=", and one or more constructs, consisting of non-terminal and/or terminal symbol(s) at the right-hand side. For example, <block reference>, <block name> and <end> in the example above are non-terminals; block and referenced are terminal symbols. By convention, a rule always starts on a new line and continuation lines start with a non-printing character (usually a space or tab). The rule continues on the same line and subsequent lines until followed by the start of another rule (a line starting with non-terminal symbol followed by "::=") or the end of the text containing rules (indicated by text style, a heading, characters not allowed in rules or some other means).

Sometimes the symbol includes an underlined part. This underlined part stresses a semantic aspect of that symbol. For example, <block name> is syntactically identical to <name>, but semantically it requires the name to be a block name.

At the right-hand side of the $"::="$ symbol, several alternative productions for the non-terminal can be given, separated by vertical bars ("|" ). For example,
<diagram in package> ::=

> <package diagram>
> <package reference area>
> <entity in agent diagram>
> <data type reference area>
> <signal reference area>
> <procedure reference area>
> <interface reference area>
> <create line area>
> <option area>
expresses that a <diagram in package> is a <package diagram>, or a <package reference area>, or an <entity in agent diagram>, or a <data type reference area>, or a <signal reference area>, or a <procedure reference area>, or an <interface reference area>, or a <create line area> or an <option area>.
Syntactic elements may be grouped together by using curly brackets ("\{" and "\}"), similar to the parentheses in the textual presentation of abstract syntax (see clause 5.4.1.1). A curly bracketed group may contain one or more vertical bars, indicating alternative syntactic elements.
Repetition of syntactic elements or curly bracketed groups is indicated by:

- $\quad$ suffix "[n..m]" the group is repeated between $n$ and $m$ items, where $n$ is an integer $>=0, m$ is a positive integer $>=\mathrm{n}$;
- suffix "[n]" the group is repeated exactly n items, where n is a positive integer - equivalent to the suffix [n..n];
- suffix "[n..*]" the group is repeated at least n items, where n is a positive integer;
- suffix "*" the group is optional and repetition of any number of times is allowed - equivalent to the suffix [0..*];
- suffix " + " the group shall be present and repetition of any number of times is allowed equivalent to the suffix [1..*].

For example,
<operation definitions> ::=
$\left\{\begin{array}{l}\text { <operation definition> } \\ \text { <operation reference> } \\ \text { <external operation definition>> }\end{array}\right\}+$
The example above expresses that <operation definitions> may contain zero or more definitions of <operation definition> or <operation reference> or <external operation definition>, and may contain more than one of any of these.
A repetition suffix can include a list separator after the repetition. If the repetition uses square brackets ("[" and "]"), a vertical bar ("|" ) is placed before the closing square bracket ("]") and the separator is placed between the vertical bar ("|") and the closing square bracket ("]"). If the repetition suffix is an asterisk ("*") or plus sign ("+"), it is followed by an opening square bracket and vertical line pair ("[|") the separator and then a closing square bracket ("]"). For example,
<two or more declarations> ::=

$$
\{\quad<\text { name }>+[\mid,]: \text { <type> }\}[2, * \mid ;]
$$

The example above expresses that <declaration list> contains at least two and possibly many repetitions separated by semicolons of the sequence: <name> list separated by commas, colon and <type>. This is equivalent to the following syntax that uses plain repetition (without separators) where <name> is repeated:
<two or more declarations> ::=

```
    <name> {, <name> }*:<type>
{; <name> {, <name> }*:<type> }+
```

This can also be expressed using recursion instead of repetition as:

```
<two or more declarations> ::=
                            <name list> : <type> ; { <name list> : <type> | < two or more declarations>
}
<name list> ::=
```

```
{ <name> | <name>, <name list> }
```

If syntactic elements are grouped using square brackets ("[" and "]"), then the group is optional. For example,
<valid input signal set> ::=
signalset [<signal list>] <end>
expresses that a <valid input signal set> may, but need not, contain <signal list>. An optional grouping is equivalent to a repetition suffix [0..1].
If there is any ambiguity between an optional group and the use of square brackets ("[" and "]") for a repetition suffix, it is a repetition unless there is at least one layout character (such as a space) before the "[" in which case it is an optional group.
To support the graphical grammar, the metalanguage has the following metasymbols:
a) set
b) contains
c) is associated with
d) is followed by
e) is connected to
f) is attached to

The set metasymbol is a postfix operator operating on the immediately preceding syntactic elements within curly brackets, and indicating an (unordered) set of items. Each item may be any group of syntactic elements, in which case it must be expanded before applying the set metasymbol.

Example:
\{ <operation text area>* <operation body area> \} set
is a set consisting of zero or more <operation text area>s, and one <operation body area>. The set metasymbol is used when the position of the syntactic elements relative to one another in the diagram is irrelevant and the elements can be considered in any order.

All the other metasymbols are infix operators, having a graphical non-terminal symbol as the left-hand argument. The right-hand argument is either a group of syntactic elements within curly brackets or a single syntactic element. If the right-hand side of a production rule has a graphical non-terminal symbol as the first element and contains one or more of these infix operators, then the graphical non-terminal symbol is the left-hand argument of each of these infix operators. A graphical non-terminal symbol is a non-terminal ending with the word "symbol".

The metasymbol contains indicates that its right-hand argument should be placed within its left-hand argument and the attached <text extension symbol>, if any. The right-hand argument is expanded within the symbol, should not cross the symbol boundaries and is distinct from any occurrence of the same syntax in another rule. For example,
<package use area> ::= <text symbol> contains <package use clause>
<text symbol> ::=

means the following


The metasymbol is associated with indicates that its right-hand argument is logically associated with its left-hand argument (as if it were "contained" in that argument, the unambiguous association is ensured by appropriate drawing rules). The right-hand argument is expanded and is distinct from any occurrence of the same syntax in another rule.

The metasymbol is followed by means that its right-hand argument follows (both logically and in drawing) its left-hand argument. The right-hand argument is expanded at the end of the implied symbol, and is distinct from any occurrence of the same syntax in another rule.

The metasymbol is connected to means that its right-hand argument is connected (both logically and in drawing) to its left-hand argument. The right-hand argument is expanded, and is distinct from any occurrence of the same syntax in another rule (in contrast to is attached to below).

The metasymbol is attached to expresses syntax requirements but not syntax productions. The metasymbol is attached to requires its right-hand argument and left-hand argument be attached to each other (both logically and in drawing), but one argument is not expanded with the syntax for the other argument, but each shall exist as separate expansions from syntax rules (in contrast to is
connected to above). Being attached is mutual, so that A is attached to B is always matched in the syntax by another rule where B is attached to A , though this need not be directly expressed on B. For example, B may have alternatives B1 and B2 each of which is attached to A. Being attached will usually mean that the abstract syntax for each side contains the identifier of the other side.

## Annex A

## Use of graphical meta-symbols

(This annex forms an integral part of this Recommendation.)

The meta symbols for the notation of the graphical grammar should always be used according to their definition.

- It is required that the operators (contains, is associated with, is followed by, is connected to, is attached to) have a graphical non terminal as the left hand argument, that is a nonterminal ending with the word "symbol".
- The left-hand-side should never be empty.
- The right-hand-side is allowed to be a set that is possibly empty, such as the form ... is connected to $\left\{<\right.$ gate $\left.>^{*}\right\}$ set ; that is connected to an unordered (logically and graphically), possibly empty set of <gates>.
- The post-fix operator set should be applied consistently. The form $\{\langle x x x\rangle *\}$ set is an unordered (logically and graphically/textually), possibly empty set of <xxx>. The form $\{\langle x x x\rangle *\}$ set should be used rather than $\{\{\langle x x x\rangle *\}$ set $\}$ or $\{\langle x x x\rangle\}^{*}$ set as right hand side of a graphical meta operator, where the meaning that the right hand side is an unordered set of $\langle x x x\rangle$. In particular $\{\langle x x x\rangle\}^{*}$ set means an ordered list of $\langle x x x\rangle$ that is in a set with just one element. In $\left\{\left\{\langle x x\rangle^{*}\right\}\right.$ set $\}$ the outer braces are (probably) redundant. The form $\{\{\langle x x x>*\}<y y>\}$ set means a set consisting of an ordered list of <xxx> items with <yy> required in the set but with the placing unordered: that is, either a list of $\langle x x>\rangle$ items followed by <yy>, or <yy> followed by a list of <xxx> items. The form $\{\{\langle x x x\rangle *\}$ set <yy> \}set is a set containing any number of unordered <xxx> items and one <yy> item.
- If set is in a graphical context the items are unordered vertically or horizontally but should not overlap.
- If the right hand side of an operator is $\{\langle x x\rangle\langle y y\rangle\langle z z\rangle\}$ this is an ordered list $\langle x x\rangle\langle y y\rangle\langle z z\rangle$ and $\{<x x><y y><z z>\}$ set should be used if <xx> and <yy> and <zz> are graphical \{<xx>*\} is an ordered list of $\langle\mathrm{xx}\rangle$ and $\{\langle\mathrm{xx}\rangle *\}$ set should be used if $\langle\mathrm{xx}\rangle$ is graphical
- For a line symbol that connects two endpoints the form <xxx line symbol> is connected to \{<end1>\} is connected to \{<end2>\} should be used where possible, though the same meaning is given by<xxx line symbol> is connected to $\{<$ end $1><e n d 2>\}$ set.


## Appendix I

## Mapping from the metamodel presentation to the textual presentation: High-level description

(This appendix does not form an integral part of this Recommendation.)
The metamodel presentation and the BNF-based textual presentation of the abstract syntax each have benefits and drawbacks. Metamodels are appealing due to their graphical nature (where associations and inherited concepts are explicit), whereas BNF grammars are easily analysable by tools and there is considerable experience in their use. A procedure is outlined here that converts an abstract syntax expressed as a metamodel into an abstract syntax expressed textually. Leveraging this procedure, an abstract syntax can be expressed as a metamodel for convenience, but can be subject to validation through grammar-oriented tools.

The following gives the mapping from the concepts in the metamodel presentation of the abstract syntax to the textual representation of the abstract syntax.

A Class corresponds to a Definition-rule. The name corresponds to the Name. The class referenced by inheritsFrom corresponds to the Non-terminal-domain, if any. The set of attribute and association end corresponds to the Expression.
If the Attribute or AssociationEnd is named, this corresponds to an Equivalence-rule, where Name is derived from name and the current class name and the attribute or association end constructs the Expression of the alias as in the following, and name corresponds also to Symbol in the outer Expression. If the name is unique, it corresponds to Token. Otherwise, the Expression is constructed as follows.

If multiplicityMin of an Attribute or AssociationEnd is 0 and multiplicityMax of this Attribute or AssociationEnd is 1, and isNavigable is true, this corresponds to an Option, where Expression is a Symbol corresponding to the type of the attribute or association end.

If multiplicityMin of an Attribute or AssociationEnd is 0 , and multiplicityMax of this Attribute or AssociationEnd is *, and isOrdered of this Attribute or AssociationEnd is true, and isNavigable is true, this corresponds to a List, where Expression is a Symbol corresponding to the type of the attribute or association end.

If multiplicityMin of an Attribute or AssociationEnd is 1, and multiplicityMax of this Attribute or AssociationEnd is *, and isOrdered of this Attribute or AssociationEnd is true, and isNavigable is true, this corresponds to a List with a minimum of one element, where Expression is a Symbol corresponding to the type of the attribute or association end.
If multiplicityMin of an Attribute or AssociationEnd is 0 , and multiplicityMax of this Attribute or AssociationEnd is *, and isOrdered of this Attribute or AssociationEnd is false, and isNavigable is true, this corresponds to a Set, where Expression is a Symbol corresponding to the type of the attribute or association end.

If multiplicityMin of an Attribute or AssociationEnd is 1, and multiplicityMax of this Attribute or AssociationEnd is *, and isOrdered of this Attribute or AssociationEnd is false, and isNavigable is true, this corresponds to a Set with a minimum of one element where Expression is a Symbol corresponding to the type of the attribute or association end.

If multiplicityMin of an Attribute or AssociationEnd is 1 and multiplicityMax of this Attribute or AssociationEnd is 1, and isNavigable is true, this corresponds to a Symbol, such that the Symbol is corresponding to the type of the attribute or association end.
If isComposite of an AssociationEnd is false, then the Symbol is a Name; otherwise, the Symbol is a Non-terminal-domain.

If an Attribute has a default, then this generates an Assignment, where DefaultValue corresponds to Elementary-domain and the name of the type of the Attribute corresponds to Token.
Any multiplicityMin and multiplicityMax on an AssociationEnd where isNavigable is false and which is not both 1 in the case where isComposite is true or is otherwise 0 and ${ }^{*}$, respectively, generates a constraint on the respective element.

An Enumeration corresponds to a Non-terminal-domain where the name corresponds to Token, and Expression is an Alternative where each Symbol is an Elementary-domain corresponding to the literals.

## Bibliography

[b-ITU-T Z.101] Recommendation ITU-T Z. 101 (2016), Specification and Description Language: Basic SDL-22010. [http://www.itu.int/rec/T-REC-Z.101](http://www.itu.int/rec/T-REC-Z.101)
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[^0]:    For further details, please refer to the list of ITU-T Recommendations.

[^1]:    * To access the Recommendation, type the URL http://handle.itu.int/ in the address field of your web browser, followed by the Recommendation's unique ID. For example, http://handle.itu.int/11.1002/1000/1 1830-en.

