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SERIES Z: LANGUAGES AND GENERAL SOFTWARE ASPECTS FOR TELECOMMUNICATION SYSTEMS

Formal description techniques (FDT) – Specification and Description Language (SDL)

Specification and Description Language – Data and action language in SDL-2010

Recommendation ITU-T Z.104

International Telecommunication Union

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Recommendation ITU-T Z.104

Specification and Description Language – Data and action language in SDL-2010

Summary

Scope/Objective

This Recommendation defines the data features of the Specification and Description Language so that data definitions and expressions are well defined. Together with Recommendations ITU-T Z.100, ITU-T Z.101, ITU-T Z.102, ITU-T Z.103, ITU-T Z.105 and ITU-T Z.106, this Recommendation is part of a reference manual for the language. The language defined in this document partially overlaps features of the language included in Basic SDL-2010 in Recommendation ITU-T Z.101 and used in Comprehensive SDL-2010 in Recommendation ITU-T Z.102 and the features of Recommendation ITU-T Z.103.

Coverage

The Specification and Description Language has concepts for behaviour, data description and (particularly for larger systems) structuring. The basis of behaviour description is extended finite state machines communicating by messages. Data description is based on data types for values. The basis for structuring is hierarchical decomposition and type hierarchies. Though a distinctive feature of the Specification and Description Language is the graphical representation, the data and expression language is textual. This Recommendation covers the features of the language used to encode and decode data communicated by channels, define data types with values and operations and variables (including parameters) based on data types and expression actions that use the data types. This Recommendation does not always provide a canonical syntax, but by applying the *Model* descriptions given, a specification is transformed to Basic SDL-2010 defined in ITU-T Z.101 except in those cases where additional abstract syntax is added in this Recommendation.

Applications

The Specification and Description Language is applicable within standards bodies and industry. The main application areas for which the Specification and Description Language has been designed are stated in Recommendation ITU-T Z.100, but the language is generally suitable for describing reactive systems. The range of applications is from requirement description to implementation. The features of the language defined in Recommendation ITU-T Z.104 are essential for the data within a system.

History

Edition	Recommendation	Approval	Study Group
1.0	ITU-T Z.104	2004-10-07	17
2.0	ITU-T Z.104	2011-12-22	17

FOREWORD

The International Telecommunication Union (ITU) is the United Nations specialized agency in the field of telecommunications, information and communication technologies (ICTs). The ITU Telecommunication Standardization Sector (ITU-T) is a permanent organ of ITU. ITU-T is responsible for studying technical, operating and tariff questions and issuing Recommendations on them with a view to standardizing telecommunications on a worldwide basis.

The World Telecommunication Standardization Assembly (WTSA), which meets every four years, establishes the topics for study by the ITU-T study groups which, in turn, produce Recommendations on these topics.

The approval of ITU-T Recommendations is covered by the procedure laid down in WTSA Resolution 1.

In some areas of information technology which fall within ITU-T's purview, the necessary standards are prepared on a collaborative basis with ISO and IEC.

NOTE

In this Recommendation, the expression "Administration" is used for conciseness to indicate both a telecommunication administration and a recognized operating agency.

Compliance with this Recommendation is voluntary. However, the Recommendation may contain certain mandatory provisions (to ensure, e.g., interoperability or applicability) and compliance with the Recommendation is achieved when all of these mandatory provisions are met. The words "shall" or some other obligatory language such as "must" and the negative equivalents are used to express requirements. The use of such words does not suggest that compliance with the Recommendation is required of any party.

INTELLECTUAL PROPERTY RIGHTS

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As of the date of approval of this Recommendation, ITU had not received notice of intellectual property, protected by patents, which may be required to implement this Recommendation. However, implementers are cautioned that this may not represent the latest information and are therefore strongly urged to consult the TSB patent database at http://www.itu.int/ITU-T/ipr/.

Table of Contents

Introduction

Status/Stability

This Recommendation is part of the ITU-T Z.100 to ITU-T Z.106 series of Recommendations that give the complete language reference manual for SDL-2010. The text of this Recommendation is stable. For more details see Recommendation ITU-T Z.100.

Recommendation ITU-T Z.104

Specification and Description Language – Data and action language in SDL-2010

1 Scope and objective

This Recommendation defines the data and action language features of the Specification and Description Language. The features defined in this document include: encoding and decoding of data communicated over channels, the definition of data types (including interfaces), the definition of variables (including parameters), the definition and application of operations, the evaluation of expressions and assignment of values to variables. Without these features the language would be of limited use. For this reason, even very simple specifications in the Specification and Description Language use ITU-T Z.104 features. Together with [ITU-T Z.100], [ITU-T Z.101], [ITU-T Z.102], [ITU-T Z.103], [ITU-T Z.105] and [ITU-T Z.106], this Recommendation forms a reference manual for the language.

1.1 Objective

The objective of this Recommendation is to fully define the data related features of the Specification and Description Language. Some of the ITU-T Z.104 material is also in [ITU-T Z.100] and [ITU-T Z.101]. These definitions should be entirely consistent. The difference between the descriptions in [ITU-T Z.100] and [ITU-T Z.101] and the definition in Recommendation ITU-T Z.104 is that the former descriptions are overviews to enable data to be used with other features of SDL-2010, whereas the ITU-T Z.104 definition is intended to be complete (with the exception of definition of data types in ASN.1 modules – see [ITU-T Z.105]).

Recommendation ITU-T Z.104 also includes a mechanism for binding data definitions and actions expressed in the concrete syntax from other notations as described in Annex C to the abstract grammar and semantics of SDL-2010 and data binding in clause 7.2. By default the concrete grammar is as given in the main body of ITU-T Z.104 or [ITU-T Z.101], [ITU-T Z.102], [ITU-T Z.103], [ITU-T Z.105] and [ITU-T Z.106].

1.2 Application

This Recommendation is part of the reference manual for the Specification and Description Language.

The use of data is so fundamental, that even the simplest specifications need some data language features. Therefore, it is almost certain that some features defined in Recommendation ITU-T Z.104 will appear in any specification in SDL-2010.

2 References

The following ITU-T Recommendations and other references contain provisions which, through reference in this text, constitute provisions of this Recommendation. At the time of publication, the editions indicated were valid. All Recommendations and other references are subject to revision; users of this Recommendation are therefore encouraged to investigate the possibility of applying the most recent edition of the Recommendations and other references listed below. A list of the currently valid ITU-T Recommendations is regularly published. The reference to a document within this Recommendation does not give it, as a stand-alone document, the status of a Recommendation.

[ITU-T T.50] Recommendation ITU-T T.50 (1992), *International Reference Alphabet (IRA) (Formerly International Alphabet No. 5 or IA5) – Information technology – 7-bit coded character set for information interchange*.

[ITU-T X.680] Recommendation ITU-T X.680 (2008) | ISO/IEC 8824-1:2008, *Information technology – Abstract Syntax Notation One (ASN.1): Specification of basic notation*.

[ITU-T X.681] Recommendation ITU-T X.681 (2008) | ISO/IEC 8824-2:2008, *Information technology – Abstract Syntax Notation One (ASN.1): Information object specification*.

[ITU-T X.682] Recommendation ITU-T X.682 (2008) | ISO/IEC 8824-3:2008, *Information technology – Abstract Syntax Notation One (ASN.1): Constraint specification*.

[ITU-T X.683] Recommendation ITU-T X.683 (2008) | ISO/IEC 8824-4:2008, *Information technology – Abstract Syntax Notation One (ASN.1): Parameterization of ASN.1 specifications*.

[ITU-T X.690] Recommendation ITU-T X.690 (2008) | ISO/IEC 8825-1:2008, *Information technology – ASN.1 encoding rules: Specification of Basic Encoding Rules (BER), Canonical Encoding Rules (CER) and Distinguished Encoding Rules (DER)*.

[ITU-T X.691] Recommendation ITU-T X.691 (2008) | ISO/IEC 8825-2:2008, *Information technology – ASN.1 encoding rules: Specification of Packed Encoding Rules (PER)*.

[ITU-T X.693] Recommendation ITU-T X.693 (2008) | ISO/IEC 8825-4:2008, *Information technology – ASN.1 encoding rules: XML Encoding Rules (XER)*.

[ITU-T Z.100] Recommendation ITU-T Z.100 (2011), *Specification and Description Language 2008 – Overview of SDL-2010*.

[ITU-T Z.101] Recommendation ITU-T Z.101 (2011), *Specification and Description Language – Basic SDL-2010*.

[ITU-T Z.102] Recommendation ITU-T Z.102 (2011), *Specification and Description Language – Comprehensive SDL-2010*.

[ITU-T Z.103] Recommendation ITU-T Z.103 (2011), *Specification and Description Language – Shorthand notation and annotation in SDL-2010*.

[ITU-T Z.105] Recommendation ITU-T Z.105 (2011), *Specification and Description Language – SDL-2010 combined with ASN.1 modules*.

[ITU-T Z.106] Recommendation ITU-T Z.106 (2011), *Specification and Description Language – Common interchange format for SDL-2010*.

[ITU-T Z.111] Recommendation ITU-T Z.111 (2008), *Notations and guidelines for the definition of ITU-T languages*.

[ISO/IEC 10646] ISO/IEC 10646:2012, *Information technology – Universal Coded Character Set (UCS)*.

3 Definitions

3.1 Terms defined elsewhere

This Recommendation uses the following terms defined elsewhere: the definitions of [ITU-T Z.100] apply.

3.2 Terms defined in this Recommendation

3.2.1 decode (process): The process to construct an SDL-2010 data value from a text string or bit-pattern that is assumed to be an encoding of the SDL-2010 data value using the same encoding rules as those used in the decode.

3.2.2 decoding: The result of a decode process.

3.2.3 encode (process): The process of producing an encoding.

3.2.4 encoding: The text string or bit-pattern resulting from the application of a set of encoding rules to an SDL-2010 data value.

3.2.5 set of encoding rules: One of the sets of (ASN.1) encoding rules (defined in [ITU-T X.690], [ITU-T X.691] and [ITU-T X.693]) or the set of text encoding rules defined in clause 10.7.1 and Annex B of this Recommendation, or an implementation-dependant or application-dependant set of encoding rules defined by a procedure invoked according to this Recommendation (see clause 10.7).

4 Abbreviations and acronyms

This Recommendation uses the following abbreviations and acronyms: the abbreviations defined in [ITU-T Z.100] apply. In addition the following abbreviation applies:

ASN.1 Abstract Syntax Notation One (as defined in [ITU-T X.680], [ITU-T X.681], [ITU-T X.682] and [ITU-T X.683]).

5 Conventions

The conventions defined in [ITU-T Z.100] apply, which includes the conventions defined in [ITU-T Z.111].

Where an abstract or concrete syntax rule is defined in this Recommendation with the same name as a rule in [ITU-T Z.101], [ITU-T Z.102] or [ITU-T Z.103], the rule given here replaces the rule in [ITU-T Z.101], [ITU-T Z.102] or [ITU-T Z.103]. Any *Abstract grammar* or *Concrete grammar* conditions, *Semantics* and *Model* defined on a named rule in [ITU-T Z.101], [ITU-T Z.102] or [ITU-T Z.103] apply to the redefined rule, unless specifically defined otherwise in this Recommendation. Any contradiction between [ITU-T Z.101] or [ITU-T Z.102] or [ITU-T Z.103] and this Recommendation is an error in the definition of SDL-2010 that needs to be resolved by further study.

Some numbered clauses are included so that this Recommendation has the general structure and numbering of [ITU-T Z.101], and in the text the feature is as described in [ITU-T Z.101]. Any extensions described in [ITU-T Z.102] or [ITU-T Z.103] are also allowed.

6 General rules

The general rules of [ITU-T Z.101], [ITU-T Z.102] and [ITU-T Z.103] apply.

However, <operation definition> is an alternative textual concrete syntax to <operation diagram> and therefore defines a scope unit.

7 Organization of Specification and Description Language specifications

7.1 Framework

The framework of specifications is as defined in [ITU-T Z.101].

7.2 Package

Concrete grammar

<package use clause> ::=
[<data binding>]
use <package identifier> [/ <definition selection list>] <end>

<data binding> ::=
default <package identifier>

The <package use clause> of Basic SDL-2010 is extended to include an optional <data binding>.

Each diagram that uses the native SDL-2010 defined concrete syntax for data and actions implicitly or explicitly includes **`default`** `Predefined` in its <package use clause>. If the language <data binding> implicitly or explicitly includes a package defined in Annex C for language binding, each rule identified in Annex C as a concrete syntax variation replaces the rule of the same name wherever it is used, and the rule represents the abstract grammar as defined in Annex C. Each diagram that uses a particular syntax linked to a package name by a language <data binding>, implicitly or explicitly includes a <data binding> for that package name in its <package use clause>. A referenced diagram that is logically included in another diagram implicitly includes the <data binding> of the enclosing diagram, unless a different <data binding> is given. Therefore it is only usually necessary to give the <data binding> for the package for the system type or for the <system specification>.

If the language <data binding> is omitted for the diagram and enclosing diagrams, the implicit language <data binding> is **`default`** `Predefined` if not otherwise specified. It is allowed to specify the implicit <data binding> by means such as analysis directives or tool features not defined by SDL-2010, but the concrete grammar including how it represents the SDL-2010 abstract grammar shall be one of the concrete grammars defined in Annex C.

Model

When an ASN.1 module is used as a package, and a <definition selection> in the <package use clause> referring to the ASN.1 has the <selected entity kind> **interface** and the <name> is the name of a CHOICE data type, there is an implied interface. This implied interface has the same name as the CHOICE and defines signals equivalent to the CHOICE alternatives.

7.3 Referenced definition

The concrete syntax is extended to include operations that are defined textually.

Concrete grammar

<definition> ::=
<macro definition>
| <procedure definition>
| <operation definition>

The <definition> of [ITU-T Z.103] is extended to include <operation definition>. Each <operation definition> shall have corresponding <operation reference> in the associated <package diagram> or <system specification>.

8 Structural concepts

The structural concepts are as defined in [ITU-T Z.101] with the addition of the optional identification of a set of encoding rules for gates and channels. Otherwise, the structural concepts are as defined in [ITU-T Z.101].

8.1 Types, instances and gates

This Recommendation adds the optional identification of a set of encoding rules to the definition of gates.

8.1.1 Structural type definitions

The structural type definitions are as defined in [ITU-T Z.101].

8.1.2 Type expressions

Type expressions are as defined in [ITU-T Z.101].

8.1.3 Abstract types

Abstract types are as defined in [ITU-T Z.102].

8.1.4 Gates with encoding rules

A gate is allowed to have a set of encoding rules. An output of signal from an agent via the gate is encoded as specified by the set of encoding rules. Information received via the gate is decoded according to the set of encoding rules. If no specific set of encoding rules is given, the encoding is not defined by the SDL-2010 specification.

A gate definition defines an anonymous choice data type, which is not described in Basic SDL-2010.

Abstract grammar

Gate-definition :: *Gate-name*
[*Encoding-rules*]
*In-signal-identifier-**set***
*Out-signal-identifier-**set***

Basic SDL-2010 is extended to allow *Gate-definition* to have *Encoding-rules*.

If an external channel with a set of *Encoding-rules* is connected to the gate of an agent or composite state, the *Encoding-rules* of the *Gate-definition* for the gate shall be the same as the *Encoding-rules* for the channel. There shall be, at most, one such channel connected to the gate and the signals conveyed in a direction for the gate shall be the same as the signals conveyed in the corresponding direction in the channel.

If an internal channel of an agent or agent type is connected to the gate of the agent or agent type that has a *Gate-definition* with a set of *Encoding-rules*, the *Encoding-rules* of the channel shall be the same. The signals conveyed in a direction for the gate shall be the same as the signals conveyed in the corresponding direction in the channel. More than one such internal channel is allowed.

Concrete grammar

<gate definition> ::=
{ { <gate symbol 1> | <inherited gate symbol 1> }
is associated with
{ <gate> [<encoding rules>] [<signal list area>] } ***set***
| { <gate symbol 2> | <inherited gate symbol 2> }
is associated with
{ <gate> [<encoding rules>] [<signal list area> <signal list area>] } ***set***
} [***is connected to*** <endpoint constraint>]

Shorthand SDL-2010 <gate definition> is extended to include <encoding rules>.

A specification of <encoding rules> that is associated with an <inherited gate symbol 1> or <inherited gate symbol 2> shall specify the same set of *Encoding-rules* as the *Encoding-rules* of the corresponding gate definition in the supertype if that gate has *Encoding-rules*. If there is no set of <encoding rules> that is associated with an <inherited gate symbol 1> or <inherited gate symbol 2>,

and there is a set of *Encoding-rules* for the gate in the supertype, the inherited gate has this set of *Encoding-rules*. If there is no set of *Encoding-rules* for the gate in the supertype, the presence and value of the inherited gate *Encoding-rules* is determined by the presence and value of the <encoding rules>.

NOTE – To keep the language grammar simple only one <encoding rules> is allowed for a bi-directional channel (<gate symbol 2> or <inherited gate symbol 2>); therefore, if a different encoding is required in each direction, the communication has to be specified by two gates that each carry signals in one direction and each connected to a different channel.

<interface gate definition> ::=
{ <gate symbol 1> | <inherited gate symbol 1> }
is associated with { <interface identifier> [<encoding rules>] }
[***is connected to*** <endpoint constraint>]

Shorthand SDL-2010 <interface gate definition> is extended to include <encoding rules>.

<as gate> ::=
as gate <gate identifier>

A <gate definition> defines a *Data-type-definition* of a choice data type with a unique anonymous *Sort* name in the context that the *Gate-definition* is visible. A <basic sort> that is <as gate> where <gate identifier> identifies the *Gate-definition* represents this *Data-type-definition.* The *Data-type-definition* is equivalent to defining a choice data type with a <choice of sort> for each distinct *Signal-definition* identified by the *In-signal-identifier-**set*** and *Out-signal-identifier-**set***. Each <choice of sort> has as its <field sort> the data type denoted by <as signal> for the identified signal definition (the `NULL` sort for a signal without parameters; otherwise a structure data type with an anonymous unique name – for an identified signal `signal_id` with parameters the <choice of sort> has a <field sort> with identity as denoted by **`as signal`** `signal_id)`. If an identified *Signal-definition* has a name distinct from any other identified distinct *Signal-definition,* the <choice of sort> has the same <field name> as the name of the corresponding identified signal. If an identified *Signal-definition* has the same name as another identified *Signal-definition*, the <choice of sort> has <field name> with the same name as the anonymous unique name for the <as signal> structure data type.

Semantics

The set of *Encoding-rules* of a *Gate-definition* of an agent type or composite state type corresponds to the set of *Encoding-rules* of the channel in the enclosing scope in (the set of) instance specifications.

Model

The set of *Encoding-rules* of a *Gate-definition* of the implied agent type of an <agent diagram> (or the implied composite state type of a <composite state diagram>) is the same as the set of *Encoding-rules* of the external channel (explicit or implied) from which the *Gate-definition* is derived. There is no *Encoding-rules* item in this *Gate-definition* if there is none in the external channel.

The set of *Encoding-rules* of an implied *Gate-definition* of a system *Agent-type-definition* (whether defined by a <system type diagram> or implied from a <system diagram>) is the same as the set of *Encoding-rules* of the internal channel (explicit or implied) from which the *Gate-definition* is derived. There is no *Encoding-rules* item in this *Gate-definition* if there is none in the internal channel.

If an explicit <gate on diagram> is given for the <system type diagram>, the set of *Encoding-rules* of the corresponding *Gate-definition* of the *Agent-type-definition* is determined by the <encoding rules> of the <gate on diagram> if the <encoding rules> is present. If the <encoding rules> is absent, the set of *Encoding-rules* is determined from the internal channel in the same way as for an implied *Gate-definition*.

8.2 Type references and operation references

Type references are as defined in [ITU-T Z.101].

NOTE – A <gate property area> in a type reference is a <gate definition> or <interface gate definition> and therefore optionally contains an <encoding rules> specification, which has to be consistent with the definition given with the type.

<operation reference> ::=
{ **operator** | **method** } <operation signature> **referenced** <end>

Basic SDL-2010 <operation reference> is extended to allow the operation to be a **method**.

8.3 Context parameters

Context parameters (including gate context parameters) are as defined in [ITU-T Z.102].

NOTE – When a gate in parameterized type is defined by a formal context parameter, the set of encoding rules of the gate in a specialized type used to define instances will be determined by the actual gate definition identified by the actual context parameter.

8.4 Specialization

Specialization is as defined in [ITU-T Z.102] with the additional rules for shorthand notation in [ITU-T Z.103].

9 Agents

Agents are as defined in [ITU-T Z.101] or as in [ITU-T Z.102] or [ITU-T Z.103], if they are being used.

10 Communication

Encoding of data extends the definition of channels and connections.

10.1 Channel

A channel connecting two agents determines the encoding (if any) to be used for communication between the agents. A channel that is connected to the environment of an agent has the encoding defined (if any) for that gate connecting it with the environment.

A channel definition defines an anonymous choice data type, which is not described in Basic SDL-2010.

Abstract grammar

Channel-definition :: *Channel-name*
[*Encoding-rules*]
[**NODELAY**]
*Channel-path-**set***

The *Originating-gate* or *Destination-gate* shall have the same *Encoding-rules* as the *Channel-definition*. If the *Channel-definition* has no *Encoding-rules*, neither the *Originating-gate* nor *Destination-gate* shall have *Encoding-rules*.

NOTE 1 – In the *Concrete grammar* it is allowed to omit the <encoding rules> if the channel is connected to gates with encoding, and the *Encoding-rules* derived from the gates are as described in the *Model* below.

Concrete grammar

<channel definition area> ::=
<channel symbol 1>
is associated with
{ [<channel name> [<encoding rules>]] [<signal list area>] [<signal list area>] } ***set***
is connected to {
{ <agent area> | <state machine area> | <gate on diagram>
| <external channel identifiers> }
{ <agent area> | <state machine area> | <gate on diagram>
| <external channel identifiers> } } ***set***
| <channel symbol 2>
is associated with
{ [<channel name> [<encoding rules>]] [<signal list area>] [<signal list area>] } ***set***
is connected to {
{ <agent area> | <state machine area> | <gate on diagram>
| <external channel identifiers> }
{ <agent area> | <state machine area> | <gate on diagram>
| <external channel identifiers> } } ***set***

Shorthand SDL-2010 <channel definition area> is extended to allow <encoding rules> to be specified for a channel.

NOTE 2 – To keep the language grammar simple only one <encoding rules> is allowed for a bi-directional channel (<channel symbol 2>; therefore if a different encoding is required in each direction, the communication has to be specified by two channels that each carry signals in one direction.

<as channel> ::=
as channel <channel identifier>

A <channel definition> defines a *Data-type-definition* of a choice data type with a unique anonymous *Sort* name in the context that the *Channel-definition* is visible. A <basic sort> that is <as channel> where <channel identifier> identifies the *Channel-definition* represents this *Data-type-definition*. The *Data-type-definition* is equivalent to defining a choice data type with a <choice of sort> for each distinct *Signal-definition* identified by each *Signal-identifier-**set*** of the *Channel-path-**set***. The <field name> and <field sort> for each <choice of sort> is determined in the same way as a <choice of sort> for each distinct *Signal-definition* of the *Data-type-definition* defined by a <gate definition>.

Semantics

The *Encoding-rules* of a *Channel-definition* connected to a set of instance specifications is used in the behaviour of the instances.

Model

If the <encoding rules> item is omitted and the <channel symbol 1> or <channel symbol 2> is connected to a <gate on diagram> with <encoding rules>, the *Channel-definition* has the same *Encoding-rules* as the *Gate-definition* of the <gate on diagram>. If the *Gate-definition* has no *Encoding-rules*, the *Channel-definition* has no *Encoding-rules*.

Implicit channels have no *Encoding-rules* if the gates they are connected to have no *Encoding-rules*; otherwise, the implicit channel has the same *Encoding-rules* as each of the gates.

NOTE 3 – If the channel is connected to the frame of the enclosing diagram and there is no <gate on diagram>, this represents a connection.

10.2 Connection

Model

The *Encoding-rules* of an implicit gate of a connection is the same as the connected external channel. A channel without an <encoding rules> (or an implicit channel) that is connected to a gate that has an *Encoding-rules* has the same *Encoding-rules*.

10.3 Signal

Signal is extended so that a signal definition introduces a structure data type (optionally with field names) and use of a structure data type to define the sorts of the signal.

Concrete grammar

<signal definition> ::=
<signal name>
[<formal context parameters>]
[<virtuality constraint>]
[<specialization>]
[<sort list> | <named fields sort list> | **struct** <sort identifier>]

Comprehensive SDL-2010 <signal definition list> is extended to allow a <named fields sort list>.

The <sort identifier> of **struct** <sort identifier> of a <signal definition> shall identify the sort of a structure data type.

In a <signal definition>, a <sort identifier> (after **struct)** adds an *Aggregation-kind* and *Sort-reference-identifier* for each field of the structure as the *Aggregation-kind* and *Sort-reference-identifier* of a *Signal-parameter* to the end of the *Signal-parameter* list of the *Signal-definition*. The *Aggregation-kind* for a *Signal-parameter* corresponding to a field is the same as the *Result-aggregation* for the operator to obtain the value of that field. The *Sort-reference-identifier* for a *Signal-parameter* corresponding to a field is the same as the *Operation-result* for the operator to obtain the value of that field.

<named fields sort list> ::=
(<aggregation kind> <field name> <sort>
{ **,** <aggregation kind> <field name> <sort>}* **)**

Each <aggregation kind> and <sort> in a <named fields sort list> of the <signal definition> adds the corresponding *Aggregation-kind* and *Sort-reference-identifier* of a *Signal-parameter* to the end of the *Signal-parameter* list of the *Signal-definition.*

<as signal> ::=
as signal <signal identifier>

A <signal definition> that defines a *Signal-definition* with an empty *Signal-parameter* list (a signal definition> without a <sort list> or <named fields sort list> or **struct** <sort identifier>), defines in the same context as the *Signal-definition* a *Syntype-definition* with a unique anonymous *Syntype-name*, `NULL`, as the *Parent-sort-identifier*, and empty *Range-condition* and no *Default-initialization.* In this case, an <as signal> (for the signal definition) when used as a <basic sort> denotes the data type `NULL`.

A <signal definition> that defines a *Signal-definition* with a non-empty *Signal-parameter* list defined by a <sort list> or <named fields sort list>, defines in the same context as the *Signal-definition* a *Data-type-definition* for a structure data type with a unique anonymous *Sort* name. In this case, an <as signal> (for the signal definition) when used as a <basic sort> denotes the *Sort* of the *Data-type-definition*. The *Data-type-definition* is equivalent to defining a structure data type with an **optional** <field> for each <aggregation kind> and <sort> item (in order) of the <sort list> or <named fields sort list>, where the <aggregation kind> and <field sort> is the same as those of the <sort list> or <named fields sort list>. If the <signal definition> has a <sort list>, each field has

a unique anonymous name and therefore has to be identified using a <field number> and the field present operation is not accessible because its name is unknown. If the <signal definition> has a <named fields sort list>, each <field> has the name given by the <field name> of the <named fields sort list> item.

For a <signal definition> that defines a *Signal-definition* with a non-empty *Signal-parameter* list defined by **struct** <sort identifier>, an <as signal> (for the signal definition) when used as a <basic sort> denotes the same *Sort* as the <sort identifier> of **struct** <sort identifier>.

10.4 Signal list area

Signal list area is as defined in [ITU-T Z.101] and as in [ITU-T Z.102] and [ITU-T Z.103] if those are being used.

10.5 Remote procedures

Remote procedures are as defined in [ITU-T Z.102].

10.6 Remote variables

Remote variables are as defined in [ITU-T Z.102].

10.7 Communication path encoding rules, encode and decode

The set of encoding rules specifies which set of encoding rules is used to encode and decode data conveyed by a particular channel or gate.

Abstract grammar

Encoding-rules	::	*Rules-identifier*
Encoding-expression	::	*Signal-identifier* [*Expression*]* *Encoding-path*

Encoding-path	::	{
		Gate-identifier |
		| *Data-type-identifier Rules-identifier*
		}

Decoding-expression	::	*Expression* *Encoding-path*
Rules-identifier	::	*Literal-identifier*

The *Rules-identifier* shall be one of the literal identifiers of the data type `Encoding` (see below in *Semantics*). If the actual rule identified corresponds to an encoding defined in [ITU-T X.690], [ITU-T X.691] and [ITU-T X.693], the set of signals carried by an *Encoding-path* shall correspond to elements of an ASN.1 CHOICE type that is carried by the channel, or the channel shall define a choice type that is equivalent to an ASN.1 CHOICE type. The *Encoding-path* and ASN.1 CHOICE type correspond in one direction if each signal name corresponds to a CHOICE name and for each signal the (single) parameter of the signal is the same as the data type of the corresponding CHOICE.

The data type `Encoding` shall be the Predefined data type `Encoding` or a data type with the name `Encoding` that is a specialization (direct or indirect) of the Predefined data type `Encoding`. The specialization shall only add literal names to the Predefined data type `Encoding` and shall not change any other properties.

If the *Rules-identifier* corresponds (by *Name*) to an `Encoding` literal defined in the **`package`** `Predefined`, built-in procedures implied by the standardized sets of encoding rules are invoked (and other procedures – even if visible with names corresponding as below – are ignored).

If the *Rules-identifier* corresponds to an additional literal added to a specialization of the predefined data type `Encoding`, there shall be a visible procedure with the appropriate signature for each invocation (implicit or explicit) of the *Encoding-expression* or *Decoding-expression*.

The name of the procedure for encode is the name `encode` concatenated with the name of an additional `Encoding` literal (for example, `encodemyprotocol` where the additional `Encoding` literal is `myprotocol`). This procedure shall have one parameter of the implicit choice type for the relevant path for the invocation (see below in *Semantics*). The procedure shall return a `Charstring`, `Bitstring` or `Octetstring`.

The name of the procedure for decode is the name `decode` concatenated with the name of an additional `Encoding` literal (for example, `decodemyprotocol` where the additional `Encoding` literal is `myprotocol`). This procedure shall have one parameter of the same type (`Charstring`, `Bitstring` or `Octetstring`) as the corresponding procedure for encode, and shall return the choice type for the relevant path for the invocation (see below in *Semantics*).

Each encode or decode procedure shall be functional (that is, it shall not contain states and shall not change the value of any variable external to the procedure when it is interpreted).

The length of the optional *Expression* list of an *Encoding-expression* shall be the same as the number of *Sort-reference-identifier*s in the *Signal-definition* denoted by the *Signal-identifier*.

Each *Expression* in an *Encoding-expression* shall be sort compatible to the corresponding (by position) *Sort-reference-identifier* in the *Signal-definition* denoted by the *Signal-identifier*.

For a *Gate-identifier* of an *Encoding-path* of an agent's *Encoding-expression,* the gate shall be a gate of the agent or reachable via a channel with the *Signal-identifier* of the *Encoding-expression* from the agent, and the *Out-signal-identifier-**set*** of the gate shall include the *Signal-identifier*.

The *Data-type-identifier* of an *Encoding-path* shall identify an *Interface-definition.*

For an *Encoding-path* of an *Encoding-expression* with a *Data-type-identifier* that identifies an *Interface-definition,* the *Signal-identifier* shall identify one of the signals defined or inherited by the interface identified.

The *Expression* in a *Decoding-expression* shall be compatible with the sort generated by an *Encoding-expression* using the same *Encoding-path* in a context where the context of the *Decoding-expression* is reachable via the *Encoding-path.*

Concrete grammar

<encoding rules> ::=
encode <rules identifier>

<rules identifier> ::=
<literal>

<encoding expression> ::=
encode { <signal identifier> [(<actual parameters>)] | <expression> }
[<encoding path>]

<encoding path> ::=
as { [**channel**] <channel identifier>
| [**gate**] <gate identifier>
| [**interface**] <interface identifier> **with** <rules identifier> }

<decoding expression> ::=
decode <expression> [<encoding path>]

The set of encoding rules of the path identified by the <channel identifier> or <gate identifier> (or in the case of an <interface identifier>, by the <rules identifier>) is used. A <channel identifier> in an <encoding path> is a shorthand that represents a (possibly anonymous) gate.

If an <encoding expression> contains an <expression> (rather than a signal), the actual signal is derived from the <expression> as described in the model below. The sort of the <expression> shall be the sort of the choice data type corresponding to the set of signals for the <encoding path>: normally the choice data type defined for the channel, gate or interface.

The <encoding path> shall only be omitted from <encoding expression> if there is exactly one path (channel or gate) with encoding for output of the signal in the context of the <encoding expression> and, in this case, the encoding for that path is used.

The <encoding path> shall only be omitted from <decoding expression> if there is exactly one path (channel or gate) with encoding for input in the context of the decoding expression and, in this case, the encoding for that path is used.

The <interface identifier> and <rules identifier> of an <encoding path> represents the *Data-type-identifier* and *Rules-identifier* of an *Encoding-path*.

Semantics

The *Encoding-rules* item determines the set of rules used to change the implementation-dependent encoding for internal data:

– either to a standardized implementation-independent encoding (for one of the data type `Encoding` literals `text` to `EXER`, or `PER` as a synonym for `APER`),

– or to a defined implementation or application encoding (for a literal added to a specialization of the data type `Encoding`).

Encode is invoked whenever a signal is output via a path with a set of encoding rules specified and the corresponding encode procedure is called. When a signal is input from such a path, decode of the data into the internal encoding is invoked by calling the corresponding decode procedure. This encode and decode therefore has no impact on the semantics of SDL, as defined in [ITU-T Z.101], but requires specific encoding of signals for the specified paths and, therefore, enables different parts of the system to be implemented separately.

When an *Encoding-expression* or *Decoding-expression* is interpreted, the appropriate procedure is called.

Encode relative to the set of encoding rules for a path is invoked in an *Encoding-expression* to produce a `Charstring`, `Bitstring` or `Octetstring` depending on the context and the set of encoding rules used. The `Charstring`, `Bitstring` or `Octetstring` produced by encode is decoded by a *Decoding-expression* using the same set of encoding rules for the same path. The set of encoding rules of the encoding path identified by the *Gate-identifier* or *Rules-identifier* (after an interface *Data-type-identifier*) is used.

For an *Encoding-expression*, the data is encoded as if it were going to be output on that path. The result is a data type corresponding to the set of encoding rules for the path.

For a *Decoding-expression*, the data is decoded as if it had been received on the specified path. The result is an expression corresponding to the implicit data type for input from that channel in the decode context as defined below. If decoding fails, the InvalidReference exception is raised.

For a channel, gate or interface that conveys the signals `signal1`, `signal2` and `signal3`, for a choice data type is with a unique, anonymous name defined that corresponds to the SDL:

```
value type UniqueAnon /* representing the unique anononymous name */
{ choice
        signal1 as signal signal1;
        signal2 as signal signal2;
        signal3 NULL;
}
```

where

```
as signal signal1
```

represents the structure data type with a unique, anonymous name

```
value type signal1UniqueAnon
{ struct
   Ua11 Sort11 optional;
   Ua12 Sort12 optional;
   Ua13 Sort13 optional;
   /* ... and so on for each parameter of signal1 */
 } ;
```

as a consequence of the signal definition

```
signal signal1 (Sort11, Sort12, Sort13 /* ... and so on */);
```

and `Ua11`, `Ua12` and `Ua13` are unique, anonymous field names for the structure data type, and similarly

```
as signal signal2
```

represents the structure data type with a unique, anonymous name

```
value type signal2UniqueAnon
{ struct
   Ua21 Sort21 optional;
   Ua22 Sort22 optional;
 } ;
```

as a consequence of the signal definition

```
signal signal2 (Sort21, Sort22);
```

and `signal3` is defined by

```
signal signal3;
```

The anonymous identifier of the choice data type associated with a path is denoted by the **as channel** <channel identifier> or **as gate** <gate identifier> for the path, so that a legal variable declaration is:

```
dcl message as gate user_input;
```

where `user_input` is the name of a channel or gate with encoding and a valid assignment is:

```
message := decode encoded_value as user_input;
```

Similarly, the anonymous identifier of the data type associated with an interface is denoted by the **as interface** <interface identifier>, so that a legal variable declaration is:

```
dcl if_message as interface user_interface;
```

where `user_interface` is the name of an interface and a valid assignment (assuming the information was encoded using PER) is:

```
if_message := decode encoded_value as user_interface with PER;
```

The enumerated data type for the standardized set of encoding rules is defined in the package `Predefined` (see clause 14).

Model

If an <encoding expression> contains an <expression>, the choice present is determined from the expression, and the signal with the same name as the choice is output with the values of the signal parameters given by the expression.

If an <encoding path> contains a <channel identifier>, this is transformed to the (possibly anonymous) <gate identifier> for the gate of the identified channel nearest to the agent containing the <encoding path>.

10.7.1 The set of text encoding rules

The set of text encoding rules is provided so that it is possible to convey information on communication paths by means of text strings between elements in the system, and between the system and the environment. The encoding of characters is not defined. Though the text strings are, in general, readable by humans, that is not the purpose of the text encoding rules.

Semantics

If the set of encoding rules is specified as `text` the result of an encoding is a `Charstring`.

The actual string is determined as follows:

– LEFT CURLY BRACKET and RIGHT CURLY BRACKET { } delimit the values of data types and show where the values start and stop, except where they occur within the encoding of a `Charstring` in which case they represent the actual <left curly bracket> or <right curly bracket> characters of [ITU-T Z.101];

– COMMA characters are used to delimit elements within a list (for example, in a **struct** encoding);

– Otherwise the actual string of characters is determined for each data type as defined in Annex B and illustrated as the 'Generated `Charstring'` in the examples in Annex B.

A complete signal is encoded as a list of values and is treated as a **choice** encoding of the data type for the path with encoding.

The full details of text encoding are given in Annex B.

10.7.2 The sets of encoding rules standardized in the ITU-T X.69x series

The use of one of the names `BER`, `CER`, `DER`, `PER`, `APER`, `UPER`, `CAPER`, `CUPER`, `BXER`, `CXER` or `EXER` (corresponding to [ITU-T X.690], [ITU-T X.691] and [ITU-T X.693] encoding rules – see *Semantics* in clause 10.7 above and clause 14.17 below) shall only be specified if the signals carried by the path or included in the interface are defined by an ASN.1 CHOICE data type or are able to be expressed as an ASN.1 CHOICE data type. The values are encoded according to this data type treated as an ASN.1 ABSTRACT-SYNTAX.

NOTE – Whether it is possible to express a set of signals not defined by an ASN.1 CHOICE as an ASN.1 CHOICE is not currently further defined.

11 Behaviour

11.1 Start

Start is as defined in [ITU-T Z.101].

11.2 State

Semantics

If during the consideration of the signals on the input port, the signal being considered does not correspond to any of the signals valid for any of the states of the agent because it is not possible to decode the message received (as in clause 10.7) to a valid signal, the `InvalidReference` exception is raised.

11.3 Input

Input is extended to allow a choice value to be stored for the signal that was input, and is also extended so that it is possible to store signals received on paths with encoding in the encoded form.

If an <encoded input> is given for a path (a channel or gate), the messages received from that path are received and stored in the variable given in their encoded form. It is allowed to specify these signals in other inputs or saves of the same state only if a <via path> is specified. Although the model given below describes the signals as first being decoded and then encoded again, it is expected that real implementations will optimize this to copying the encoded value to the variable given in the <encoded input>.

Concrete grammar

<input list> ::=
<stimulus> [<in choice>] [<priority clause>]
{ **,** <stimulus> [<in choice>] [<priority clause>] }*
| <asterisk input list> [<in choice>] [<priority clause>]
| <encoded input>

The <input list> syntax of Comprehensive SDL-2010 is extended to allow an <encoded input>.

<in choice> ::=
in <choice variable>

The <choice variable> of an <in choice> shall have a choice data type that includes fields where all the signals that are mentioned in the <stimulus> are covered: that is, for each signal (for example, with the identity `signal_id`) there shall be a field with the unique name denoted by an <as signal> for the signal (**`as signal`** `signal_id`) and the data type defined by <as signal> for the signal (**`as signal`** `signal_id`).

NOTE 1 – If the <encoding path> is specified by an <interface identifier> with <rules identifier>, because the <encoded input> is transformed into inputs for the signals of the identified interface, no other input or save of the state is allowed to mention the same signals.

<encoded input> ::=
encode <variable> [<encoding path>]

If the <encoding path> is omitted from <encoded input>, there shall be exactly one path with encoding leading to the context of the input that has a set of encoding rules that produces the sort (`Charstring`, `Octetstring` or `Bitstring`) of the <variable>. Otherwise the <encoding path> shall have a set of encoding rules that produces the sort (`Charstring`, `Octetstring` or `Bitstring`) of the <variable>.

Semantics

If the signal specified in the input is received via a channel that has a set of encoding rules specified, the signal is decoded according to that set of encoding rules.

Model

Before the model for <in choice> is applied, other models for input are applied so that each input contains only one signal. When applying the model for <asterisk input list> the <in choice> is added to each input. For a <stimulus> with an <in choice>, the values conveyed by signal are assigned to the choice variable as detailed below instead of to the variables given in the <stimulus> followed by a transformed transition that is an implied task and then followed by the original transition. The field of the choice variable corresponding to the signal (the choice field) receives the signal: that is, each value conveyed by the signal is assigned to the corresponding field of the choice field. If the signal conveys no values, the implicit task contains an assignment of NULL to the field of the choice variable corresponding to the signal. If the signal conveys values, for each variable in the <stimulus> the task contains an assignment of the corresponding field of the choice field to the variable in the <stimulus>.

If the <encoding path> is omitted from <encoded input>, the encoding path is the unique path for the <encoded input>.

For each of the signals that are receivable from the gate or channel specified in an <encoded input>, there is an implied input with a <via path> for the gate or channel. For each of the signals of an interface identified by an <interface identifier> of an <encoding path> of an <encoded input>, there is an implied input for the signal.

NOTE 2 – If the <encoding path> is specified by an <interface identifier> **with** <rules identifier>, because the <encoded input> is transformed into inputs for the signals of the identified interface, no other input or save of the state is allowed to mention the same signals.

Each implied input for <encoded input> is equivalent to assigning the values conveyed by the signal to implicit local variables of appropriate types followed by transformed transition that is an implied task followed by the original transition for the <encoded input>. The implied task assigns the <variable> the value of an encoding expression for the signal and values received from the implicit variables using the set of encoding rules for the path of the <encoded input>.

11.4 Priority input

Priority input is as defined in [ITU-T Z.102] and is not allowed to contain an <encoded input>.

11.5 Continuous signal

Continuous signal is as defined in [ITU-T Z.102].

11.6 Enabling condition

Semantics

If the signal specified is received via a channel that has a set of encoding rules specified, the signal is decoded according to that set of encoding rules.

11.7 Save

Save is as defined in [ITU-T Z.101].

11.8 Implicit transition

Implicit transition is as defined in [ITU-T Z.101].

11.9 Spontaneous transition

Spontaneous transition is as defined in [ITU-T Z.101].

11.10 Label

Label is as defined in [ITU-T Z.101].

11.11 State machine and composite state

State machine and composite state are as defined in [ITU-T Z.101].

11.12 Transition

Transition is as defined in [ITU-T Z.101].

11.13 Action

11.13.1 Task

Task is as defined in [ITU-T Z.101].

11.13.2 Create

Create is as defined in [ITU-T Z.101].

11.13.3 Procedure call

Procedure call is as defined in [ITU-T Z.101].

11.13.4 Output

If an <expression output> is given for a path (a channel or gate), the expression is a choice value used to derive the signal to output. The choice sort corresponds to the set of signals for the path.

If an <encoded output> is given for a path (a channel or gate), the expression given is used as the signal to output. In the model given below the expression is decoded to check the signal being sent.

Concrete grammar

<output body item> ::=
    <signal identifier> [<actual parameters>]
  | <expression output>
  | <encoded output>

<expression output> ::=
    <expression>

<encoded output> ::=
    encode <expression>

The sort of an <expression output> shall be the sort of a choice data type, where each choice field name corresponds to an outgoing signal carried from the local agent by one of the gates of the agent. A choice field of the choice data type corresponds if its field name is the same as a signal name that unambiguously identifies an outgoing signal, or its field name is the name for an <as signal> with a <signal identifier> that identifies an outgoing signal.

When an <encoded output> is used, there shall be exactly one <via path> in the <communication constraints> and this <via path> shall specify a gate or channel that has a set of encoding rules that corresponds to the sort (`Charstring`, `Octetstring` or `Bitstring`) of the <expression> of the <encoded output>.

Semantics

If a signal is output via a path that has a set of encoding rules specified, the signal is encoded according to that set of encoding rules.

Model

If the <output body item> is an <expression output>, the choice present is determined from the expression and the signal with the same name as the choice is output with the values of the signal parameters given by the expression.

If the <output body item> is an <encoded output>, the signal that is output is the signal obtained from decoding the contents of the expression according to the decoding rule for receiving signals sent on the specified path. The signal selected is the signal with the same name as the choice present in the decoding. If the signal is not valid for the path, or decoding fails for any reason (for example, if the string of the expression does not validly correspond to one of the signals for the path), the OutOfRange exception is raised and no signal is sent; otherwise the value of the decoding is output as the signal.

NOTE – Because the expression is already encoded as a string for the path, it is possible to send the string value with no further conversion. It is not therefore expected that an implementation would normally decode and re-encode an <encoded output>.

11.13.5 Decision

Decision is as defined in [ITU-T Z.101].

11.14 Statement list

Statement list is as defined in [ITU-T Z.101].

11.15 Timer

Timer is as defined in [ITU-T Z.101].

12 Data

The concept of data in SDL-2010 is defined in this clause. This includes the data terminology, the concepts to define new data types and the predefined data.

Data in SDL-2010 is principally concerned with data types. A data type defines a set of elements or data items, referred to as sort, and a set of operations that are allowed to be applied to these data items. The sorts and operations define the properties of the data type. These properties are defined by data type definitions.

A data type consists of a set, which is the *sort* of the data type, and one or more *operations*. As an example, consider the predefined data type `Boolean`. The sort `Boolean` of the data type Boolean consists of the elements `true` and `false`. Among the operations of the data type `Boolean` are `"="` (equal), `"/="` (not equal), `"not"`, `"and"`, `"or"`, `"xor"`, and `"=>"` (implies). As a further example, consider the predefined data type Natural. It has the sort Natural consisting of the elements `0`, `1`, `2`, etc., and the operations `"="`, `"/="`, `"+"`, `"-"`, `"*"`, `"/"`, `"mod"`, `"rem"`, `"<"`, `">"`, `"<="`, `">="`, and `power`.

The Specification and Description Language provides several predefined data types, which are familiar in both their behaviour and syntax. The predefined data types are described in clause 13.

Variables are objects that are associated with an element of a sort by assignment. When the variable is accessed, the associated data item is returned.

The elements of the sort of a data type are either *values*, or *pids*, which are references to agents. The sort of a data type is defined in the following ways:

a) Explicitly enumerating the elements of the sort.

b) Forming the Cartesian product of sorts S_1, S_2, … S_n; the sort is equal to the set that consists of all tuples formed by taking the first element from sort S_1, taking the second element from sort S_2, … , and finally, taking the last element from sort S_n.

c) For the sorts of pids, by defining an interface (see clause 12.1.2).

d) As one of several sorts that are predefined and form the basis of the predefined data types described in Annex D. The predefined sort `Pid` is described in clause 12.1.5.

The elements of a value sort are values. The elements of a pid sort are pids.

Operations are defined from and to elements of sorts. For instance, the application of the operation for summation (`"+"`) from and to elements of the `Integer` sort is valid, whereas summation of elements of the `Boolean` sort is not.

Each data item belongs to exactly one sort. That is, sorts never have data items in common.

For most sorts there are literal forms to denote elements of the sort: for example, for Integers, "`2`" is used rather than "`1 + 1`". It is allowed that more than one literal denote the same data item; for example, `12` and `012` are used to denote the same Integer data item. It is also allowed that the same literal denotation is used for more than one sort; for example, `'A'` is both a Character and a Character String of length one. Some sorts have no literal forms to denote the elements of the sort; for example, the sorts that are formed as the Cartesian product of other sorts. In that case, the elements of these sorts are denoted by operations that construct the data item from elements of the component sort(s).

An expression denotes a data item. If an expression does not contain a variable or an imperative expression, e.g., if it is a literal of a given sort, each occurrence of the expression will always denote the same data item. These "passive" expressions correspond to a functional use of the language.

An expression that contains variables or imperative expressions is interpreted as having (usually) different results during the interpretation of a Specification and Description Language system, depending on the data item associated with the variables. The active use of data includes assignment

to variables, use of variables, and initialization of variables. The difference between active and passive expressions is that the result of a passive expression is independent of when it is interpreted, whereas an active expression has (usually) different results depending on the current values or pids associated with variables or the current system state.

12.1 Data definitions

Data definitions are used to define data types. The basic mechanisms to define data are data type definitions (see clause 12.1.1) and interfaces (see clause 12.1.2). Definition of additional operations is described in clause 12.1.4. The definition of the sort of the data type as well as operations implied for the sort are given by data type constructors (see 12.1.6). Clause 12.1.7 shows how to define the behaviour of the operations of a data type. Clause 12.1.8 details a variety of features related to data definitions including syntypes, synonyms and visibility restriction. Specialization (see clause 12.1.9) allows the definition of a data type to be based on another data type, referred to as its supertype.

Concrete grammar

```
<sort> ::=
                        <basic sort> [ ( <range condition> ) ]
                |       <anchored sort>
                |       <pid sort>
                |       <inline data type definition>
                |       <inline syntype definition>

<inline data type definition> ::=
                        value [<data type specialization>]
                        [ [ <comment body> ] <left curly bracket> <data type definition body>
                        <right curly bracket> ]

<inline syntype definition> ::=
                        syntype <basic sort>
                        [ [ <comment body> ] <left curly bracket>
                                { <default initialization> [ [<end>] <constraint> ] | <constraint> } <end>
                          <right curly bracket> ]

<basic sort> ::=
                        <datatype type expression>
                |       <as signal>
                |       <as interface>
                |       <as channel>
                |       <as gate>
                |       <syntype>
```

An <as signal> represents the sort of the structure data type introduced by the identified signal definition.

An <as interface> represents the sort of the choice data type introduced by the identified interface definition.

An <as channel> represents the sort of the choice data type introduced by the identified channel definition.

An <as gate> represents the sort of the choice data type introduced by the identified gate definition.

```
<anchored sort> ::=
                        this [<basic sort>]
                |       parent [<basic sort>]

<pid sort> ::=
                        <sort identifier>
```

An <anchored sort> with <basic sort> is only allowed within the definition of <basic sort>.

An <anchored sort> is legal concrete syntax only if it occurs within a <data type definition>. The <basic sort> in the <anchored sort> shall name the <sort> introduced by the <data type definition>.

The meaning of an <anchored sort> is given in *Model* of clause 12.1.9.

Model

An <anchored sort> without a <basic sort> is a shorthand for specifying a <basic sort> with the name of the data type definition or syntype definition in the context of which the <anchored sort> occurs.

A <sort> that is a <basic sort> with a <range condition> is derived concrete syntax for a <syntype> of an implied <syntype definition> having an anonymous name. This anonymously named <syntype definition> is defined with its elements restricted by the <range condition>, if the <basic sort> has been constructed using the literal data type constructor; otherwise, the <range condition> is part of a <size constraint>.

An <inline data type definition> is derived concrete syntax for a <basic sort> of an implied <data type definition> having an anonymous name. This anonymously named <data type definition> is derived from the <inline data type definition> by inserting **type** and the anonymous name after **value** in the <inline data type definition>. Each <inline data type definition> defines a different implied <data type definition>.

An <inline syntype definition> is derived concrete syntax for a <basic sort> of an implied <syntype definition> having an anonymous name. This anonymously named <syntype definition> is derived from the <inline syntype definition> by inserting the anonymous name and <equals sign> after syntype in the <inline syntype definition>.

12.1.1 Data type definition

Although SDL-2010 does not include generators for data types, parameterized data types in **package** Predefined of SDL-2010 replace the generators such as Array that were included in the **package** Predefined in SDL-92. The legacy data type definition below includes legacy generators syntax so that these parameterized data types are allowed to be used with SDL-92 syntax.

Concrete grammar

```
<entity in data type> ::=
                        <data type definition>
                    |   <syntype definition>
                    |   <synonym definition>
                    |   <legacy data type definition>
                    |   <legacy syntype definition>
```

The alternatives <synonym definition>, <legacy syntype definition> and <legacy data type definition> extend <data definition> compared with Basic SDL-2010. The legacy forms are alternative syntax to represent a *Data-type-definition* or *Syntype-definition*. The alternative <synonym definition> is explained in clause 12.1.8.3.

```
<data type definition> ::=
                        {<package use clause>}*
                        <type preamble> <data type heading> [<data type specialization>]
                        {           <end>
                            |       [ <comment body> ] <left curly bracket> <data type definition body>
                                    <right curly bracket> }
```

The <data type specialization> is added to <data type definition> so that it is possible to define a data type inheriting and specializing another data type (see clause 12.1.9). If the <type preamble> is **virtual** or **redefined**, a virtual data type is defined.

```
<data type heading> ::=
                        value type <data type name>
                            [ <formal context parameters> ] [<virtuality constraint>]
```

The <data type heading> is extended compared with Basic SDL-2010 to include <formal context parameters>, to allow use of a data type with context parameters.

```
<legacy data type definition> ::=
                        newtype <sort name>
                            [<formal context parameters>]
                            [<data type specialization>
                            |<legacy generators>
                            |<structure definition>]
                            [<literal list> ]
                            [<legacy operator signatures> ]
                            {       <legacy operator definition>
                                  | <legacy operator reference>
                                  | <legacy external operator definition> }*
                            [ <default initialization> [ <end> ] ]
                            [ constants <range condition> ]
                        endnewtype [ <sort name> ]

<legacy generators> ::=
                        <sort identifier> (<legacy generator actual> { , <legacy generator actual> }* )
```

The <sort identifier> of <legacy generators> should identify one of the parameterized data types in **package** `Predefined`. The <legacy generator actual> should be an appropriate actual parameter for the parameterized data types.

```
<legacy generator actual> ::=
                        <sort>
                    |   <literal signature>
                    |   <operator name>
                    |   <constant expression>
```

To be consistent with SDL-92, the <data type specialization> in a <legacy data type definition> should contain a <legacy data inheritance>.

To be consistent with SDL-92, the <structure definition> in a <legacy data type definition> should not contain <visibility>, **optional** or <field default initialization>.

To be consistent with SDL-92, the <literal list> in a <legacy data type definition> should not contain <visibility> or <named number>.

If a <legacy data type definition> definition contains a <range condition>, it represents the definition of syntype and an anonymous parent data type.

The definition of a legacy operator shall be defined either by the <legacy operator definition> or the operator referenced by a <legacy operator reference> or <legacy external operator definition>.

A <formal context parameter> of <formal context parameters> of a <data type heading> or <legacy data type definition> shall be either a <sort context parameter> or a <synonym context parameter list>.

Basic SDL-2010 is extended so that in an <operation signature> of <operation signatures> there shall be one and only one corresponding definition (<operation reference> or <external operation definition>) in the <operation definitions> of the <operations>. An <external operation definition> is an additional alternative to <operation reference> compared with Basic SDL-2010.

Semantics

A *Value-data-type-definition* describes the elements of the sort and operations induced by the way the sort is constructed, and operations each characterized by an operation signature and the corresponding *Procedure-definition*. The operations are the set of operations that are allowed for the elements of a sort (see clause 12.1.4). The data type identified by the *Data-type-identifier* of the

data type definition is the inherited data type. That is, the defined data type is a specialization (see clause 12.1.3) of the identified data type.

12.1.2 Interface definition

An interface is used to define a set of signals, remote procedures and remote variables provided by or used on a gate of an agent. The defining context of entities defined in the interface is the scope unit of the interface, and the entities defined are visible where the interface is visible. An interface is also allowed to refer to signals, remote procedures, or remote variables defined outside the interface through the <interface use list>.

An interface is used in a signal list to denote that the signals, remote procedure calls and remote variables of the interface definition are included in the signal list.

An interface definition defines an anonymous choice data type, which is not described in Basic SDL-2010.

Concrete grammar

```
<interface definition> ::=
                        {<package use clause>}*
                        [<virtuality>] <interface heading>
                        {       <interface specialization>  <end>
                            |   [<interface specialization>] [ <comment body> ]
                                <left curly bracket>
                                    <entity in interface>*
                                <right curly bracket>
                        }
                    |   <signal list definition>
```

Basic SDL-2010 is extended to allow <virtuality> to be specified for an <interface definition>, and <interface specialization> to allow the interface definition to be a specialization of another interface definition. Basic SDL-2010 is also extended with <end> for an empty set of entities defined in the interface, in which case there has to be an <interface specialization>; otherwise the interface has no identified signals, remote procedures or remote variables.

```
<interface heading> ::=
                        interface <interface name>
                            [<formal context parameters>] [<virtuality constraint>]
```

The <formal context parameters> shall only contain parameters specified by a <signal context parameter list>, or a <remote procedure context parameter>, or a <remotevariable context parameter list> or a <sort context parameter>.

```
<entity in interface> ::=
                        <signal definition list>
                    |   <interface use list>
                    |   <interface variable definition>
                    |   <interface procedure definition>
```

The Basic SDL-2010 <entity in interface> is extended to allow <interface variable definition> entities for remote variables and <interface procedure definition> entities for remote procedures.

```
<interface variable definition> ::=
                        dcl <remote variable name> { , <remote variable name>}* <sort> <end>

<interface procedure definition> ::=
                        procedure <remote procedure name> <procedure signature> <end>
```

NOTE – The <remote variable name> or <interface procedure definition> each has to be a <name>, whereas in SDL-2000 each is allowed to be a name or a number (an <integer name> or a <real name>).

The semantics of <virtuality> is defined in clause 8.4.2 of [ITU-T Z.102].

The content of an interface is the set of all signals, remote procedures and remote variables that are defined in an <entity in interface> of the interface or referenced in the <interface use list> or included in the interface by specialization (that is, inheritance or context parameterization).

<as interface> ::=

as interface <interface identifier>

An <interface definition> defines a *Data-type-definition* of a choice data type with a unique anonymous *Sort* name in the context that the *Interface-definition* is visible. A <basic sort> that is <as interface> where <interface identifier> identifies the *Interface-definition* represents this *Data-type-definition*. The *Data-type-definition* is equivalent to defining a choice data type with a <choice of sort> for each distinct *Signal-definition* identified by the *Signal-identifier-**set*** of the *Interface-definition*. The <field name> and <field sort> for each <choice of sort> is determined in the same way as a <choice of sort> for each distinct *Signal-definition* of the *Data-type-definition* defined by a <gate definition>.

Model

The inclusion of an <interface identifier> in a <signal list> means the content of the interface (that is, all signal identifiers, all remote procedure identifiers and remote variable identifiers forming part of the <interface definition>) are included in the <signal list>.

Internally connected gates of an agent (or agent type) are explicit or implicit gates of the agent (or agent type respectively) that are connected via implicit or explicit channels to the gates of either the state machine of the agent (or agent type respectively) or a contained agent. The interface defined by an agent or agent type contains in its <interface specialization> all interfaces given in the incoming signal list associated with internally connected gates. The interface contains in its <interface use list> all signals, remote variables and remote procedures given in the incoming signal list associated with internally connected gates. In addition, the interface for an agent type that inherits another agent type also contains in its <interface specialization> the implicit interface defined by the inherited agent type.

The implicit interface for an agent type that inherits another agent type also contains in its interface specialization the implicit interface defined by the inherited agent type.

If the containing entity is an agent type that inherits another agent type, then the interface of the state machine of the agent type also contains in its interface specialization the implicit interface of the state machine of the inherited agent type.

12.1.3 Operation signature

Basic SDL-2010 is extended to include methods, legacy syntax and operation visibility. A method is an operation that is applied to a variable and is able (but does not have to) modify the value associated with a variable.

Concrete grammar

<operation signatures> ::=

[<operator list>] [<method list>]

The Basic SDL-2010 <operation signatures> is extended to allow a <method list>.

<legacy operator signatures> ::=

operators

<legacy operator signature> { <end> <legacy operator signature> }* [<end>]

A <legacy data type definition> has <legacy operator signatures> instead of <operation signatures> for the concrete syntax of operation signatures.

<legacy operator signature> ::=

<operator name> : <arguments> -> <sort>

A <legacy operator signature> represents an *Operation-signature*.

The <sort> of a <legacy operator signature> represents the *Result* of the *Operation-signature*.

<method list> ::=
methods <operation signature> { <end> <operation signature> }* <end>

<operation signature> ::=
<operation preamble>
<operation name>
[<arguments>] [<result>]

The Basic SDL-2010 <operation signature> is extended to include an <operation preamble> that defines the visibility scope of the operation. Omitting the <result> is allowed for a method, as explained below.

<operation preamble> ::=
[<visibility>]

<result> in <operation signature> shall be omitted only if the <operation signature> occurred in a <method list>.

Model

If <operation signature> is contained in a <method list>, this is derived syntax and is transformed as follows: an <argument> is constructed from the <parameter kind> **in/out**, and the <sort identifier> of the sort being defined by the enclosing <data type definition>. If there are no <arguments> (that is, the original argument list was empty), <arguments> is formed from the constructed <argument> and inserted into the <operation signature>. If there are <arguments>, the constructed <argument> is added to the start of the original list of <argument>s in the <arguments>. If the <result> was omitted, the <result> is the <sort identifier> of the sort being defined by the enclosing <data type definition>.

NOTE – An empty <parameter kind> is allowed and has the same meaning as the <parameter kind> **in** (see Procedure in [ITU-T Z.103]).

12.1.4 Generic data type operations

These are as defined in Basic SDL-2010.

12.1.5 `Pid` and pid sorts

These are as defined in Basic SDL-2010, except specialization is allowed so the output compatibility check (b) in clause 12.1.5 of [ITU-T Z.101] applies.

12.1.6 Data type constructors

12.1.6.1 Literals

Concrete grammar

<literal list> ::=
[<visibility>] **literals** <literal signature> { **,** <literal signature> }* <end>

The Basic SDL-2010 <literal list> is extended to allow restricting visibility of the literals.

Semantics

The meaning of <visibility> in <literal list> is explained in clause 12.1.8.4.

12.1.6.2 Structure data types

The concrete syntax of structure data types is extended to include visibility constraints and a shorthand for several consecutive fields of the same sort.

Concrete grammar

<structure definition> ::=
[<visibility>] **struct** [<field list>] <end>

The Basic SDL-2010 <structure definition> is extended to allow restricting <visibility> of the structure data type.

<fields of sort> ::=
[<visibility>] <field name> { **,** <field name> }* <field sort>

The Basic SDL-2010 <fields of sort> is extended to allow restricting <visibility> of fields of the structure. The Basic SDL-2010 <fields of sort> is also extended to allow two or more adjacent fields of the same sort to be defined more concisely.

The meaning of <visibility> in <fields of sort> and <structure definition> is explained in clause 12.1.8.4.

The generic operations to create structure values are applied as operators.

The generic operations for structure values that have an initial **in/out** parameter (to modify fields, access fields, test for the presence of field and test for the structure value being undefined) are treated as methods.

Model

A <field list> containing a <field> that has a <fields of sort> with <field name> list is derived concrete syntax where this <field> is replaced by a list of <field>s separated by <end>, one for each <field name> in the order of occurrence of each <field name>. Each <field> in the replacement list has a <fields of sort> with same <visibility> and <field sort> as the original <fields of sort>. Each <field> in the replacement list is **`optional`** if the original <field> was **`optional`**, or has the <field default initialization> of the original <field> if there was one.

After transforming any <field name> list as above, a <structure definition> introducing a sort named `S` implies for each field `fn` the following in the <method list> of `S`:

```
fnExtract -> fs;
fnModify ( fs ) -> S;
fnPresent -> <<package Predefined>>Boolean;
```

These are transformed according to the model for methods in clause 12.1.4 to:

```
fnExtract ( in/out S ) -> fs;
fnModify  ( in/out S, fs ) -> S;
fnPresent ( in/out S ) -> <<package Predefined>>Boolean;
```

12.1.6.3 Choice data types

The concrete syntax of choice data types is extended to include visibility constraints.

Concrete grammar

<choice definition> ::=
[<visibility>] **choice** [<choice list>] <end>

The Basic SDL-2010 <choice definition> is extended to allow restricting visibility of the choice data type.

<choice of sort> ::=
[<visibility>] <field name> <field sort>

The Basic SDL-2010 <choice of sort> is extended to allow restricting visibility of the choice field.

The meaning of <visibility> in <choice of sort> and <choice definition> is explained in clause 12.1.8.4.

The generic operations to create choice values are applied as operators.

The generic operations for choice values that have an initial **in/out** parameter (to modify fields, access fields, test for the presence of a particular field, test for which field is present in the choice value, and test for a choice value being undefined) are treated as methods.

Model

A <choice definition> introducing a sort named `C` implies for each field `fn` the following in the <operator list> of `C`:

```
fn ( fs ) -> C;
```

A <choice definition> introducing a sort named `C` implies for each field `fn` the following in the <method list> of `C`:

```
fnExtract -> fs;
fnModify ( fs ) -> C;
fnPresent -> <<package Predefined>>Boolean;
```

These are transformed according to the model for methods in clause 12.1.4 to:

```
fnExtract ( in/out C ) -> fs;
fnModify  ( in/out C, fs ) -> C;
fnPresent ( in/out C ) -> <<package Predefined>>Boolean;
```

12.1.7 Behaviour of operations

The syntax for the definition of operations is extended compared with Basic SDL-2010.

Concrete grammar

<operation definitions> ::=
{ <operation definition>
| <operation reference>
| <external operation definition> }*

The Basic SDL-2010 <operation definitions> is extended to allow operations to be defined textually and to allow operations to be defined external to the <sdl specification> possibly in another notation.

<operation definition> ::=
{<package use clause>}*
<operation heading>
[<end> <entity in operation>+]
[<comment body>] <left curly bracket>
<statements>
<right curly bracket>

<operation heading> ::=
{ **operator** | **method** } <operation preamble> [<qualifier>] <operation name>
[<formal operation parameters>]
[<operation result>]

The Basic SDL-2010 <operation heading> is extended to allow operations to be methods. If an <operation heading> contains **operator**, each item in the <formal operation parameters> represents an item (in the same order) in the *Procedure-formal-parameter* list of the *Procedure-definition* for the operator. If an <operation heading> contains **method**, the first item in the *Procedure-formal-parameter* list of the *Procedure-definition* for the method is an *Inout-parameter* with an anonymous *Variable-name* and a *Sort-reference-identifier* that identifies the sort for the data type in which the method is defined. Each item in the <formal operation parameters> of a method represents a subsequent item (in the same order) in the *Procedure-formal-parameter* list of the *Procedure-definition* for the method.

NOTE 1 – <operation preamble> is placed after the keyword **operator** or **method** to avoid ambiguity with the optional <operation signatures> which also sometimes starts with an <operation preamble>. The initial keyword of <operation definitions> is never the same as the initial keyword of an <operation signature>.

If the <operation heading> begins with the keyword **operator**, then <operation definition> defines the behaviour of an operator. If the <operation heading> begins with the keyword **method**, then <operation definition> defines the behaviour of a method. Whether an operation is an operator or method is part of the operation signature and therefore part of the identity of the operation.

<formal operation parameters> ::=
(<operation parameters> {**,** <operation parameters> }* **)**
| [<end>] **fpar** <operation parameters> {**,** <operation parameters> }*

The Basic SDL-2010 <formal operation parameters> is extended to allow a bracketed list as an alternative.

<entity in operation> ::=
<data definition>
| <variable definition>
| <select definition>
| <macro definition>

<operation result> ::=
{ <result sign> | **returns** } <result aggregation> [<variable name>] <sort>

NOTE 2 – The <variable name> has to be a <name>, whereas in SDL-2000 it is allowed to be a name or a number (an <integer name> or a <real name>).

The Basic SDL-2010 <operation result> is extended to allow the use of <result sign> instead of **returns** and a <variable name> for the result. If the result is named, the derived procedure has a result named with this name.

<external operation definition> ::=
{ **operator** | **method** } <operation signature> external <end>

It is allowed to omit <arguments> and <result> in <external operation definition> if there is no other <external operation definition> within the same sort which has the same name, and an <operation signature> is present. In this case, the <arguments> and the <result> are derived from the <operation signature>.

An <external operation definition> is an operator or method whose behaviour body is not included in the <sdl specification> and is not necessarily defined in the SDL-2010 or any previous version of the Specification and Description Language. An <external operation definition> is treated as if written in SDL-2010, and the <operation signature> represents *Operation-signature* and *Procedure-definition* for the operation. How the abstract grammar for external operations is derived is not further defined by SDL-2010.

For each <operation signature>, at most one corresponding <operation definition> shall be given. No operation shall have both an <operation diagram> and an <operation definition>.

The <statement>s in <operation definition> shall contain neither an <imperative expression> nor an <identifier> defined outside the enclosing <operation definition> or <operation diagram>, respectively, except for <synonym identifier>s, <operation identifier>s, <literal identifier>s and <sort>s.

The list of <variable name>s is considered to bind tighter than the list of <operation parameters> within <formal operation parameters>.

<operation diagram> ::=
{ <operation page> }+

The Basic SDL-2010 <operation diagram> is extended to allow more than one page. The same <qualifier> shall be present on all pages or be omitted in the <operation heading> on every page. If the operation is unambiguously identified from the <operation heading> by <operation name> or <qualifier> and <operation name>, the <formal operation parameters> and <operation result> need to be included in the <operation heading> of at least one page. The <formal operation parameters> and <operation result> of the <operation heading> shall be the same on every page on which they

are included. Similarly, the <package use area> needs to be included for at least one page, and shall be the same for every page for which it is included.

<operation text area> ::=
<text symbol> ***contains***
{ <data definition>
| <variable definition>
| <macro definition>
| <select definition> }*

The Basic SDL-2010 <operation text area> is extended to allow <macro definition> and <select definition>.

The <start area> in <operation diagram> shall not contain <virtuality>.

The restriction on an <operation body area> containing an <identifier> (in Basic SDL-2010), does not apply for any <synonym identifier>.

NOTE 3 – <synonym identifier> is not mentioned in [ITU-T Z.101] because synonym is not included in Basic SDL-2010.

Any <statement> in <operation definition> shall contain neither an <imperative expression> nor an <identifier> defined outside the enclosing <operation definition>, except for <synonym identifier>s, <operation identifier>s, <literal identifier>s and <sort>s.

The following syntax provides compatibility with SDL-92 models: <legacy operator definition>, <legacy operator reference> and <legacy external operator definition>. These are used instead of <operation definitions> within a <legacy data type definition> (see clause 12.1.8).

<legacy operator definition> ::=
{<package use clause>}*
<operation heading> <end>
{ <entity in operation> }*
<start>
endoperator
[{<operation identifier> | <operation name> }] <end>

<start> is defined in [ITU-T Z.106].

The body of the transition for <start> in a <legacy operator definition> shall contain only items that are allowed in an operation definition.

<legacy operator reference> ::=
<operation heading> **referenced** <end>

<legacy external operator definition> ::=
operator <operation name> [<legacy procedure signature>] **external** <end>

An <operation heading> in a <legacy operator definition> or <legacy operator reference> shall use the keyword **operator**.

A <legacy operator definition> corresponds to an <operation definition>.

A <legacy operator reference> corresponds to an <operation reference>.

A <legacy external operator definition> corresponds to an <external operation definition>.

Semantics

An operator is a constructor for elements of the sort identified by the result. It returns a value, or an agent identity. A method acts on a variable (through the implicit in/out parameter of the procedure for the method) and therefore is able to modify the variable it acts on. In addition, a method optionally produces a result.

An operator shall not modify items that are reachable by the actual parameters. An item is considered modified in an operator if there is a potential control flow inside the operator resulting in modification. Therefore an operator with an *Inout-parameter* (in the procedure for the operator) shall not assign a value to the parameter or call a method or procedure that modifies the parameter.

NOTE 4 – For an operator there is no practical difference between an *In-parameter* and an *Inout-parameter* (though they will be probably have different implementations); an *Out-parameter* has no practical use, and an operator returns only a single value (whereas for a procedure each *Inout-parameter* or *Out-parameter* returns a value so a procedure is able to return multiple values).

Model

If an <operation diagram> has more than one <operation page>, this is transformed to a single page that contains every <operation text area> and <operation body area> on the pages. This transformation takes place before the contents of the page are transformed.

For every <operation definition> or <operation diagram> which does not have a corresponding <operation signature>, an <operation signature> is constructed.

A <operation definition> is derived syntax for an <operation diagram> that is a single <operation page> having the same <operation heading> and the <package use area> of the <operation diagram> containing the <package use clause> of the <operation definition>. The <transition area> of the <procedure start area> of the <operation page> consists of a <task area> containing the <statements> of the <operation definition> followed by an unlabelled <return area>. The <entity in operation> items of the <operation definition> are inserted into an <operation text area> of the <operation diagram> and allow local items (such as variables) of the operation to be defined. This transformation takes place after handling any <select definition> or <macro definition> in the <operation definition>.

12.1.8 Additional data definition constructs

This clause covers further constructs for data.

12.1.8.1 Syntypes

Concrete grammar

<syntype> ::=
<syntype identifier>

<syntype definition> ::=
{<package use clause>}*
syntype <syntype name> <equals sign> <parent sort identifier>
[<comment body>] <left curly bracket>
[{ <default initialization> [[<end>] <constraint>] | <constraint> } <end>]
<right curly bracket>
| {<package use clause>}*
<type preamble> <data type heading> [<data type specialization>]
[<comment body>] <left curly bracket>
<data type definition body> <constraint> <end>
<right curly bracket>

The Basic SDL-2010 <syntype definition> is extended to allow definition in the form of a data type definition with a constraint. Such a <syntype definition> with a <data type heading> has the keywords **value type** and is derived syntax defined below. In addition there is a legacy form used on legacy data type definitions.

<legacy syntype definition> ::=
syntype
<syntype name> = <parent sort identifier>
[<default initialization> [<end>]]
[**constants** <range condition>]
endsyntype [<syntype name>]

NOTE – See also <legacy data type definition> for syntype combined with a **newtype**.

Semantics

Specifying a syntype identifier is the same as specifying the parent sort identifier of the syntype, except for the cases listed for Basic SDL-2010 in [ITU-T Z.101] and following cases:

a) remote variable or remote procedure definition if one of the sorts for derivation of implicit signals is a syntype;

b) <any expression>, where the result will be within the range (see clause 12.3.4.6).

Model

A <syntype definition> with the keywords **value type** or **newtype** is distinguished from a <data type definition> by the inclusion of a <constraint>. Such a <syntype definition> is shorthand for introducing a <data type definition> with an anonymous name without the <constraint>, followed by a <syntype definition> with the keyword **syntype** based on this anonymously named sort and including <constraint>.

12.1.8.2 Constraint

Constraint on data types is defined in Basic SDL-2010.

12.1.8.3 Synonym definition

A synonym allows the use of a name for a constant that represents one of the data items of a sort. This is achieved by storing the value in a read-only variable.

Concrete grammar

<synonym definition> ::=
synonym <synonym definition item> { **,** <synonym definition item> }*<end>

<synonym definition item> ::=
<internal synonym definition item>
| <external synonym definition item>

<internal synonym definition item> ::=
<synonym name> [<sort>] <equals sign> <constant expression>

<external synonym definition item> ::=
<synonym name> <predefined sort> = **external**

NOTE 1 – The <synonym name> has to be a <name>, whereas in SDL-2000 it is allowed to be a name or a number (an <integer name> or a <real name>).

A <synonym definition> represents a *Variable-definition* in the context in which the synonym definition appears with special property that the variable is read-only. The <synonym name> represents the *Variable-name*.

NOTE 2 – Writing to the variable is not possible, because <synonym> is not allowed where assignments could take place.

The <constant expression> in the concrete syntax denotes a *Constant-expression* in the abstract syntax as defined in [ITU-T Z.101].

If sort of the <constant expression> is unique (that is, the expression belongs to only one sort), it is allowed to omit the <sort> in the <synonym definition> and the *Sort-reference-identifier* is derived from the constant expression sort.

If a <sort> is specified, this determines the *Sort-reference-identifier* of the *Variable-definition*, and the sort of the <constant expression> has to be compatible with this sort.

The <constant expression> shall not refer to the synonym defined by the <synonym definition>, either directly or indirectly (via another synonym).

An <external synonym definition item> defines a <synonym> whose result is not defined in a specification (see clause 13).

Semantics

The result that the synonym represents is determined by the context in which the synonym definition appears.

A synonym has a value, which is the value of the constant expression associated with the variable for the synonym definition.

A synonym has a sort, which is the sort of the variable for the synonym definition.

12.1.8.4 Restricted visibility

Concrete grammar

<visibility> ::=
public | **protected** | **private**

A <visibility> keyword shall not precede a <literal list>, <structure definition>, or <choice definition> in a <data type definition> containing <data type specialization>. A <visibility> keyword shall not be used in an <operation signature> that redefines an inherited operation signature.

Semantics

<visibility> controls visibility of a literal name or operation name.

When a <literal list> is preceded by <visibility>, this <visibility> applies to all <literal signature>s. When a <structure definition> or <choice definition> is preceded by <visibility>, then this <visibility> applies to all implied <operation signatures>.

When a <fields of sort> or <choice of sort> is preceded by a <visibility>, this <visibility> applies to all implied <operation signatures>.

If a <literal signature> or <operation signature> contains the keyword **private** in <visibility>, then the *Operation-name* derived from this <operation signature> is only visible within the scope of the <data type definition> that contains the <operation signature>. When a <data type definition> containing such <operation signature> is specialized, the <operation name> in <operation signature> is implicitly renamed to an anonymous name. Every occurrence of this <operation name> within the <operation definitions> or <operation diagram>s corresponding to this <operation signature> is renamed to the same anonymous name, when the <operation signature> and the corresponding operation definition are inherited by specialization.

NOTE 1 – As a consequence, the operator or method defined by this <operation signature> is usable only in operation applications within the data type definition that originally defined this <operation signature>, but not in any subtype thereof.

If a <literal signature> or <operation signature> contains the keyword **protected** in <visibility>, then the *Operation-name* derived from this <operation signature> is only visible within the scope of the <data type definition> that contains the <operation signature>.

NOTE 2 – Because inherited operators and methods are copied into the body of the subtype, the operator or method defined by this <operation signature> is accessible within the scope of any <data type definition> that is a subtype of the <data type definition> that originally defined this <operation signature>.

NOTE 3 – If a <literal signature> or <operation signature> does not contain <visibility>, the *Operation-name* derived from this <operation signature> is visible everywhere where the <sort name> that is defined in the enclosing <data type definition> is visible.

Model

If a <literal signature> or <operation signature> contains the keyword **public** in <visibility>, this is derived syntax for a signature having no visibility protection.

12.1.9 Specialization of data types

Specialization allows the definition of a data type based on another (super) type.

Concrete grammar

<data type specialization> ::=
inherits <data type type expression> [<renaming> | <legacy data inheritance>] [**adding**]

<interface specialization> ::=
inherits <interface type expression> { **,** <interface type expression> }* [**adding**]

<renaming> ::=
(<rename list> **)**

<rename list> ::=
<rename pair> { **,** <rename pair> }*

<rename pair> ::=
<operation name> <equals sign> <base type operation name>
| <literal name> <equals sign> <base type literal name>

The *Data-type-identifier* in the *Data-type-definition* represented by the <data type definition> in which <data type specialization> (or <interface specialization>) is contained identifies the data type represented by the <data type type expression> in its <data type specialization> (see also clause 8.1.2).

An *Interface-definition* is allowed to contain a list of *Data-type-identifier* items. The interface denoted by an *Interface-definition* is a specialization of all the interfaces denoted by the *Data-type-identifier* list items.

The resulting content of a specialized interface definition (<interface specialization>) consists of the content of the supertypes with the content of the specialized definition. This implies that the set of signals, remote procedures and remote variables of the specialized interface definition is the union of those given in the specialized definition itself and those of the supertypes. The resulting set of definitions shall obey the rules for visibility and naming.

The <data type constructor> with the <data type specialization> shall not be a <choice definition>.

The <data type constructor> with the <data type specialization> shall be of the same kind as the <data type constructor> used in the <data type definition> of the sort referenced by <data type type expression> in the <data type specialization>: if the <data type constructor> used in a (direct or indirect) supertype was a <literal list>, the <data type constructor> shall be a <literal list>; if the supertype <data type constructor> was a <structure definition>, the subtype <data type constructor> shall be a <structure definition>.

Renaming changes the name of inherited literals, operators and methods in the derived data type.

All <literal name>s and all <base type literal name>s in a <rename list> shall be distinct.

All <operation name>s and all <base type operation name>s in a <rename list> shall be distinct.

A <base type operation name> specified in a <rename list> shall be an operation with <operation name> defined in the data type definition defining the <base type> of <data type type expression>.

A subtype with <renaming> of a sort is not allowed as an actual sort parameter constrained to be **atleast** that sort.

<legacy data inheritance> ::=
[<legacy literal renaming>]
[[**operators**] { **all** | (<legacy inheritance list>) } [<end>]]

<legacy literal renaming> ::=
literals <rename pair> { , <rename pair> }* <end>

<legacy inheritance list> ::=
<legacy inherited operator> { , <legacy inherited operator> }*

<legacy inherited operator> ::=
<operation name> | <rename pair>

To be consistent with SDL-92, the <operation name> as <legacy inherited operator> should refer to an operation defined in the base type. Specifying **operators all** or a named operation without renaming has no influence on the inherited operations, which are determined according to the SDL-2010 rules.

Semantics

The sort defined by the specialization is a subsort of the sort defined by the base type. The sort defined by the base type is a supersort of the sort defined by the specialization.

Sort compatibility determines when a sort is allowed to be used in place of another sort, and when it is not allowed. This relation is used for assignments (see clause 12.3.3), for parameter passing (see clause 12.2.7 and clause 9.4 in [ITU-T Z.101]), for re-declaration of parameter types during inheritance (see clause 12.1.2), and for actual context parameters.

Let P and R be two sorts. P is sort compatible with R if and only if one of the following apply:

a) P and R are the same sort;

b) P is directly sort compatible with R (see below);

c) R is denoted by a <sort identifier>, and the <sort identifier> is defined by an interface (a pid sort) and, for some sort Q, P is sort compatible with Q, and Q is sort compatible with R.

NOTE 1 – Sort compatibility is transitive only for sorts defined by interfaces, not for sorts defined by value types.

Let Y and Z be different sorts. Y is directly sort compatible with Z if and only if any of the following applies:

a) Y is denoted by an <anchored sort> of the form **this** Z;

b) Y is denoted by a <pid sort> (see clause 12.1.2 and clause 12.1.2 of [ITU-T Z.101]) and Zis a supersort of Y.

Model

The model for specialization in clause 8.4 of [ITU-T Z.102] is used, augmented as follows.

A specialized data type is based on another (base) data type by using a <data type definition> in combination with a <data type specialization>. The sort defined by the specialization is disjoint from the sort defined by the base type.

If the sort defined by the base type has literals defined, the literal names are inherited as names for literals of the sort defined by the specialized type unless literal renaming has taken place for that literal. Literal renaming has taken place for a literal if the base type literal name appears as the second name in a <rename pair>, in which case the literal is renamed to the first name in that pair.

If the base type has operators or methods defined, the operation names are inherited as names for operators or methods of the sort being defined, unless the operator or method has been declared as private (see clause 12.1.8.4) or operation renaming has taken place for that operator or method. Operation renaming has taken place for an operator or method if the inherited operation name

appears as the second name in a <rename pair>, in which case the operator or method is renamed to the first name in that pair.

When several operators or methods of the <base type> of <sort type expression> have the same name as the <base type operation name> in a <rename pair>, all of these operators or methods are renamed.

In the following paragraphs ST is a subsort that is a specialization of a data type T, and Tid is an <identifier> for data type T.

For every occurrence of a <basic sort> of the form Tid in a specialization of T, the <basic sort> is replaced by the <basic sort> of the subsort: that is, the <identifier> for the subsort ST. Every occurrence of an <anchored sort> of the form **parent** T in a specialization of T, the <anchored sort> is replaced by an unambiguous qualified <basic sort> for the base type T: that is, the qualified <identifier> for T. Every occurrence of an <anchored sort> of the form **this** T in a specialization of T, the <anchored sort> is replaced by the <basic sort> of the subsort: that is, the <identifier> for the subsort ST. The signature of the operator is considered only after applying these models.

NOTE 2 – An <anchored sort> of the form **this** T has the same meaning as a <basic sort> of the form T in a specialization of T.

The argument sorts and result of an inherited operator or method are the same as those of the corresponding operator or method of the base type, except that every <argument> in T that is a <basic sort> T or an <anchored sort> of the form **this** T, in the inherited operator or method the <anchored sort> is replaced by the subsort ST.

NOTE 3 – From the specialization semantics in clause 8.4 of [ITU-T Z.102], an operator of T is inherited if the signature of the specialized operator is not the same as the signature of the base type operator, which is only the case if the signature is not already present in the subtype (a generic operator or matching user defined operator) or renaming had taken place, or contains one <sort> that is: a <basic sort> of the form T, or an <anchored sort> of the form **this** T.

12.2 Use of data

The following clauses define how sorts, literals, operators, methods and synonyms are interpreted in expressions.

12.2.1 Expression and expressions as actual parameters

Abstract grammar

Active-expression	=	*Variable-access*
	|	*Conditional-expression*
	|	*Operation-application*
	|	*Equality-expression*
	|	*Imperative-expression*
	|	*Range-check-expression*
	|	*Value-returning-call-node*
	|	*Encoding-expression*
	|	*Decoding-expression*

The Basic SDL-2010 *Active-expression* is extended to include *Encoding-expression* and *Decoding-expression* to allow the use of encoding rules in expressions.

Concrete grammar

For simplicity of description, no distinction is made between the concrete syntax of *Constant-expression* and *Active-expression.*

<expression0> ::=
<operand>
| <create expression>
| <value returning procedure call>
| <decoding expression>
| <encoding expression>

The Basic SDL-2010 <expression0> is extended to allow a <create expression>, <decoding expression> and <encoding expression>. An <expression0> that contains a <create expression>, <decoding expression> or <encoding expression> is not a <constant expression0> and represents an *Active-expression*.

<primary> ::=
<operation application>
| <literal>
| **(** <expression> **)**
| <conditional expression>
| <extended primary>
| <active primary>
| <synonym>

The Basic SDL-2010 <primary> is extended to allow a <synonym>.

12.2.2 Literal

Literal is as defined in Basic SDL-2010.

12.2.3 Extended primary

Extended variable for an indexed primary is as defined in Basic SDL-2010, but for a field primary it is extended.

Concrete grammar

<field primary> ::=
<primary> <exclamation mark> { <field name> | <field number> | <as signal> }
| <primary> <full stop> { <field name> | <field number> | <as signal> }
| <field name>

The <field primary> of Basic SDL-2010 is extended to allow the field to be identified by a <field number> or <as signal>, and to allow the short form <field name> (without a <primary>) in the context of the data type definition of a structure or choice data type.

<field number> ::=
<integer name>

The <field number> shall represent a non-zero, positive integer less than or equal to the number of fields in the sort of the variable.

Model

A <field number> is a shorthand for the <field name> of a field of the sort of <field primary> or <field variable>, where the value 1 represents the first field, the value 2 the second field and the value n the nth field.

NOTE 1 – A <field number> has to be used to denote a field of a structure data type with a unique anonymous name introduced by a <signal definition> with a <sort list>.

An <as signal> in a <field primary> is a shorthand for the <field name> of a field of the sort of the <field primary>, where the <field name> is the same as the unique anonymous name of the structure data type defined by the <signal definition> identified by the <signal identifier> of <as signal>.

NOTE 2 – An <as signal> has to be used to denote a field of a choice data type introduced by a <gate definition>, <channel definition area> or <interface definition>.

When the <field primary> has the form <field name>, this is derived syntax for

this ! <field name>

12.2.4 Equality expression

Equality expression is as defined in Basic SDL-2010.

12.2.5 Conditional expression

Conditional expression is as defined in Basic SDL-2010.

12.2.6 Operation application

Operation application is extended to include the syntax for method application

Concrete grammar

<operation application> ::=
<operator application>
| <method application>

<method application> ::=
<primary> <full stop> <operation identifier> [<actual parameters>]

A <method application> is legal concrete syntax only if <operation identifier> represents a method.

When the operation application in an operation has the syntactical form:

<operation identifier> [<actual parameters>]

then, during derivation of the *Operation-identifier* from context, the form:

this <full stop> <operation identifier> [<actual parameters>]

is also considered if the <operation identifier> could identify a method. The model in clause 12.3.2 is applied before resolution by context is attempted.

An <expression> in <actual parameters> corresponding to an *Inout-parameter* or *Out-parameter* in the *Procedure-definition* associated with the *Operation-signature* shall not be omitted and shall be a <variable access> or <extended primary>.

NOTE – Omission of <actual parameters> is allowed in an <operation application> if all actual parameters have been omitted.

Model

The concrete syntax form:

<expression> <full stop> <operation identifier> [<actual parameters>]

is derived concrete syntax for:

<operation identifier> *new-actual-parameters*

where *new-actual-parameters* is <actual parameters> containing only <expression>, if <actual parameters> was not present; otherwise, *new-actual-parameters* is obtained by inserting <expression> before the first optional expression in <actual parameters>.

If the <primary> of a <method application> is not a variable or **this**, there is an implicit assignment of the <primary> to an implicit variable with the sort of the first parameter of the operation (that is, the method sort). The assignment is placed before the action in which the <method application> occurs. The implicit variable replaces the <primary> in the <method application>.

12.2.7 Range check expression

Range check expression is as defined in Basic SDL-2010.

12.2.8 Synonym

Concrete grammar

<synonym> ::=
<synonym identifier>

A <synonym> represents the *Variable-identifier* for the identified synonym.

Semantics

The synonym gives the *Constant-expression* associated with the variable.

12.3 Active use of data

This clause extends the active use of data defined in Basic SDL-2010.

12.3.1 Variable definition

Variable definition is as defined in Basic SDL-2010 and (for exported) Comprehensive SDL-2010 with the extension to include remote variables defined in an interface definition.

Concrete grammar

The <remote variable identifier> of an <exported as> is defined by a <remote variable definition> (as in Comprehensive SDL-2010), or in an <interface variable definition> of an explicit <interface definition>.

12.3.2 Variable access

Variable access is extended from Basic SDL-2010 and (for import expression) Comprehensive SDL-2010 to allow the use of **this**.

Concrete grammar

<variable access> ::=
<variable identifier>
| <import expression>
| **this**

this shall only occur in method definitions.

Model

The transformation for <import expression> is given in Comprehensive SDL-2010.

A <variable access> using the keyword **this** is replaced by the anonymous name introduced as the name of the leading parameter in <arguments> according to clause 12.1.3.

12.3.3 Assignment

Assignment is as defined in Basic SDL-2010.

12.3.3.1 Extended variable

Extended variable for an indexed variable is as defined in Basic SDL-2010, but for a field variable it is extended to allow the field to be identified by an integer instead of a name.

Concrete grammar

<field variable> ::=
<variable> <exclamation mark> { <field name> | <field number> | <as signal> }
| <variable> <full stop> { <field name> | <field number> | <as signal> }

The <field variable> of Basic SDL-2010 is extended to allow the field to be identified by a <field number> or <as signal>.

Model

An <as signal> in a <field variable> is a shorthand for the <field name> of a field of the sort of the <field variable>, where the <field name> is the same as the unique anonymous name of the structure data type defined by the <signal definition> identified by the <signal identifier> of <as signal>.

NOTE – An <as signal> has to be used to denote a field of a choice data type introduced by a <gate definition>, <channel definition area> or <interface definition>.

12.3.3.2 Default initialization

Default initialization of Basic SDL-2010 is extended to include virtuality.

Default initialization of instances of a data type is specified to be virtual by means of the keyword **virtual** in <virtuality>. Redefinition of a virtual default initialization is allowed in specializations. This is indicated by a default initialization with the keyword **redefined** or **finalized** in <virtuality>.

If the derived type contains no <constant expression> in its default initialization, then the derived type does not have a default initialization.

Concrete grammar

<default initialization> ::=
default [<virtuality>] [<constant expression>]

The <constant expression> shall only be omitted if <virtuality> is **redefined** or **finalized**. In that case the *Data-type-definition* or *Syntype-definition* for the <data type definition> or <syntype definition> including the <default initialization> has no *Default-initialization*.

Semantics

Redefinition of a virtual default initialization corresponds closely to redefinition of virtual types (see clause 8.4.2 of [ITU-T Z.102]).

12.3.4 Imperative expression

The transformations described in the *Model* of this clause are made at the same time as the expansion for import is made. A label attached to an action in which an imperative expression appears is moved to the first task inserted during the described transformation. If several imperative expressions appear in an expression, the tasks are inserted in the same order as the imperative expressions appear in the expression.

Abstract grammar

Imperative-expression	=	*Now-expression*
	|	*Pid-expression*
	|	*Timer-active-expression*
	|	*Timer-remaining-duration*
	|	*Active-agents-expression*
	|	*Any-expression*
	|	*State-expression*

The grammar is extended to include *Any-expression* and *State-expression*.

Concrete grammar

<imperative expression> ::=
<now expression>
| <pid expression>
| <timer active expression>
| <timer remaining duration>
| <active agents expression>
| <any expression>
| <state expression>

12.3.4.1 Now expression

Now expression is as defined in Basic SDL-2010.

12.3.4.2 Pid expression

The `Pid` expression is extended to describe the <create expression> that returns the offspring pid value.

Concrete grammar

<create expression> ::=
create <create body>

A <create body> of a <create expression> represents the *Create-request-node* of the *Create-request-node* as a *Graph-node* as described in [ITU-T Z.101] in the model for <create expression>.

Model

The use of <create expression> in an expression is shorthand for inserting a *Create-request-node* just before the action where the <create expression> occurs, followed by an assignment of **offspring** to an implicitly declared anonymous variable with a `Pid` sort. An access to the implicit variable then replaces the <create expression> in the expression. If <create expression> occurs several times in an expression, one distinct variable is used for each occurrence. In this case, the order of the inserted create requests and variable assignments is the same as the order of the <create expression> items.

If the <create expression> contains an <agent type identifier>, then the transformations that are applied to a create statement that contains an <agent type identifier> are also applied to the implicit create statements resulting from the transformation of a <create expression>.

12.3.4.3 Timer active expression and timer remaining duration

The timer active expression and timer remaining duration are as defined in Basic SDL-2010.

12.3.4.4 Active agents expression

The active agents expression is as defined in Basic SDL-2010.

12.3.4.5 Import expression

Concrete grammar

The concrete syntax for an import expression is defined in [ITU-T Z.102].

Semantics

In addition to the semantics defined in [ITU-T Z.102], an import expression is interpreted as a variable access (see clause 12.3.2) to the implicit variable for the import expression.

Model

The import expression has implied syntax for the importing of the result as defined in [ITU-T Z.102] and also has an implied *Variable-access* of the implied variable for the import in the context where the <import expression> appears.

The use of <import expression> in an expression is shorthand for inserting a task just before the action, where the expression occurs which assigns to an implicit variable the result of the <import expression> and then uses that implicit variable in the expression. If <import expression> occurs several times in an expression, one variable is used for each occurrence.

12.3.4.6 Any expression

Any-expression is useful for modelling behaviour, where stating a specific data item would imply over-specification. From a result returned by an *Any-expression*, no assumption should be made on other results returned by subsequent or previous interpretation of the *Any-expression*.

Abstract grammar

Any-expression :: *Sort-reference-identifier*

If the *Sort-reference-identifier* references the *Sort* of a *Value-data-type-definition* there shall be at least one element in the set of values for the data type.

If the *Sort-reference-identifier* denotes *a Syntype-identifier*, the *Parent-sort-identifier* of the identified *Syntype-definition* shall identify the *Sort* of a *Value-data-type-definition* for which there is at least one element in the set of values for the data type.

Concrete grammar

<any expression> ::=
any (<sort> **)**

Semantics

An *Any-expression* returns an unspecified element of the sort or syntype designated by *Sort-reference-identifier*, if that sort or syntype is a value sort. If *Sort-reference-identifier* denotes *a Syntype-identifier*, the result will be within the range of that syntype. If the sort or syntype designated by *Sort-reference-identifier* is a pid sort, the *Any-expression* returns `Null`.

12.3.4.7 State expression

Abstract grammar

State-expression :: { }

Concrete grammar

<state expression> ::=
state

Semantics

A state expression is interpreted as a `Charstring` that contains the spelling of the name of the most recently entered state of the nearest enclosing scope unit. If there is no such state, <state expression> denotes the empty string (").

12.3.5 Value returning procedure call

The value returning procedure call is extended to cover remote procedures, additional constraints and transformation models.

Concrete grammar

<value returning procedure call> ::=
[**call**] <procedure call body>
| [**call**] <remote procedure call body>

A <value returning procedure call> shall not occur in the <Boolean expression> of a <continuous signal area> or <enabling condition area>.

An <expression> in <actual parameters> corresponding to a formal **in**/**out** or **out** parameter shall not be omitted and shall be a <variable identifier>.

After the *Model* for **this** has been applied, the <procedure identifier> shall denote a procedure that contains a start transition.

If **this** is used, <procedure identifier> shall denote an enclosing procedure.

The <remote procedure call body> represents a *Value-returning-call-node*, where *Procedure-identifier* contains only the *Procedure-identifier* of the procedure implicitly defined by the *Model* below.

Model

If the <procedure identifier> is not defined within the enclosing agent, the procedure call is transformed into a call of a local, implicitly created, subtype of the procedure.

this implies that when the procedure is specialized, the <procedure identifier> is replaced by the identifier of the specialized procedure.

When the <value returning procedure call> contains a <remote procedure call body>, a procedure with an anonymous name is implicitly defined, where <sort> in <procedure result> of the procedure definition denoted by the <procedure identifier> is the return sort of this anonymous procedure. This anonymous procedure has a single <start area> containing a <return area> with <remote procedure call body> as its <expression>.

NOTE – This transformation is not applied again to the implicit procedure definition.

13 Generic system definition

See [ITU-T Z.102] and [ITU-T Z.103].

14 Package Predefined

In the following definitions, all references to names defined in the package Predefined are considered to be treated as prefixed by the qualification <<**package** Predefined>>. To increase readability, this qualification is omitted.

The extensions defined in Annex A are used in this package. It is not allowed to use these extensions in a user defined package.

```
/* */
package Predefined
/*
```

14.1 Boolean sort

14.1.1 Definition

```
*/
value type Boolean;
  literals true, false;
  operators
      "not" ( this Boolean )                -> this Boolean;
      "and" ( this Boolean, this Boolean )  -> this Boolean;
      "or"  ( this Boolean, this Boolean )  -> this Boolean;
      "xor" ( this Boolean, this Boolean )  -> this Boolean;
      "=>"  ( this Boolean, this Boolean )  -> this Boolean;
axioms
     not( true  )     == false;
     not( false )     == true ;
/* */
     true  and true   == true ;
     true  and false  == false;
     false and true   == false;
     false and false  == false;
/* */
```

```
     true or true    == true ;
     true or false   == true ;
     false or true   == true ;
     false or false  == false;
/* */
     true xor true   == false;
     true xor false  == true ;
     false xor true  == true ;
     false xor false == false;
/* */
     true => true    == true ;
     true => false   == false;
     false => true   == true ;
     false => false  == true ;
endvalue type Boolean;
/*
```

14.1.2 Usage

The Boolean sort is used to represent `true` and `false` values. Often it is used as the result of a comparison.
The Boolean sort is used widely throughout SDL.

14.2 Character sort

14.2.1 Definition

```
*/
value type Character;
  literals
     NUL, SOH, STX, ETX, EOT, ENQ, ACK, BEL,
     BS,  HT,  LF,  VT,  FF,  CR,  SO,  SI,
     DLE, DC1, DC2, DC3, DC4, NAK, SYN, ETB,
     CAN, EM,  SUB, ESC, IS4, IS3, IS2, IS1,
     ' ', '!', '"', '#', '$', '%', '&', '''',
     '(', ')', '*', '+', ',', '-', '.', '/',
     '0', '1', '2', '3', '4', '5', '6', '7',
     '8', '9', ':', ';', '<', '=', '>', '?',
     '@', 'A', 'B', 'C', 'D', 'E', 'F', 'G',
     'H', 'I', 'J', 'K', 'L', 'M', 'N', 'O',
     'P', 'Q', 'R', 'S', 'T', 'U', 'V', 'W',
     'X', 'Y', 'Z', '[', '\', ']', '^', '_',
     '`', 'a', 'b', 'c', 'd', 'e', 'f', 'g',
     'h', 'i', 'j', 'k', 'l', 'm', 'n', 'o',
     'p', 'q', 'r', 's', 't', 'u', 'v', 'w',
     'x', 'y', 'z', '{', '|', '}', '~', DEL;
/* '''' is an apostrophe, ' ' is a space, '~' is a tilde */
/* */
  operators
     chr  ( Integer ) -> this Character;
/*   "<", "<=", ">", ">=", and "num" are implicitly defined (see clause 12.1.6.1). */
axioms
     for all a,b in Character (
     for all i in Integer (
/* definition of Chr */
     chr(num(a)) == a;
     chr(i+128)  == chr(i);
     ));
endvalue type Character;
/*
```

14.2.2 Usage

The Character sort is used to represent characters of the International Reference Alphabet [ITU-T T.50].

14.3 String sort

14.3.1 Definition

```
*/
value type String < type Itemsort >;
/* Strings are "indexed" from one */
  operators
     emptystring                                    -> this String;
     mkstring   ( Itemsort                    )  -> this String;
     Make       ( Itemsort                    )  -> this String;
     length     ( this String                    ) -> Integer;
     first      ( this String                    ) -> Itemsort;
     last       ( this String                    ) -> Itemsort;
     "//"       ( this String, this String       ) -> this String;
     Extract    ( this String, Integer           ) -> Itemsort raise InvalidIndex;
     Modify     ( this String, Integer, Itemsort ) -> this String;
     substring  ( this String, Integer, Integer ) -> this String raise InvalidIndex;
     /* substring (s,i,j) gives a string of length j starting from the ith element */
     remove     ( this String, Integer, Integer ) -> this String;
     /* remove (s,i,j) gives a string with a substring of length j starting from
        the ith element removed */
axioms
     for all e in Itemsort ( /*e - element of Itemsort*/
     for all s,s1,s2,s3 in String (
     for all i,j in Integer (
/* constructors are emptystring, mkstring, and "//" */
/* equalities between constructor terms */
     s // emptystring== s;
     emptystring // s== s;
     (s1 // s2) // s3== s1 // (s2 // s3);
/* */
/* definition of length by applying it to all constructors */
     <<type String>>length(emptystring)== 0;
     <<type String>>length(mkstring(e))== 1;
     <<type String>>length(s1 // s2)    == length(s1) + length(s2);
     Make(s)                              == mkstring(s);
/* */
/* definition of Extract by applying it to all constructors,
        error cases handled separately */
     Extract(mkstring(e),1)                    == e;
     i <= length(s1) ==> Extract(s1 // s2,i)== Extract(s1,i);
     i > length(s1) ==> Extract(s1 // s2,i)  == Extract(s2,i-length(s1));
     i<=0 or i>length(s) ==> Extract(s,i)   == raise InvalidIndex;
/* */
/* definition of first and last by other operations */
     first(s)   == Extract(s,1);
     last(s)    == Extract(s,length(s));
/* */
/* definition of substring(s,i,j) by induction on j,
        error cases handled separately */
     i>0 and i-1<=length(s)            ==>
                    substring(s,i,0)   == emptystring;
/* */
     i>0 and j>0 and i+j-1<=length(s)  ==>
                    substring(s,i,j)   == substring(s,i,j-1) // mkstring(Extract(s,i+j-1));
/* */
     i<=0 or j<0 or i+j-1>length(s)    ==>
                    substring(s,i,j)   == raise InvalidIndex;
/* */
/* definition of Modify by other operations */
     Modify(s,i,e) == substring(s,1,i-1) // mkstring(e) // substring(s,i+1,length(s)-i);
/* definition of remove */
     remove(s,i,j)   == substring(s,1,i-1) // substring(s,i+j,length(s)-i-j+1);
     )));
endvalue type String;
/*
```

14.3.2 Usage

The `Make`, `Extract`, and `Modify` operators will typically be used with the shorthand forms defined in clause 12.2.3 of [ITU-T Z.101] and clause 12.3.3.1 of [ITU-T Z.101] for accessing the values of strings and assigning values to strings.

14.4 Charstring sort

14.4.1 Definition

```
*/
value type Charstring
     inherits String < Character > ( '' = emptystring )
     adding ;
        operators ocs in nameclass
             '''' ( (' ':'&') or '''''' or ('(': '~') )+ '''' -> this Charstring;
/* character strings of any length of any characters from a space ' ' to a tilde '~' */
axioms
     for all c in Character nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)                    ==> cs == mkstring(c);
/* string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1                     ==> cs  == cs1 // cs2;
     ));
endvalue type Charstring;
/*
```

14.4.2 Usage

The Charstring sort defines strings of characters.

A Charstring literal is allowed to contain printing characters and spaces.

A non-printing character is usable as a string by using `mkstring`, for example, `mkstring(DEL)`.

Example:

```
synonym newline_prompt Charstring = mkstring(CR) // mkstring(LF) // '$>';
```

14.5 Integer sort

14.5.1 Definition

```
*/
value type Integer;
  literals unordered nameclass (('0':'9')*) ('0':'9'));
  operators
     "-"    ( this Integer                ) -> this Integer;
     "+"    ( this Integer, this Integer  ) -> this Integer;
     "-"    ( this Integer, this Integer  ) -> this Integer;
     "*"    ( this Integer, this Integer  ) -> this Integer;
     "/"    ( this Integer, this Integer  ) -> this Integer raise DivisionByZero;
     "mod"  ( this Integer, this Integer  ) -> this Integer raise DivisionByZero;
     "rem"  ( this Integer, this Integer  ) -> this Integer;
     "<"    ( this Integer, this Integer  ) -> Boolean;
     ">"    ( this Integer, this Integer  ) -> Boolean;
     "<="   ( this Integer, this Integer  ) -> Boolean;
     ">="   ( this Integer, this Integer  ) -> Boolean;
     power  ( this Integer, this Integer  ) -> this Integer;
     integer ( Integer                    ) -> this Integer;
     num     ( this Integer               ) -> Integer;
     bs in nameclass '''' ( (('0' or '1')*'''B') or ((('0':'9') or ('A':'F'))*'''H') )
              -> this Integer;
axioms  noequality
     for all a,b,c in this Integer (
/* constructors are 0, 1, +, and unary - */
/* equalities between constructor terms */
     (a + b) + c          == a + (b + c);
     a + b                == b + a;
     0 + a                == a;
     a + (- a)            == 0;
     (- a) + (- b)      == - (a + b);
```

```
     <<type Integer>> - 0 == 0;
     - (- a)              == a;
/* */
/* definition of binary "-" by other operations */
     a - b                == a + (- b);
/* */
/* definition of "*" by applying it to all constructors */
     0 * a                == 0;
     1 * a                == a;
     (- a) * b            == - (a * b);
     (a + b) * c          == a * c + b * c;
/* */
/* definition of "<" by applying it to all constructors */
     a < b                                    == 0 < (b - a);
     <<type Integer>> 0 < 0                   == false;
     <<type Integer>> 0 < 1                   == true ;
     0 < a           == true ==> 0 < (- a)    == false;
     0 < a and 0 < b == true ==> 0 < (a + b)  == true ;
/* */
/* definition of ">", "equal", "<=", and ">=" by other operations */
     a > b                == b < a;
     equal(a, b)          == not(a < b or a > b);
     a <= b               == a < b or a = b;
     a >= b               == a > b or a = b;
/* */
/* definition of "/" by other operations */
     a / 0                                          == raise DivisionByZero;
     a >= 0 and b > a           == true ==> a / b == 0;
     a >= 0 and b <= a and b > 0  == true ==> a / b == 1 + (a-b) / b;
     a >= 0 and b < 0           == true ==> a / b == - (a / (- b));
     a < 0 and b < 0              == true ==> a / b == (- a) / (- b);
     a < 0 and b > 0              == true ==> a / b == - ((- a) / b);
/* */
/* definition of "rem" by other operations */
     a rem b == a - b * (a/b);
/* */
/* definition of "mod" by other operations */
     a >= 0 and b > 0               ==> a mod b  == a rem b;
     b < 0                            ==> a mod b  == a mod (- b);
     a < 0 and b > 0 and a rem b = 0 ==> a mod b  == 0;
     a < 0 and b > 0 and a rem b < 0 ==> a mod b  == b + a rem b;
     a mod 0 == raise DivisionByZero;
/* */
/* definition of power by other operations */
     power(a, 0)            == 1;
     b > 0 ==> power(a, b)  == a * power(a, b-1);
     b < 0 ==> power(a, b)  == power(a, b+1) / a; );
/* */
/* definition of literals */
     <<type Integer>> 2 == 1 + 1;
     <<type Integer>> 3 == 2 + 1;
     <<type Integer>> 4 == 3 + 1;
     <<type Integer>> 5 == 4 + 1;
     <<type Integer>> 6 == 5 + 1;
     <<type Integer>> 7 == 6 + 1;
     <<type Integer>> 8 == 7 + 1;
     <<type Integer>> 9 == 8 + 1;
/* */
/* literals other than 0 to 9 */
     for all a,b,c in this Integer nameclass (
     spelling(a) == spelling(b) // spelling(c),
     length(spelling(c)) == 1               ==> a == b * (9 + 1) + c;
     );
/* */
/* hex and binary representation of Integer */
     for all b in Bitstring nameclass (
     for all i in bs nameclass (
     spelling(i) == spelling(b)             ==> i == <<type Bitstring>>num(b);
     ));
/* integer makes parent sort into subsort */
/* num makes subsort into parent sort */
```

```
     for all i in this Integer (
     for all p in Integer (
     spelling(i) == spelling(p)             ==> num(i) == p ;
     ));
     integer(num(i)) == i;
endvalue type Integer;
/*
```

14.5.2 Usage

The Integer sort is used for mathematical integers with decimal, hex, or binary notation.

14.6 Natural syntype

14.6.1 Definition

```
*/
syntype Natural = Integer constants >= 0; endsyntype Natural;
/*
```

14.6.2 Usage

The natural syntype is used when positive integers only are required. All operators will be the integer operators but when a value is used as a parameter or assigned, the value is checked. A negative value will be an error.

14.7 Real sort

14.7.1 Definition

```
*/
value type Real;
  literals unordered nameclass
        ( ('0':'9')* ('0':'9') ) or ( ('0':'9')* '.'('0':'9')+ );
  operators
     "-"   ( this Real             ) -> this Real;
     "+"   ( this Real, this Real  ) -> this Real;
     "-"   ( this Real, this Real  ) -> this Real;
     "*"   ( this Real, this Real  ) -> this Real;
     "/"   ( this Real, this Real  ) -> this Real raise DivisionByZero;
     "<"   ( this Real, this Real  ) -> Boolean;
     ">"   ( this Real, this Real  ) -> Boolean;
     "<="  ( this Real, this Real  ) -> Boolean;
     ">="  ( this Real, this Real  ) -> Boolean;
     float ( Integer               ) -> this Real;
     fix   ( this Real             ) -> Integer;
axioms noequality
     for all r,s in Real (
     for all a,b,c,d in Integer (
/* constructors are float and "/" */
/* equalities between constructor terms allow to reach always a form
        float(a) / float(b) where b > 0 */
     r / float(0)                                    == raise DivisionByZero;
     r / float(1)                                    == r;
     c /= 0 ==> float(a) / float(b)                  == float(a*c) / float(b*c);
     b /= 0 and d /= 0 ==>
     (float(a) / float(b)) / (float(c) / float(d))   == float(a*d) / float(b*c);
/* */
/* definition of unary "-" by applying it to all constructors */
     - (float(a) / float(b))                        == float(- a) / float(b);
/* */
/* definition of "+" by applying it to all constructors */
     (float(a) / float(b)) + (float(c) / float(d))  ==float(a*d + c*b) / float(b*d);
/* */
/* definition of binary "-" by other operations */
     r - s                                          == r + (- s);
/* */
/* definition of "*" by applying it to all constructors */
     (float(a) / float(b)) * (float(c) / float(d))  == float(a*c) / float(b*d);
/* */
```

```
/* definition of "<" by applying it to all constructors */
     b > 0 and d > 0 ==>
     (float(a) / float(b)) < (float(c) / float(d))  == a * d < c * b;
/* */
/* definition of ">", "equal", "<=", and ">=" by other operations */
     r > s       == s < r;
     equal(r, s) == not(r < s or r > s);
     r <= s      == r < s or r = s;
     r >= s      == r > s or r = s;
/* */
/* definition of fix by applying it to all constructors */
     a >= b and b > 0 ==> fix(float(a) / float(b)) == fix(float(a-b) / float(b)) + 1;
     b > a  and a >= 0  ==> fix(float(a) / float(b)) == 0;
     a < 0  and b > 0 ==> fix(float(a) / float(b)) == - fix(float(-a)/float(b)) - 1;));
/* */
     for all r,s in Real nameclass (
     for all i,j in Integer nameclass (
        spelling(r) == spelling(i)          ==> r == float(i);
/* */
        spelling(r) == spelling(i)          ==> i == fix(r);
/* */
        spelling(r) == spelling(i) // spelling(s),
        spelling(s) == '.' // spelling(j) ==> r == float(i) + s;
/* */
        spelling(r) == '.' // spelling(i),
        length(spelling(i)) == 1            ==> r == float(i) / 10;
/* */
        spelling(r) == '.' // spelling(i) // spelling(j),
        length(spelling(i)) == 1,
        spelling(s) == '.' // spelling(j) ==> r == (float(i) + s) / 10;
     ));
endvalue type Real;
/*
```

14.7.2 Usage

The real sort is used to represent real numbers.

The real sort represents all numbers which are able to be represented as one integer divided by another.

Numbers which are not able to be represented in this way (irrational numbers – for example, the square root of 2) are not part of the real sort. However, for practical engineering a sufficiently accurate approximation is usually usable.

14.8 Array sort

14.8.1 Definition

```
*/
value type Array < type Index; type Itemsort >;
  operators
     Make                                -> this Array ;
     Make   ( Itemsort                 ) -> this Array ;
     Modify ( this Array,Index,Itemsort ) -> this Array ;
     Extract( this Array,Index          ) -> Itemsort raise InvalidIndex;
axioms
     for all item, itemi, itemj in Itemsort (
     for all i, j in Index (
     for all a, s in Array (
        <<type Array>>Extract(make,i)                    == raise InvalidIndex;
        <<type Array>>Extract(make(item),i)              == item ;
        i = j ==> Modify(Modify(s,i,itemi),j,item)       == Modify(s,i,item);
        i = j ==> Extract(Modify(a,i,item),j)            == item ;
        i = j == false ==> Extract(Modify(a,i,item),j)   == Extract(a,j);
        i = j == false ==> Modify(Modify(s,i,itemi),j,itemj)==
                                                  Modify(Modify(s,j,itemj),i,itemi);
```

```
/*equality*/
        <<type Array>>Make(itemi) = Make(itemj)             == itemi = itemj;
        a=s == true, i=j == true, itemi = itemj ==>
                 Modify(a,i,itemi) = Modify(s,j,itemj)      == true;
/* */
        Extract(a,i) = Extract(s,i) == false ==>  a = s     == false;)));
endvalue type Array;
/*
```

14.8.2 Usage

An array is used to define one sort which is indexed by another. For example:

```
        value type indexbychar inherits Array< Character, Integer >
        endvalue type indexbychar;
```

defines an array containing integers and indexed by characters.

Arrays are usually used in combination with the shorthand forms of `Make`, `Modify`, and `Extract` defined in clause 12.2.3 of [ITU-T Z.101] and clause 12.3.3.1 of [ITU-T Z.101]. For example:

```
        dcl charvalue indexbychar;
        task charvalue := (. 12 .);
        task charvalue('A') := charvalue('B')-1;
```

14.9 Vector

14.9.1 Definition

```
*/
value type Vector < type Itemsort; synonym MaxIndex >
     inherits Array< Indexsort, Itemsort >;
syntype Indexsort = Integer constants 1:MaxIndex endsyntype;
endvalue type Vector;
/*
```

14.10 Powerset sort

14.10.1 Definition

```
*/
value type Powerset < type Itemsort >;
  operators
     empty                                      -> this Powerset;
     "in"   ( Itemsort, this Powerset       ) -> Boolean;       /* is member of          */
     incl   ( Itemsort, this Powerset       ) -> this Powerset; /* include item in set   */
     del    ( Itemsort, this Powerset       ) -> this Powerset; /* delete item from set  */
     "<"    ( this Powerset, this Powerset  ) -> Boolean;       /* is proper subset of   */
     ">"    ( this Powerset, this Powerset  ) -> Boolean;       /* is proper superset of */
     "<="   ( this Powerset, this Powerset  ) -> Boolean;       /* is subset of          */
     ">="   ( this Powerset, this Powerset  ) -> Boolean;       /* is superset of        */
     "and"  ( this Powerset, this Powerset  ) -> this Powerset; /* intersection of sets  */
     "or"   ( this Powerset, this Powerset  ) -> this Powerset; /* union of sets         */
     length ( this Powerset                 ) -> Integer;
     take   ( this Powerset                 ) -> Itemsort raise Empty;
axioms
     for all i,j in Itemsort (
     for all p,ps,a,b,c in Powerset (
/* constructors are empty and incl */
/* equalities between constructor terms */
     incl(i,incl(j,p))                    == incl(j,incl(i,p));
     i = j ==> incl(i,incl(j,p))          == incl(i,p);
/* definition of "in" by applying it to all constructors */
     i in <<type Powerset>>empty          == false;
     i in incl(j,ps)                      == i=j or i in ps;
/* definition of del by applying it to all constructors */
     <<type Powerset>>del(i,empty)        == empty;
     i =  j ==> del(i,incl(j,ps))       == del(i,ps);
     i /= j ==> del(i,incl(j,ps))       == incl(j,del(i,ps));
/* definition of "<" by applying it to all constructors */
```

```
     a <  <<type Powerset>>empty          == false;
     <<type Powerset>>empty < incl(i,b)== true;
     incl(i,a) < b                      == i in b and del(i,a) < del(i,b);
/* definition of ">" by other operations */
     a > b == b < a;
/* definition of "=" by applying it to all constructors */
     empty = incl(i,ps)                 == false;
     incl(i,a) = b      == i in b and del(i,a) = del(i,b);
/* definition of "<=" and ">=" by other operations */
     a <= b                             == a < b or a = b;
     a >= b                             == a > b or a = b;
/* definition of "and" by applying it to all constructors */
     empty and b                        == empty;
     i in b     ==> incl(i,a) and b     == incl(i,a and b);
     not(i in b) ==> incl(i,a) and b    == a and b;
/* definition of "or" by applying it to all constructors */
     empty or b                         == b;
     incl(i,a) or b                     == incl(i,a or b);
/* definition of length */
     length(<<type Powerset>>empty)     == 0;
     i in ps ==> length(ps)             == 1 + length(del(i, ps));
/* definition of take */
     take(empty)                        == raise Empty;
     i in ps ==> take(ps)               == i;
     ));
endvalue type Powerset;
/*
```

14.10.2 Usage

Powersets are used to represent mathematical sets. For example:

```
value type Boolset inherits Powerset< Boolean > endvalue type Boolset;
```

is used for a variable which is allowed to be empty or contain (`true`), (`false`) or (`true, false`).

14.11 Duration sort

14.11.1 Definition

```
*/
value type Duration;
  literals unordered nameclass ('0':'9')+ or (('0':'9')* '.' ('0':'9')+);
  operators
     protected duration ( Real            ) -> this Duration;
     "+"   ( this Duration, this Duration ) -> this Duration;
     "-"   ( this Duration                ) -> this Duration;
     "-"   ( this Duration, this Duration ) -> this Duration;
     ">"   ( this Duration, this Duration ) -> Boolean;
     "<"   ( this Duration, this Duration ) -> Boolean;
     ">="  ( this Duration, this Duration ) -> Boolean;
     "<="  ( this Duration, this Duration ) -> Boolean;
     "*"   ( this Duration, Real          ) -> this Duration;
     "*"   ( Real, this Duration          ) -> this Duration;
     "/"   ( this Duration, Real          ) -> Duration;
axioms noequality
/* constructor is duration(Real)*/
     for all a, b in Real nameclass (
     for all d, e in Duration nameclass (
/* definition of "+" by applying it to all constructors */
     duration(a) + duration(b)    == duration(a + b);
/* */
/* definition of unary "-" by applying it to all constructors */
     - duration(a)             == duration(-a);
/* */
/* definition of binary "-" by other operations */
     d - e                        == d + (-e);
/* */
/* definition of "equal", ">", "<", ">=", and "<=" by applying it to all constructors */
     equal(duration(a), duration(b)) == a = b;
     duration(a) > duration(b)    == a > b;
     duration(a) < duration(b)    == a < b;
```

```
     duration(a) >= duration(b) == a >= b;
     duration(a) <= duration(b) == a <= b;
/* */
/* definition of "*" by applying it to all constructors */
     duration(a) * b              == duration(a * b);
     a * d                        == d * a;
/* */
/* definition of "/" by applying it to all constructors */
     duration(a) / b           == duration(a / b);
/* */
     spelling(d) == spelling(a) ==>
                           d   == duration(a);
     ));
endvalue type Duration;
/*
```

14.11.2 Usage

The duration sort is used for the value to be added to the current time to set timers. The literals of the sort duration are the same as the literals for the real sort. The meaning of one unit of duration will depend on the system being defined.

Duration values are allowed be multiplied and divided by real values. Unless otherwise specified, the time unit is 1 second.

14.12 Time sort

14.12.1 Definition

```
*/
value type Time;
  literals unordered nameclass ('0':'9')+ or (('0':'9')* '.' ('0':'9')+);
  operators
     protected time  ( Duration   ) -> this Time;
     "<"   ( this Time, this Time  ) -> Boolean;
     "<="  ( this Time, this Time  ) -> Boolean;
     ">"   ( this Time, this Time  ) -> Boolean;
     ">="  ( this Time, this Time  ) -> Boolean;
     "+"   ( this Time, Duration) -> this Time;
     "+"   ( Duration, this Time) -> this Time;
     "-"   ( this Time, Duration) -> this Time;
     "-"   ( this Time, this Time  ) -> Duration;
axioms noequality
/* constructor is time */
     for all t, u in Time nameclass (
     for all a, b in Duration nameclass (
/* definition of ">", "equal" by applying it to all constructors */
     time(a) > time(b)          == a > b;
     equal(time(a), time(b))    == a = b;
/* */
/* definition of "<", "<=", ">=" by other operations */
     t < u                      == u > t;
     t <= u                     == (t < u) or (t = u);
     t >= u                     == (t > u) or (t = u);
/* */
/* definition of "+" by applying it to all constructors */
     time(a) + b                == time(a + b);
     a + t                      == t + a;
/* */
/* definition of "-" : Time, Duration by other operations */
     t - b                      == t + (-b);
/* */
/* definition of "-" : Time, Time by applying it to all constructors */
     time(a) - time(b)          == a - b;
/* */
     spelling(a) == spelling(t) ==>
                           a   == time(t);
     ));
endvalue type Time;
/*
```

14.12.2 Usage

The now expression returns a value of the time sort. It is allowed that time value has a duration added or subtracted from it to give another time value. A time value subtracted from another time value gives a duration. Time values are used to set the expiry time of timers.

The origin of time is system dependent. A unit of time is the amount of time represented by adding one duration unit to a time. Unless otherwise specified, the time unit is 1 second.

14.13 Bag sort

14.13.1 Definition

```
*/
value type Bag < type Itemsort >;
  operators
     empty                              -> this Bag;
     "in"   ( Itemsort, this Bag  ) -> Boolean;/* is member of  */
     incl   ( Itemsort, this Bag  ) -> this Bag;  /* include item in set     */
     del    ( Itemsort, this Bag  ) -> this Bag;  /* delete item from set    */
     "<"    ( this Bag, this Bag  ) -> Boolean;/* is proper subbag of       */
     ">"    ( this Bag, this Bag  ) -> Boolean;/* is proper superbag of     */
     "<="   ( this Bag, this Bag  ) -> Boolean;/* is subbag of  */
     ">="   ( this Bag, this Bag  ) -> Boolean;/* is superbag of            */
     "and"  ( this Bag, this Bag  ) -> this Bag;  /* intersection of bags    */
     "or"   ( this Bag, this Bag  ) -> this Bag;  /* union of bags           */
     length ( this Bag            ) -> Integer;
     take   ( this Bag            ) -> Itemsort raise Empty;
axioms
     for all i,j in Itemsort (
     for all p,ps,a,b,c in Bag (
/* constructors are empty and incl */
/* equalities between constructor terms */
     incl(i,incl(j,p))                  == incl(j,incl(i,p));
/* definition of "in" by applying it to all constructors */
     i in <<type Bag>>empty             == false;
     i in incl(j,ps)                    == i=j or i in ps;
/* definition of del by applying it to all constructors */
     <<type Bag>>del(i,empty)           == empty;
     i =  j ==> del(i,incl(j,ps))  == ps;
     i /= j ==> del(i,incl(j,ps))  == incl(j,del(i,ps));
/* definition of "<" by applying it to all constructors */
     a <  <<type Bag>>empty             == false;
     <<type Bag>>empty < incl(i,b)   == true;
     incl(i,a) < b                  == i in b and del(i,a) < del(i,b);
/* definition of ">" by other operations */
     a > b == b < a;
/* definition of "=" by applying it to all constructors */
     empty = incl(i,ps)             == false;
     incl(i,a) = b                  == i in b and del(i,a) = del(i,b);
/* definition of "<=" and ">=" by other operations */
     a <= b                             == a < b or a = b;
     a >= b                             == a > b or a = b;
/* definition of "and" by applying it to all constructors */
     empty and b                        == empty;
     i in b     ==> incl(i,a) and b == incl(i,a and b);
     not(i in b) ==> incl(i,a) and b == a and b;
/* definition of "or" by applying it to all constructors */
     empty or b                         == b;
     incl(i,a) or b                     == incl(i,a or b);
/* definition of length */
     length(<<type Bag>>empty)          == 0;
     i in ps    ==> length(ps)          == 1 + length(del(i, ps));
/* definition of take */
     take(empty)                        == raise Empty;
     i in ps    ==> take(ps)            == i; ));
endvalue type Bag;
/*
```

14.13.2 Usage

Bags are used to represent multi-sets. For example:

```
value type Boolset inherits Bag< Boolean > endvalue type Boolset;
```

is used for a variable which allowed to be empty or contain (`true`), (`false`), (`true, false`) (`true, true`), (`false, false`),... .

Bags are used to represent the `SET OF` construction of ASN.1.

14.14 ASN.1 Bit and Bitstring sorts

14.14.1 Definition

```
*/
value type Bit
     inherits Boolean ( 0 = false, 1 = true );
     adding;
  operators
     num ( this Bit  ) -> Integer;
     bit ( Integer) -> this Bit raise OutOfRange;
axioms
     <<type Bit>>num (0)           == 0;
     <<type Bit>>num (1)           == 1;
     <<type Bit>>bit (0)           == 0;
     <<type Bit>>bit (1)           == 1;
     for all i in Integer (
        i > 1 or i < 0 ==> bit (i)== raise OutOfRange;
        )
endvalue type Bit;
/* */
value type Bitstring
  operators
     bs in nameclass
     '''' ( (('0' or '1')*'''B') or ((('0':'9') or ('A':'F'))*'''H') )-> this Bitstring;
/*The following operators that are the same as String except Bitstring
     is indexed from zero*/
     mkstring   (Bit                                ) -> this Bitstring;
     Make       (Bit                                ) -> this Bitstring;
     length     ( this Bitstring                    ) -> Integer;
     first      ( this Bitstring                    ) -> Bit;
     last       ( this Bitstring                    ) -> Bit;
     "//"       ( this Bitstring, this Bitstring    ) -> this Bitstring;
     Extract    ( this Bitstring, Integer           ) -> Bit raise InvalidIndex;
     Modify     ( this Bitstring, Integer, Bit      ) -> this Bitstring;
     substring  ( this Bitstring, Integer, Integer) -> this Bitstring raise InvalidIndex;
     /* substring (s,i,j) gives a string of length j starting from the ith element */
     remove     ( this Bitstring, Integer, Integer) -> this Bitstring;
     /* remove (s,i,j) gives a string with a substring of length j starting from
        the ith element removed */
/*The following operators are specific to Bitstrings*/
     "not" ( this Bitstring                  ) ->  this Bitstring;
     "and" ( this Bitstring, this Bitstring  ) ->  this Bitstring;
     "or"  ( this Bitstring, this Bitstring  ) ->  this Bitstring;
     "xor" ( this Bitstring, this Bitstring  ) ->  this Bitstring;
     "=>"  ( this Bitstring, this Bitstring  ) ->  this Bitstring
     num   ( this Bitstring                  ) ->  Integer;
     bitstring ( Integer                     ) ->  this Bitstring raise OutOfRange;
     octet ( Integer                         ) ->  this Bitstring raise OutOfRange;
axioms
/* Bitstring starts at index 0 */
/* Definition of operators with the same names as String operators*/
     for all b in Bit ( /*b is bit in string*/
     for all s,s1,s2,s3 in Bitstring (
     for all i,j in Integer (
/* constructors are ''B, mkstring, and "//" */
/* equalities between constructor terms */
     s // ''B          == s;
     ''B// s           == s;
     (s1 // s2) // s3== s1 // (s2 // s3);
```

```
/* definition of length by applying it to all constructors */
     <<type Bitstring>>length(''B)                == 0;
     <<type Bitstring >>length(mkstring(b))       == 1;
     <<type Bitstring >>length(s1 // s2)          == length(s1) + length(s2);
     Make(s)                                      == mkstring(s);
/* definition of Extract by applying it to all constructors,
        with error cases handled separately */
     Extract(mkstring(b),0)                       == b;
     i < length(s1)       ==> Extract(s1 // s2,i) == Extract(s1,i);
     i >= length(s1)      ==> Extract(s1 // s2,i) == Extract(s2,i-length(s1));
     i<0 or i=>length(s)  ==> Extract(s,i)        == raise InvalidIndex;
/* definition of first and last by other operations */
     first(s)   == Extract(s,0);
     last(s)    == Extract(s,length(s)-1);
/* definition of substring(s,i,j) by induction on j,
        error cases handled separately */
     i>=0 and i < length(s)             ==>
          substring(s,i,0)  == ''B;
/* */
     i>=0 and j>0 and i+j<=length(s) ==>
          substring(s,i,j)  == substring(s,i,j-1) // mkstring(Extract(s,i+j));
/* */
     i<0 or j<0 or i+j>length(s)       ==>
          substring(s,i,j)  == raise InvalidIndex;
/* */
/* definition of Modify by other operations */
     Modify(s,i,b) == substring(s,0,i) // mkstring(b) // substring(s,i+1,length(s)-i);
/* definition of remove */
     remove(s,i,j)   == substring(s,0,i) // substring(s,i+j,length(s)-i-j);
     )));
/*end of definition of string operators indexed from zero*/
/* */
/* Definition of ''H and 'x'H in terms of ''B, 'xxxx'B for Bitstring*/
     <<type Bitstring>>''H  == ''B;
     <<type Bitstring>>'0'H == '0000'B;
     <<type Bitstring>>'1'H == '0001'B;
     <<type Bitstring>>'2'H == '0010'B;
     <<type Bitstring>>'3'H == '0011'B;
     <<type Bitstring>>'4'H == '0100'B;
     <<type Bitstring>>'5'H == '0101'B;
     <<type Bitstring>>'6'H == '0110'B;
     <<type Bitstring>>'7'H == '0111'B;
     <<type Bitstring>>'8'H == '1000'B;
     <<type Bitstring>>'9'H == '1001'B;
     <<type Bitstring>>'A'H == '1010'B;
     <<type Bitstring>>'B'H == '1011'B;
     <<type Bitstring>>'C'H == '1100'B;
     <<type Bitstring>>'D'H == '1101'B;
     <<type Bitstring>>'E'H == '1110'B;
     <<type Bitstring>>'F'H == '1111'B;
/* */
/* Definition of Bitstring specific operators*/
     <<type Bitstring>>mkstring(0)          == '0'B;
     <<type Bitstring>>mkstring(1)          == '1'B;
/* */
     for all s, s1, s2, s3 in Bitstring (
     s  =  s                                == true;
     s1 =  s2                               == s2 = s1;
     s1 /= s2                               == not ( s1 = s2 );
     s1 =  s2 == true            ==> s1     == s2;
     ((s1 = s2) and (s2 = s3))  ==> s1 = s3  == true;
     ((s1 = s2) and (s2 /= s3)) ==> s1 = s3  == false;
/* */
     for all b, b1, b2 in Bit (
     not(''B)                       == ''B;
     not(mkstring(b) // s)          == mkstring( not(b) ) // not(s);
/* definition of or */
/* The length of or-ing two strings is the maximal length of both strings */
     ''B or ''B                     == ''B;
     length(s) > 0 ==> ''B or s == mkstring(0) or s;
     s1 or s2                       == s2 or s1;
```

```
     (b1 or b2) // (s1 or s2)     ==(mkstring(b1) // s1) or (mkstring(b2) // s2);
/* */
/* definition of remaining operators based on "or" and "not" */
     s1 and s2                    == not (not s1 or not s2);
     s1 xor s2                    == (s1 or s2) and not(s1 and s2);
     s1 => s2                     == not (s1 and s2);
     ));
/* */
/*Definition of 'xxxxx'B literals */
     for all s in Bitstring (
     for all b in Bit (
     for all i in Integer (
     <<type Bitstring>>num (''B)            == 0;
     <<type Bitstring>>bitstring (0)      == '0'B;
     <<type Bitstring>>bitstring (1)      == '1'B;
     num (s // mkstring (b))              == num (b) + 2 * num (s);
     i > 1               ==> bitstring (i) == bitstring (i / 2) // bitstring (i mod 2);
     i >= 0 and i <= 255  ==> octet (i)   == bitstring (i) or '00000000'B;
     i < 0               ==> bitstring (i) == raise OutOfRange;
     i < 0 or i > 255   ==> octet (i)    == raise OutOfRange;
     )))
/*Definition of 'xxxxx'H literals */
     for all b1,b2,b3,h1,h2,h3 in bs nameclass (
     for all bs1, bs2, bs3, hs1, hs2, hs3 in Charstring (
     spelling(b1) = '''' // bs1 // '''B',
     spelling(b2) = '''' // bs2 // '''B',
     bs1 /= bs2                           ==> b1 = b2    == false;
/* */
     spelling(h1) = '''' // hs1 // '''H',
     spelling(h2) = '''' // hs2 // '''H',
     hs1 /= hs2 ==> h1 = h2 == false;
     spelling(b1) = '''' // bs1 // '''B',
     spelling(b2) = '''' // bs2 // '''B',
     spelling(b3) = '''' // bs1 // bs2 // '''B',
     spelling(h1) = '''' // hs1 // '''H',
     spelling(h2) = '''' // hs2 // '''H',
     spelling(h3) = '''' // hs1 // hs2 // '''H',
     length(bs1) = 4,
     length(hs1) = 1,
     length(hs2) > 0,
     length(bs2) = 4 * length(hs2),
     h1 = b1                              ==> h3 = b3    == h2 = b2;
/* */
/* connection to the String generator */
     for all b in Bit literals (
     spelling(b1)  = '''' // bs1 // bs2 // '''B',
     spelling(b2)  = '''' // bs2 // '''B',
     spelling(b)   = bs1                  ==> b1         == mkstring(b) // b2;
     )));
endvalue type Bitstring;
/*
```

14.15 ASN.1 Octet and Octetstring sorts

14.15.1 Definition

```
*/
syntype Octet = Bitstring size (8);
endsyntype Octet;
/* */
value type Octetstring
     inherits String < Octet > ( ''B = emptystring )
     adding
     operators
     os in nameclass
          '''' ( ((('0' or '1')8)*'''B') or ((('0':'9') or ('A':'F'))2)*'''H') )
           -> this Octetstring;
     bitstring   ( this Octetstring ) -> Bitstring;
     octetstring ( Bitstring ) -> this Octetstring;
```

```
axioms
     for all b,b1,b2 in Bitstring (
     for all s in Octetstring (
     for all o in Octet(
     <<type Octetstring>> bitstring(''B)            == ''B;
     <<type Octetstring>> octetstring(''B)          == ''B;
     bitstring( mkstring(o) // s )                  == o // bitstring(s);
/* */
     length(b1) > 0,
     length(b1) < 8,
     b2 == b1 or '00000000'B ==> octetstring(b1)  == mkstring(b2);
/* */
     b == b1 // b2,
     length(b1) == 8           ==> octetstring(b)== mkstring(b1) // octetstring(b2);
     )));
/* */
     for all b1, b2 in Bitstring (
     for all o1, o2 in os nameclass (
     spelling( o1 ) = spelling( b1 ),
     spelling( o2 ) = spelling( b2 ) ==> o1 = o2 == b1 = b2
     ));
endvalue type Octetstring;
/*
```

14.16 Pid sort

14.16.1 Definition

The data type `Pid` is the basis of all interface types.

```
*/
interface Pid { literals Null }
/*
```

14.17 Encoding sort

The following enumerated data type introduces the literals to support the encoding of data according to the standardized set of encoding rules. The set of literals is allowed to be extended by specializing `Encoding`.

```
*/
value type Encoding
     { literals text, BER, CER, DER, APER, UPER, CAPER, CUPER, BXER, CXER, EXER }
/*
```

which is used to denote the required set of encoding rules as follows:

`text` for the set of text encoding rules defined in this Recommendation and produces a `Charstring`;

`BER` for the set of Basic Encoding Rules of ASN.1 (see [ITU-T X.690]) and produces an `Octetstring`;

`CER` for the set of Canonical Encoding Rules of ASN.1 (see [ITU-T X.690]) and produces an `Octetstring`;

`DER` for the set of Distinguished Encoding Rules of ASN.1 (see [ITU-T X.690]) and produces an `Octetstring`;

`APER` for the basic ALIGNED variant of the Packed Encoding Rules of ASN.1 (see [ITU-T X.691]) and produces an `Octetstring`;

`UPER` for the basic UNALIGNED variant of the Packed Encoding Rules of ASN.1 (see [ITU-T X.691]) and produces a `Bitstring`;

`CAPER` for the ALIGNED variant of Canonical Packed Encoding Rules (CANONICAL-PER ALIGNED) of ASN.1 (see [ITU-T X.691]) and produces an `Octetstring`;

`CUPER` for the UNALIGNED variant of Canonical Packed Encoding Rules (CANONICAL-PER ALIGNED) of ASN.1 (see [ITU-T X.691]) and produces a `Bitstring`;

`BXER` for the Basic XML Encoding Rules (BASIC-XER) of ASN.1 (see [ITU-T X.693]) and produces a `Charstring`;

`CXER` for the Canonical XML Encoding Rules (CANONICAL-XER) of ASN.1 (see [ITU-T X.693]) and produces a `Charstring`;

`EXER` for the Extended XML Encoding Rules (EXTENDED-XER) of ASN.1 and (see [ITU-T X.693]) produces a `Charstring`.

The synonym PER is added to the **package** Predefined as an alternative for APER as follows:

```
*/
synonym PER Encoding = APER;
/*
```

The operator `last` of the data type `Encoding` is redefined to an unknown name, and therefore is inaccessible, so that an application or implementation is able to extend the data type `Encoding` without any change where only the above rules are used. It is possible to specialize the data type `Encoding` by adding additional literals.

14.18 Support for ASN.1 character, symbol string and NULL types

The following definitions are specifically to support character, string types and the NULL type of ASN.1. They have the same names in SDL-2010 and ASN.1.

```
*/
syntype NumericChar = Character constants
' ', '0', '1', '2', '3', '4',  '5', '6',
'7', '8', '9' endsyntype;
/* */

/*      NumericString sort      */
/*      Definition      */
value type NumericString
     inherits String < NumericChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
        '''' ( ('0':'9') or '''''' or (' ') )+ '''' -> this NumericString;
/*   character strings of any length of any characters from a space ' ' to a '9'      */
axioms
     for all c in NumericChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)              ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1             ==> cs  == cs1 // cs2;
     ));
endvalue type NumericString;
/* */
syntype PrintableChar = Character constants
' ', '0', '1', '2', '3', '4',  '5', '6',
'7', '8', '9', 'A', 'B', 'C',  'D', 'E',
'F', 'G', 'H', 'I', 'J', 'K',  'L', 'M',
'N', 'O', 'P', 'Q', 'R', 'S',  'T', 'U',
'V', 'W', 'X', 'Y', 'Z', 'a',  'b', 'c',
'd', 'e', 'f', 'g', 'h', 'i',  'j', 'k',
'l', 'm', 'n', 'o', 'p', 'q',  'r', 's',
't', 'u', 'v', 'w', 'x', 'y',  'z', '''',
'(', ')', '+', ',', '-', '.',  '/', ':',
'=', '?'
constants;
/* */

/*   PrintableString sort    */
/*   Definition      */
value type PrintableString
     inherits String < PrintableChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
             '''' ( (' ':'&') or '''''' or ('(': '?') )+ '''' -> this PrintableString;
```

```
/*   character strings of any length of any characters from a space ' ' to a '?'      */
axioms
     for all c in PrintableChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type PrintableString;
/* */
syntype TeletexChar = Character constants
/* TeletexString characters as in [ITU-T X.680] clause 41 Table 8 */
endsyntype;
/* */

/*   TeletexString sort */
/*   Definition   */
value type TeletexString
     inherits String < TeletexChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/*   characters specified in [ITU-T X.680] clause 41 Table 8 for TeletexString  */
         -> this TeletexString;
axioms
     for all c in TeletexChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type TeletexString;

syntype VideotexChar = Character constants
/* VideotexString characters as in [ITU-T X.680] clause 41 Table 8 */
endsyntype;
/* */

/*   VideotexString sort    */
/*   Definition     */
value type VideotexString
     inherits String < VideotexChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/*   VideotexString characters as in [ITU-T X.680] clause 41 Table 8 */
         -> this VideotexString;
axioms
     for all c in VideotexChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type VideotexString;

syntype IA5Char = Character endsyntype;

syntype IA5String = Charstring endsyntype;

value type UniversalChar
     literals /* see [ITU-T X.680] clause 41.6 */
operators
     uchr   ( Integer ) -> this UniversalChar;
endvalue type;
/* */

/*   UniversalCharString sort   */
/*   Definition     */
value type UniversalCharString
```

```
     inherits String < UniversalChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/*   see [ITU-T X.680] clause 41.6 */
        -> this UniversalCharString;
axioms
     for all c in UniversalChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type UniversalCharString;
/* */

/*   UTF8String sort */
syntype UTF8String = UniversalCharString endsyntype;
/* */

value type GeneralChar
     literals /* All G and all C sets + SPACE + DELETE
        See [ITU-T X.680] clause 41 Table 8 */
operators
     gchr   ( Integer ) -> this GeneralChar;
endvalue type;

/*   GeneralCharString sort  */
/*   Definition      */
value type GeneralCharString
     inherits String < GeneralChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/*   All G and all C sets + SPACE + DELETE
        See [ITU-T X.680] clause 41 Table 8 */
        -> this GeneralCharString;
axioms
     for all c in GeneralChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type GeneralCharString;
/* */
syntype GraphicChar = GeneralChar constants
/*   All G sets +SPACE as specified in [ITU-T X.680] clause 41 Table 8 */
endsyntype;
/* */

/*   GraphicCharString sort  */
/*   Definition      */
value type GraphicCharString
     inherits String < GraphicChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/*   All G sets + SPACE as specified in [ITU-T X.680] clause 41 Table 8 */
        -> this GraphicCharString;
axioms
     for all c in GraphicChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type GraphicCharString;

syntype VisibleChar = Character constants
/* VisibleString characters specified in [ITU-T X.680] clause 41 Table 8 */
endsyntype;
```

```
/* */

/*   VisibleString sort */
/*   Definition   */
value type VisibleString
     inherits String < VisibleChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/*   VisibleString characters specified in [ITU-T X.680] clause 41 Table 8*/
        -> this VisibleString;
axioms
     for all c in VisibleChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/*   string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type VisibleString;

syntype BMPChar = UniversalChar constants
/* see [ITU-T X.680] clause 41.15 */
endsyntype;
/* */
/*   BMPCharString sort   */
/*    Definition     */
value type BMPCharString
     inherits String < BMPChar > ( '' = emptystring )
     adding
        operators ocs in nameclass
/* see [ITU-T X.680] clause 41.15 */
        -> this BMPCharString;
axioms
     for all c in BMPChar nameclass (
     for all cs, cs1, cs2 in ocs nameclass (
     spelling(cs) == spelling(c)             ==> cs == mkstring(c);
/* string 'A' is formed from character 'A' etc. */
     spelling(cs) == spelling(cs1) // spelling(cs2),
     length(spelling(cs2)) == 1              ==> cs  == cs1 // cs2;
     ));
endvalue type BMPCharString;
/* */

value type NULL
literals NULL
endvalue type;
/*
```

14.19 Predefined exceptions

```
*/
exception
     OutOfRange,          /* A range check has failed. */
     InvalidReference,    /* Null was used incorrectly. Wrong Pid for this signal. */
     NoMatchingAnswer,    /* No answer matched in a decision without else part. */
     UndefinedVariable,   /* A variable was used that is "undefined". */
     UndefinedField,      /* An undefined field of a choice or struct was accessed. */
     InvalidIndex,        /* A String or Array was accessed with an incorrect index. */
     InvalidSort,         /* Dynamic sort failure */
     InvalidCall,         /* Dynamic call failure */
     DivisionByZero;      /* An Integer or Real division by zero was attempted. */
     Empty;               /* No element could be returned. */
/* */
endpackage Predefined;
```

Annex A

Abstract data types and axioms

(This annex forms an integral part of this Recommendation.)

A.1 Introduction

Predefined data items in the Specification and Description Language are based on abstract data types, which are defined in terms of their abstract properties rather than in terms of some concrete implementation. Even though the definition of an abstract data type gives one possible way of implementing that data type, an implementation is not mandated to choose that way of implementing the abstract data type, as long as the same abstract behaviour is preserved.

The predefined data types, including the Boolean sort that defines properties for two literals true and false, are defined in this annex. The two Boolean *terms* true and false shall not be (directly or indirectly) defined to be equivalent. Every Boolean constant expression that is used outside data type definitions shall be interpreted as either true or false. If it is not possible to reduce such an expression to true or false, then the specification is incomplete and allows more than one interpretation of the data type.

A.2 Notation

For this purpose, this annex extends the concrete syntax of the Specification and Description Language by means of describing the abstract properties of the operations added by a data type definition. However, this additional syntax is used for explanation only and does not extend the syntax defined in the main text. A specification using the syntax defined in this annex is therefore not valid SDL-2010.

The abstract properties described here do not specify a specific representation of the predefined data. Instead, an interpretation shall conform to these properties. When an <expression> is interpreted, the evaluation of the expression produces a value (e.g., as the result of an <operation application>). Two expressions, E1 and E2, are equivalent if:

a) there is an <equation> E1 == E2; or

b) one of the equations derived from the given set of <quantified equations>s is E1 == E2; or

c) i) E1 is equivalent to EA; and

   ii) E2 is equivalent to EB; and

   iii) there is an equation EA == EB or an equation derived from the given set of quantified equations such that EA == EB; or

d) by substituting a sub-term of E1 by a term of the same class as the sub-term producing a term E1A, it is possible to show that E1A is in the same class as E2.

Otherwise, the two expressions are not equivalent.

Two expressions that are equivalent represent the same value.

Interpretation of expressions conforms to these properties if two equivalent expressions represent the same value, and two non-equivalent expressions represent different values.

A.2.1 Axioms

Axioms determine which terms represent the same value. From the axioms in a data type definition, the relationship between argument values and result values of operators is determined and hence meaning is given to the operators. Axioms are either given as Boolean axioms or in the form of algebraic equivalence equations.

An operation defined by <axiomatic operation definitions> is treated as a complete definition with respect to specialization. That is, when a data type defined by the package Predefined is specialized and an operation is redefined in the specialized type, all axioms mentioning the name of the operation are replaced by the corresponding definition in the specialized type.

Concrete grammar

<axiomatic operation definitions> ::=
axioms <axioms>

<axioms> ::=
<equation> { <end> <equation>}* [<end>]

<equation> ::=
<unquantified equation>
| <quantified equations>
| <conditional equation>
| <literal equation>
| <noequality>

<unquantified equation> ::=
<term> **==** <term>
| <Boolean axiom>

<term> ::=
<constant expression>
| <error term>

<quantified equations> ::=
<quantification> **(**<axioms> **)**

<quantification> ::=
for all <value name> { **,** <value name> }* **in** <sort>

NOTE – **for** is considered a keyword for the purpose of this Annex.

This annex changes <operations> (see clause 12.1.1) as described below.

<operations> ::=
<operation signatures>
{ <operation definitions> | <axiomatic operation definitions> }

<axiomatic operation definitions> are only permitted for the description of the behaviour of operators.

An <identifier> which is an unqualified name appearing in a <term> is one of the following:

a) an <operation identifier>;

b) <literal identifier>;

c) a <value identifier> if there is a definition of that name in a <quantification> of <quantified equations> enclosing the <term>, which then shall have a suitable sort for the context.

Semantics

A ground term is a term that does not contain any value identifiers. A ground term represents a particular, known value. For each value in a sort there exists at least one ground term which represents that value.

Each equation is a statement about the algebraic equivalence of terms. The left hand side term and right hand side term are stated to be equivalent so that where one term appears, it is allowed to substitute the other term. When a value identifier appears in an equation, simultaneous substitution is allowed in that equation of the same term for every occurrence of the value identifier. For this substitution, the term is any ground term of the same sort as the value identifier.

Value identifiers are introduced by the value names in quantified equations. A value identifier is used to represent any data values belonging to the sort of the quantification. An equation will hold if the same value is simultaneously substituted for every occurrence of the value identifier in the equation regardless of the value chosen for the substitution.

In general, there is no need or reason to distinguish between a ground term and the result of the ground term. For example, the ground term for the unity Integer element is (normally) written "1". Usually there are several ground terms which denote the same data item, e.g., the Integer ground terms "0+1", "3–2" and "(7+5)/12", and it is usual to consider a simple form of the ground term (in this case "1") as denoting the data item.

A value name is always introduced by quantified equations, and the corresponding value has a value identifier, which is the value name qualified by the sort identifier of the enclosing quantified equations. For example:

```
for all z in X (for all z in X ... )
```

introduces only one value identifier named z of sort X.

In the concrete syntax of axioms, it is not allowed to specify a qualifier for value identifiers.

Each value identifier introduced by quantified equations has a sort, which is the sort identified in the quantified equations by the <sort>.

A term has a sort, which is the sort of the value identifier or the result sort of the (literal) operator.

Unless it is possible to deduce from the equations that two terms denote the same value, each term denotes a different value.

A.2.2 Conditional equations

A conditional equation allows the specification of equations that only hold when certain restrictions hold. The restrictions are written in the form of simple equations.

Concrete grammar

<conditional equation> ::=
<restriction> { , <restriction> }* ==> <restricted equation>

<restricted equation> ::=
<unquantified equation>

<restriction> ::=
<unquantified equation>

Semantics

A conditional equation defines that terms denote the same data item only when any value identifier in the restricted equation denotes a data item, which it is possible to show from other equations that satisfy the restrictions.

The semantics of a set of equations for a data type that includes conditional equations is derived as follows.

a) Quantification is removed by generating every possible ground term equation that is derivable from the quantified equations. As this is applied to both explicit and implicit quantification, a set of unquantified equations in ground terms only is generated.

b) Let a conditional equation be called a provable conditional equation if all the restrictions (in ground terms only) are able to be proved to hold from unquantified equations that are not restricted equations. If there exists a provable conditional equation, it is replaced by the restricted equation of the provable conditional equation.

c) If there are conditional equations remaining in the set of equations and none of these conditional equations are a provable conditional equation, then these conditional equations are deleted; otherwise, return to step b).

d) The remaining set of unquantified equations defines the semantics of the data type.

A.2.3 Equality

Concrete grammar

<noequality> ::=
noequality

Model

Any <data type definition> introducing some sort named S has the following implied <operation signature> in its <operator list>, unless <noequality> is present in the <axioms>:

equal (S, S) -> Boolean;

where Boolean is the predefined Boolean sort.

Any <data type definition> introducing a sort named S such that it contains only <axiomatic operation definitions> in <operator list> has an implied equation set:

```
for all a,b,c in S (
   equal(a, a) == true;
   equal(a, b) == equal(b, a);
   equal(a, b) and equal(b, c)  ==> equal(a, c) == true;
   equal(a, b) == true          ==> a == b;)
```

and an implied <literal equation>:

```
for all L1,L2 in S literals (
   spelling(L1) /= spelling(L2) ==> L1 = L2 == false;)
```

A.2.4 Boolean axioms

Concrete grammar

<Boolean axiom> ::=
<Boolean term>

Semantics

A Boolean axiom is a statement of truth that holds under all conditions for the data type being defined.

Model

An axiom of the form:

<Boolean term>;

is derived syntax for the concrete syntax equation:

<Boolean term> == << **package** Predefined/**type** Boolean >> true;

A.2.5 Conditional term

Semantics

An equation containing a conditional term is semantically equivalent to a set of equations where all the quantified value identifiers in the Boolean term have been eliminated. It is possible to form this set of equations by simultaneously substituting, throughout the conditional term equation, each <value identifier> in the <conditional expression> by each ground term of the appropriate sort. In this set of equations, the <conditional expression> will always have been replaced by a Boolean ground term. In the following, this set of equations is referred to as the expanded ground set.

A conditional term equation is equivalent to the equation set that contains:

a) for every *equation* in the expanded ground set for which the <conditional expression> is equivalent to true, that *equation* from the expanded ground set with the <conditional expression> replaced by the (ground) <consequence expression>; and

b) for every *equation* in the expanded ground set for which the <conditional expression> is equivalent to false, that *equation* from the expanded ground set with the <conditional expression> replaced by the (ground) <alternative expression>.

Note that in the special case of an equation of the form:

ex1 == **if** a **then** b **else** c **fi**;

this is equivalent to the pair of conditional equations:

a == true ==> ex1 == b;
a == false ==> ex1 == c;

A.2.6 Error term

Errors are used to allow the properties of a data type to be fully defined even for cases when it is not possible to give a specific meaning to the result of an operator.

Concrete grammar

<error term> ::=
raise <exception name>

An <error term> shall not be used as part of a <restriction>.

It shall not be possible to derive from *equations* that a <literal identifier> is equal to <error term>.

Semantics

A term is allowed to be an <error term>, so that it is possible to specify the circumstances under which an operator produces an error. If these circumstances arise during interpretation, then the exception with <exception name> is raised.

A.2.7 Unordered literals

Concrete grammar

<unordered> ::=
unordered

This annex changes the concrete syntax for the literal list type constructor (see clause 12.1.6.17) as follows:

<literal list> ::=
[<visibility>] **literals** [<unordered>]
<literal signature> { **,** <literal signature> }* <end>

Model

If <unordered> is used, the *Model* in clause 12.1.6.17 is not applied. Consequentially, the ordering operations "<", ">","<=",">=", first, last, pred, succ, and num are not implicitly defined for this data type.

A.2.8 Literal equations

Concrete grammar

<literal equation> ::=
<literal quantification>
(<equation> { <end> <equation> }* [<end>] **)**

<literal quantification> ::=
for all <value name> { **,** <value name> }* **in** <sort> **literals**
| **for all** <value name> { **,** <value name> }* **in** { <sort> | <value identifier> } **nameclass**

Semantics

Literal mapping is a shorthand for defining a large (possibly infinite) number of axioms ranging over all the literals of a sort or all the names in a name class. The literal mapping allows the literals for a sort to be mapped onto the values of the sort.

<spelling term> is used in literal quantifications to refer to the character string that contains the spelling of the literal. This mechanism allows the `Charstring` operators to be used to define literal equations.

Model

A <literal equation> is shorthand for a set of <axioms>. In each of the <equation>s contained in a <literal equation>, the <value identifier>s defined by the <value name> in the <literal quantification> are replaced. In each derived <equation>, each occurrence of the same <value identifier> is replaced by the same <literal identifier> of the <sort> of the <literal quantification> (if literals was used) or by the same <literal identifier> of the nameclass referred to (if nameclass was used). The derived set of <axioms> contains all possible <equation>s derivable in this way.

The derived <axioms> for <literal equation>s are added to <axioms> (if any) defined after the keyword **axioms**.

If a <literal quantification> contains one or more <spelling term>s, then there is replacement of the <spelling term>s with `Charstring` literals (see clause 14.4). The `Charstring` is used to replace the <value identifier> after the <literal equation> containing the <spelling term> and is expanded as defined below, using <value identifier> in place of <operation name>.

NOTE – Literal equations do not affect nullary operators defined in <operation signature>s.

A.2.9 Name class

A name class is shorthand for a (possibly infinite) set of literal names or operator names defined by a regular expression.

A <name class literal> is an alternative way of specifying a <literal name>. A <name class operation> is an alternative way of specifying an <operation name> of a nullary operation.

Concrete grammar

```
<name class literal> ::=
                        nameclass <regular expression>

<name class operation> ::=
                        <operation name> in nameclass <regular expression>

<regular expression> ::=
                        <partial regular expression> { [or ] <partial regular expression> }*

<partial regular expression> ::=
                        <regular element> [ <integer name> | <plus sign> | <asterisk> ]

<regular element> ::=
                        ( <regular expression> )
                    |   <character string>
                    |   <regular interval>

<regular interval> ::=
                        <character string> { <colon> | <range sign> } <character string>
```

The names formed by the <regular expression> shall satisfy the lexical rules for names or <character string>, <hex string>, or <bit string>.

The <character string>s in a <regular interval> shall both have a length of one, excluding the leading and trailing <apostrophe>s.

A <name class operation> is allowed only in an <operation signature>. An <operation signature> containing <name class operation> shall only occur in an <operator list> and shall not contain <arguments>.

When a name contained in the equivalent set of names of a <name class operation> occurs as the <operation name> in an <operation application>, it shall not have <actual parameters>.

The equivalent set of names of a name class is defined as the set of names that satisfy the syntax specified by the <regular expression>. The equivalent sets of names for the <regular expression>s contained in a <data type definition> shall not overlap.

Model

A <name class literal> is equivalent to this set of names in the abstract syntax. When a <name class operation> is used in an <operation signature>, a set of <operation signature>s is created by substituting each name in the equivalent set of names for the <name class operation> in the <operation signature>.

A <regular expression> which is a list of <partial regular expression>s without an **or** specifies that the names are formed from the characters defined by the first <partial regular expression> followed by the characters defined by the second <partial regular expression>.

When an **or** is specified between two <partial regular expression>s, then the names are formed from either the first or the second of these <partial regular expression>s. **or** is more tightly binding than simple sequencing.

If a <regular element> is followed by <Natural literal name>, the <partial regular expression> is equivalent to the <regular element> being repeated the number of times specified by the <Natural literal name>.

If a <regular element> is followed by '*' the <partial regular expression> is equivalent to the <regular element> being repeated zero or more times.

If a <regular element> is followed by <plus sign> the <partial regular expression> is equivalent to the <regular element> being repeated one or more times.

A <regular element> which is a bracketed <regular expression> defines the character sequences defined by the <regular expression>.

A <regular element> which is a <character string> defines the character sequence given in the character string (omitting the quotes).

A <regular element> which is a <regular interval> defines all the characters specified by the <regular interval> as alternative character sequences. The characters defined by the <regular interval> are all the characters greater than or equal to the first character and less than or equal to the second character according to the definition of the Character sort (see clause A.2).

The names generated by a <name class literal> are defined first by length (a shorter literal comes before any longer ones) then in lexicographical order according to the ordering of the character sort. The characters are considered case sensitive.

NOTE – Examples exist in clause 14.

A.2.10 Name class mapping

A name class mapping is shorthand for defining a (possibly infinite) number of operation definitions ranging over all the names in a <name class operation>. The name class mapping allows behaviour to be defined for the operators and methods defined by a <name class operation>. A name class mapping takes place when an <operation name> that occurred in a <name class operation> within an <operation signature> of the enclosing <data type definition> is used in <operation definitions> or <operation diagram>s.

A spelling term in a name class mapping refers to the character string that contains the spelling of the name. This mechanism allows the `Charstring` operations to be used to define name class mappings.

Concrete grammar

<spelling term> ::=

spelling ({ <operation name> | <value identifier> })

The <value identifier> in a <spelling term> shall be a <value identifier> defined by a <literal quantification>.

A <spelling term> is legal concrete syntax only within an <operation definition> or <operation diagram>, if a name class mapping has taken place.

Model

A name class mapping is shorthand for a set of <operation definition>s or a set of <operation diagram>s. The set of <operation definition>s is derived from an <operation definition> by substituting each name in the equivalent set of names of the corresponding <name class operation> for each occurrence of <operation name> in the <operation definition>. The derived set of <operation definition>s contains all possible <operation definition>s that are able to be generated in this way. The same procedure is followed for deriving a set of <operation diagram>s.

The derived <operation definition>s and <operation diagram>s are considered legal even though a <string name> is not allowed as an <operation name> in the concrete syntax.

The derived <operation definition>s are added to <operation definitions> (if any) in the same <data type definition>. The derived <operation diagram>s are added to the list of diagrams where the original <operation definition> had occurred.

If an <operation definition> or <operation diagram> contains one or more <spelling term>s, each <spelling term> is replaced with a `Charstring` literal (see clause 14.4).

If, during the above transformation, the <operation name> in the <spelling term> had been replaced by an <operation name>, the <spelling term> is shorthand for a `Charstring` derived from the <operation name>. The `Charstring` contains the spelling of the <operation name>.

If, during the above transformation, the <operation name> in the <spelling term> had been replaced by a <string name>, the <spelling term> is shorthand for a `Charstring` derived from the <string name>. The `Charstring` contains the spelling of the <string name>.

Annex B

Specification of the set of text encoding rules

(This annex forms an integral part of this Recommendation.)

The data types are presented below with data types of the **package** `Predefined` in the order they occur, followed by other data types.

In this Annex characters that appear in encoded text are identified by the names (in uppercase letters) they are given in [ISO/IEC 10646]. For example: LATIN CAPITAL LETTER T identifies the character in [ISO/IEC 10646] used to encode the `Boolean` value `true`.

B.1 Boolean

`Boolean` values, False and True shall be encoded as LATIN CAPITAL LETTER F and LATIN CAPITAL LETTER T respectively.

Example

```
dcl Var_Boolean Boolean;
task Var_Boolean := true;
```

Generated `CharString`: T

B.2 Character

`Character` values shall be encoded as the actual Character with the exception of the `ESC` character, which is encoded as two ESCAPE characters. An undefined or missing Character value shall be encoded as an ESCAPE character followed by the NULL character.

NOTE – The Generated `CharString` is allowed to contain non-printing characters.

Example

```
dcl Var_Character Character;
task Var_Character := 'M';
```

Generated `CharString`: M

B.3 String

`String` has a parameter for the item sort and shall be encoded as a list of values of the item sort enclosed in a LEFT CURLY BRACKET and a RIGHT CURLY BRACKET separated by COMMA characters.

Example

```
value type IntString inherits String <Integer>;
dcl str IntString;

task str:= ''//mkstring(6)//mkstring(9)//mkstring(1948);
```

Generated `CharString`: {6,9,1948}

B.4 Charstring, IA5String, NumericString, PrintableString, VisibleString

Although `Charstring` is based on `String`, because it is a predefined data type and is commonly used, it is given a special encoding. `IA5String` and `Charstring` cover the same range of characters and have the same encoding. `NumericString`, `PrintableString` and `VisibleString` are subsets of `IA5String`.

These string types shall be encoded as a string of characters enclosed in APOSTROPHE (') characters without change except the <apostrophe> (') which shall be encoded as two APOSTROPHE (') characters.

NOTE 1 – Within a character string <comma>s, <left curly bracket>s and <right curly bracket>s are treated as normal characters.

NOTE 2 – The Generated `Charstring` therefore possibly contains non-printing characters including end of file or end of record characters.

Example

```
dcl Var_Charstring Charstring;

task Var_Charstring := 'Fred''s world;
```

Generated `Charstring`: 'Fred's world'

B.5 Integer

`Integer` shall be encoded as a decimal integer notation without leading DIGIT ZERO characters where negative numbers are immediately preceded by a HYPHEN-MINUS without leading DIGIT ZERO characters. Zero shall never be treated as a negative number.

Example

```
dcl i Integer;

task i := 2-7;
```

Generated `Charstring`: –5

B.6 Natural

`Natural` is a **syntype** of `Integer` and shall therefore have the same encoding as `Integer`.

B.7 Real

The value 0.0 shall be encoded as DIGIT ZERO followed by FULL STOP followed by DIGIT ZERO.

Any negative value shall be encoded as a HYPHEN-MINUS immediately followed by the encoding of the value negated (a positive Real value).

A positive Real value shall be encoded as a single digit in the range 1 to 9, followed by a FULL STOP, followed by at least one and up to 11 decimal fraction digits, followed by LATIN SMALL LETTER E 'e', followed by an exponent. The exponent is the Integer encoding of the exponent value: the power to the base 10 to apply to the first part of the number. A negative exponent is allowed.

It is allowed to omit trailing fractional zero digits. It is not possible to precisely encode some Real values, and it is possible that unequal Real values that have very small difference have the same encoding. In any case, whether a Real value is precisely correct within an application will depend on how Real has been implemented. For example, if the Real value 2.0/7.0 is actually stored as the ratio of two integers (2 and 7), this is absolutely precise, whereas a more conventional encoding could be 0.2857142857, which is only correct to 10 decimal places.

Examples

```
dcl p, q, r1, r2 Real := 2000.0, 7.0, 0, 0;

task r1:= p/q;
task r2:= q/p;
```

Generated `Charstring` for r1: 2.85714285714e2

Generated `Charstring` for r2: 3.5e-3

B.8 Array

`Array` has two parameters: an index sort and a component sort. Either sort in principle is any sort, but usually the index sort has a finite number of values.

Where the index sort has a finite number of ordered values an `Array` shall be encoded as a list of element encoding values, one for each element, separated by COMMA characters within a single pair of LEFT CURLY BRACKET and RIGHT CURLY BRACKET characters.

Example

```
value type ABC {literals A, B, C};
value type A1 inherits Array <ABC, Integer>;
dcl avalue, bvalue A1;

task avalue:= (. 3 .);
task bvalue[A]:= 3;
task bvalue[B]:= 5;
task bvalue[C]:= 7;
```

Generated `Charstring` for avalue:{3,3,3}

Generated `Charstring` for bvalue:{3,5,7}

It is possible the index is not ordered (that is, the "<" operator is not defined, or for example the index sort is a **structure** or **choice**). It is also possible the index sort does not have a finite number of values (that is, no possible number of values is infinite for example `Charstring` or `Real`). For unordered or infinite index sort, the `Array` shall be encoded as the most frequent value followed by pairs of values for each element. If two or more values occur with the highest frequency one of these is chosen on an arbitrary basis as the most frequent value. Each pair shall be between a LEFT CURLY BRACKET and a RIGHT CURLY BRACKET and separated by a COMMA: the first value is an index value and the second value is element value. No index values in the pairs shall be repeated.

Example

```
value type Dehashing inherits Array <Charstring, Charstring> := (.''.);
/* note that the default value is an empty string */
dcl hashtable Dehashing;

task htable ('ac'):= 'action';
task htable ('ab'):= 'ability';
task htable ('zzzz'):= 'end of document';
```

Generated `Charstring` for `htable`: {",{'ab','ability'},{'ac','action'},{'zzzz','end of document'}}

B.9 Vector

Vector is a special case of Array that always has index sort that is a subset of Natural with a range from 1 to a specified maximum value and any item sort. Vector shall therefore be encoded in the same way as an Array with a finite number of ordered index values: that is a list of element encoding values, one for each element, separated by COMMA characters within a single LEFT CURLY BRACKET and RIGHT CURLY BRACKET pair.

B.10 Powerset

A `Powerset` of a sort represents a mathematical set whose elements are all members of the sort. Where the sort has a finite number of ordered values, the `Powerset` of the sort shall be encoded as a bit string (a string of DIGIT ZERO and DIGIT 1 characters) enclosed in APOSTROPHE (') characters. Each bit position in this bit string represents if a value of the sort is present or absent in the set. A DIGIT 1 represents that the value is present and a DIGIT ZERO that the value is absent. The leftmost bit represents the smallest value of the element sort, and each other bit represents a value of the element sort larger than values for the bits to the left.

Example

```
value type Shortalpha {literals a,b,c,d,e,f,g,h,i,j,k,l,m,n,o,p,q,r,s,t,u,v};
/* note Shortalpha has 22 values */
value type Psa inherits Powerset<Shortalpha>;
dcl letters_used Psa;

task letters_used := (. h, c, u, r .);
```

Generated `Charstring`: '00100001000000000010010'

Where the sort does not have a finite number of ordered values, the `Powerset` of the sort shall be encoded as the encoding of each element value present separated by COMMA characters enclosed in a LEFT CURLY BRACKET and RIGHT CURLY BRACKET pair. The element values are allowed to be in any order.

Example

```
value type Pchrstr inherits Powerset<Charstring>;
dcl strings_used Pchrstr;

task strings_used := (. 'me', 'you', 'us', 'me', 'again', 'hey', 'you' .);
```

Generated `Charstring`: {'again','hey','you','me','us'}

B.11 Duration

`Duration` is used to denote 'a time interval'. Unless specified otherwise, the default value of the unit of `Duration` is one second.

Duration values shall be encoded as a pair of integers separated by a COMMA enclosed in a LEFT CURLY BRACKET and RIGHT CURLY BRACKET pair: the first integer gives the number of units and the second integer any fractional part in nano (10^{-9}) units. By default these are seconds and nano-seconds. A negative value is allowed, in which case the encoding of the negated value shall be used with a HYPHEN-MINUS inserted immediately before the first integer.

Example

```
dcl dvar Duration;

task dvar := -17.00000007;
```

Generated `Charstring`: {-17,700}

B.12 Time

`Time` is used to denote 'a point in time'. The value of a unit of `Time` shall be the same as the value of the unit of `Duration`. The origin of `Time` units is not specified by this Recommendation, but corresponds to the system clock being at zero: that is, NOW gives the value 0.0. It is allowed for `Time` values to be negative.

Time values shall be encoded in the same way as Duration.

Example

```
dcl tvar Time;

task tvar:= 17.0000017;
```

Generated `Charstring`: {17,17000}

B.13 Bag

`Bag` shall be encoded in the same way as a `Powerset` with an element sort that does not have a finite ordered set of values, but with each element value preceded by an integer followed by COLON. The integer is the number of times the element sort occurs in the `Bag` value.

Example

```
value type B1 inherits Bag <Integer>;

dcl Var_Bag B1;

task Var_Bag := (. 7, 4, 7 .);
```

Generated `Charstring`: {2:7,1:4}

B.14 Bit, Bitstring

`Bit` shall be encoded as a single DIGIT ZERO or DIGIT 1 representing a zero and one bit respectively.

Example

```
dcl Var_Bit Bit;

task Var_Bit := 1;
```

Generated `Charstring`: 1

`Bitstring` shall be encoded as sequence of bits represented by DIGIT ZERO and DIGIT 1 characters enclosed in APOSTROPHE (') characters.

Example

```
dcl Var_Bit Bit_String;

task Var_Bit := '01011'B;
```

Generated `Charstring`: '01011'

B.15 Octet, Octetstring

`Octet` shall be encoded as two characters corresponding to the hexadecimal notation of the `Octet`. The alphabetic characters shall be represented by lowercase letters (that is, LATIN SMALL LETER A to LATIN SMALL LETTER F).

Example

```
dcl oct Octet;

task oct := 62;
```

Generated `Charstring`: 3e

`Octetstring` is a sequence of `Octet` values and shall be encoded as a string of hexadecimal character pairs enclosed in APOSTROPHE (') characters. The alphabetic characters shall be represented by lowercase letters.

Example

```
dcl os Octetstring;

task os:= '12B32D'H;
```

Generated `Charstring`: '12b32d'

B.16 Pid, pid sorts

To allow flexibility between applications pid values shall be encoded as a CHOICE value corresponding to the choice definition:

```
{ choice
      0 ApplicationDefined;
      1 Integer;
      2 OctetString;
      3 BitString;
      4 Charstring;
      5 { struct
           identity Charstring;
           instance Natural;
        };
}
```

A value of a pid sort shall be encoded as the corresponding value of the `Pid` sort.

All pid values shall be encoded using the same choice field as defined above: that is, if one pid value in a model is encoded using the choice 1 Integer, all other pid values in the same model shall be encoded using choice 1 Integer. Otherwise, which of the above choices is selected is application dependent. The first choice allows an application defined encoding to be used. The `ApplicationDefined` sort is not defined by this Recommendation.

The same agent instance of a model shall always have the same pid value encoding. Two different agent instances shall always have different encoded pid values. Otherwise the derivation of the encoded pid values is not further defined by the Recommendation.

Example (using choice 5)

```
dcl agent_id Pid; /* in second instance of process IPS */

task agent_id := self;
```

Generated `Charstring`: {5, {IPS, 2}}

B.17 Null

The value Null (see [ITU-T Z.105]) shall be encoded as a DIGIT ZERO character.

Example

```
dcl Var_Null Null;

task Var_Null := Null;
```

Generated `Charstring`: 0

B.18 Enumerated (literal list)

An enumerated type defined by a literal list or as an ASN.1 ENUMERATED type shall be encoded as the Natural value given by the application of `num` operator of the data type to the literal representing the value to be encoded.

Example

```
value type Enum
{ literals e1, e2, e3;
};

dcl evar Enum;

task evar := e2;
```

Generated `Charstring`: 1

B.19 Structures

A structure data type value shall be encoded as a list of values for the fields within a LEFT CURLY BRACKET and RIGHT CURLY BRACKET pair separated by COMMA characters.

Example

```
value type Record { struct
f1 Integer;
f2 Charstring;
f3 Integer;};
dcl locat Record;
task locat:= (. 17, 'mid-field', 230125 .);
```

Generated `Charstring`: {17, 'mid-field', 230125}

B.20 Choice

A choice value shall be encoded as a pair of values separated by a COMMA enclosed in a LEFT CURLY BRACKET and RIGHT CURLY BRACKET pair. The first value is the selected field indicated by name of the choice. The second value is the `Charstring` for the field value according to the type of the field.

Example

```
value type C {choice
cs Charstring;
cb Boolean;
};

dcl ChoiceVar C;

task ChoiceVar.cb := true;
```

Generated `Charstring`: {cb,T}

B.21 Inherits and syntype

A **syntype** defines a new name for a data type with an optional constraint on the set values. The text encoding is therefore identical to that of the data type.

A data type that inherits from another data type has the same encoding as the parent data type, though in the case that additional literal or fields are added, these are added to the encoding.

Annex C

Language binding

(This annex forms an integral part of this Recommendation.)

The material in this Annex is the subject of further study, and a replacement is planned for the future.

This Annex defines the use of the syntax of an alternative language within the syntax rules <variable definition>, <data definition>, <statements> and <expression>, which are taken as the points of syntax variation. By default, an SDL-2010 specification is bound to the native syntax as defined in the main body of this Recommendation or other Recommendations for SDL-2010.

Each diagram or other definition that is allowed a <package use clause> is bound to a particular concrete syntax as defined in clause 7.2. Until further study has been completed and this Annex is replaced, the only allowed language <data binding> is to the **package** Predefined, which is the language binding for the native SDL-2010 concrete syntax.

Though the syntax within syntax variation elements (<variable definition>, <data definition>, <statements>, <expression>) looks like another language (C, C++, Java or some other language), the semantics are defined by SDL-2010. The binding to the SDL-2010 abstract grammar is permitted to invoke complex transformations to achieve the mapping, and constraints are allowed to exclude constructions permitted by the language concrete syntax that are not able to be reasonably mapped. A model that includes constructs that do not map to SDL-2010 does not conform to the SDL-2010 language. A tool that handles constructs that do not map to SDL-2010 abstract grammar and semantics shall provide an indication if such constructs are used.

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