

Security Level:

Decentralized Network Resource Management and Trust Model

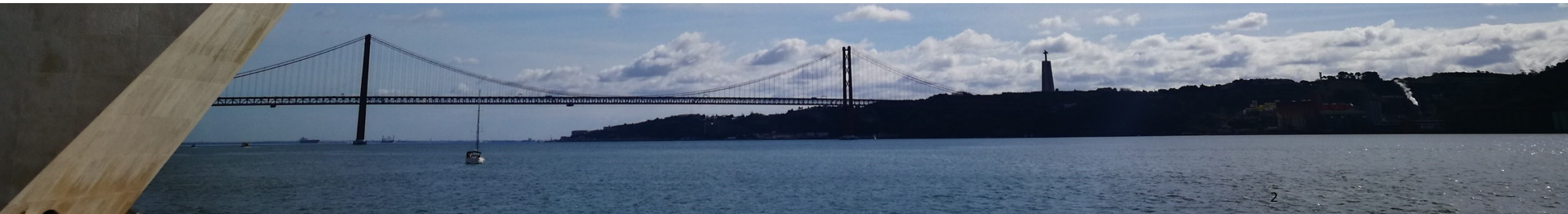
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Are we opening the Age of New Discovery.....



Network Services Rely on Trust Infrastructures

■ Infrastructure:

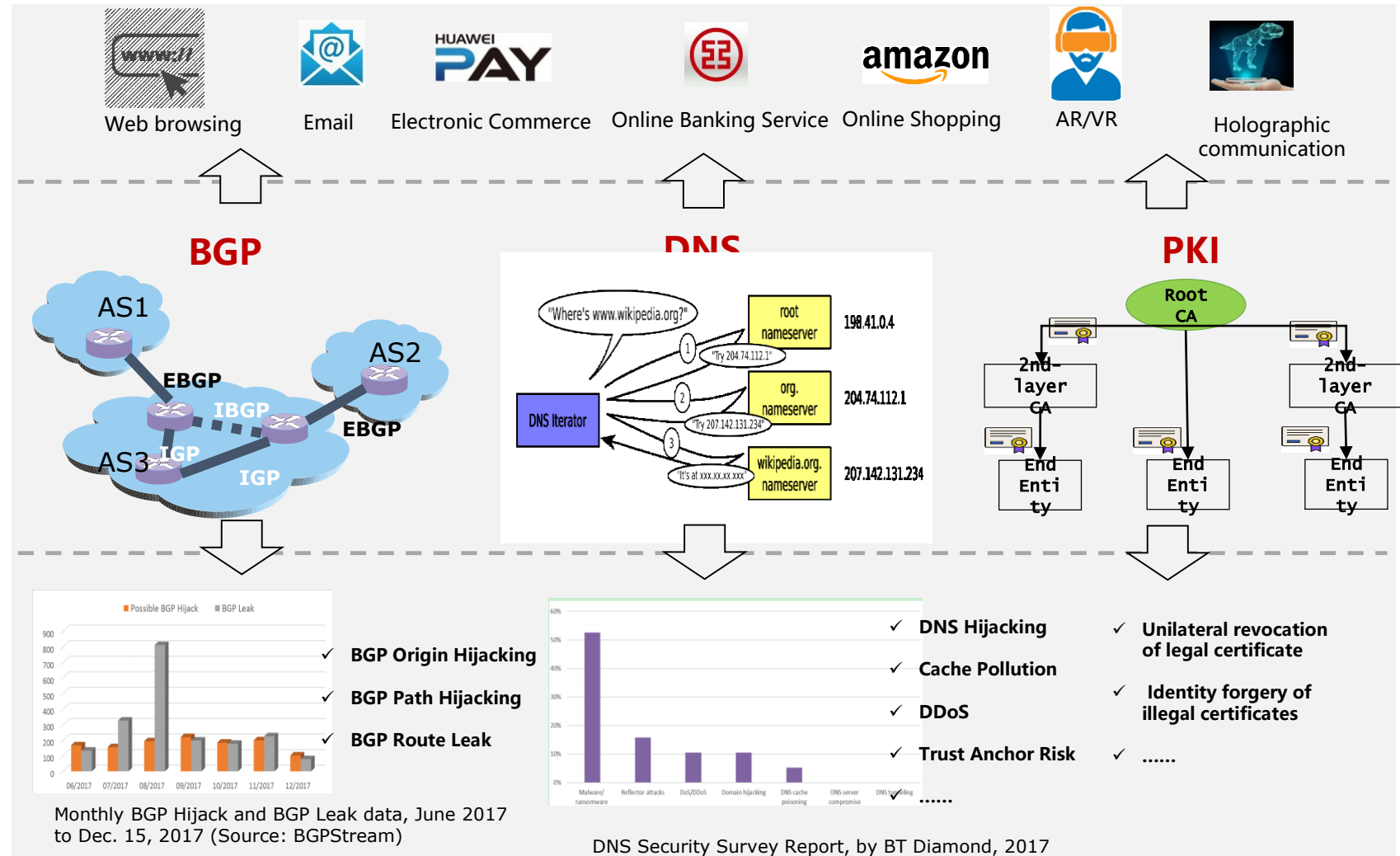
- Inter-domain Routing System (BGP)
- Name Resolution System
- Public Key Certificate System (PKI).

■ **Almost all network services rely on these infrastructure to ensure connectivity, service availability and credibility.**

■ **The current infrastructure lacks a solid, secure and credible foundation.**

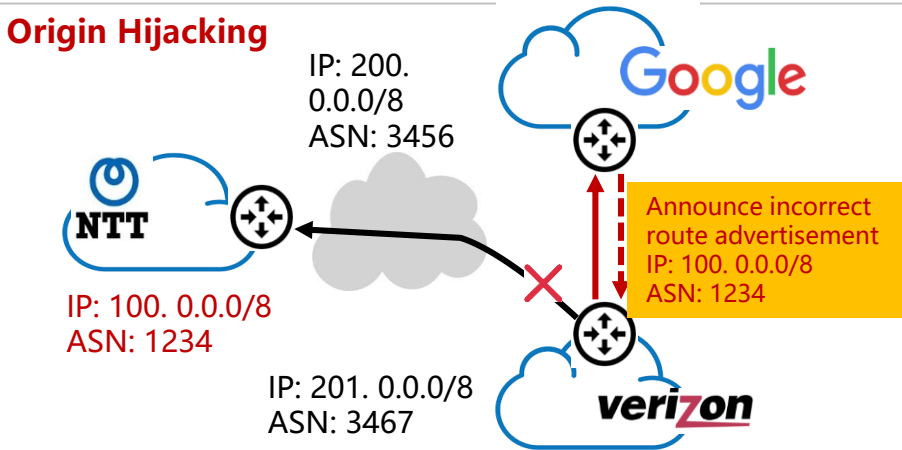
› BGP and DNS were not designed with any security and credibility at the beginning, so naturally lacked security capabilities.

› PKI relies on trust anchors for endorsement



- BGP lacks the ability to verify the validity of announcement messages, which brings many security risks.

Origin Hijacking



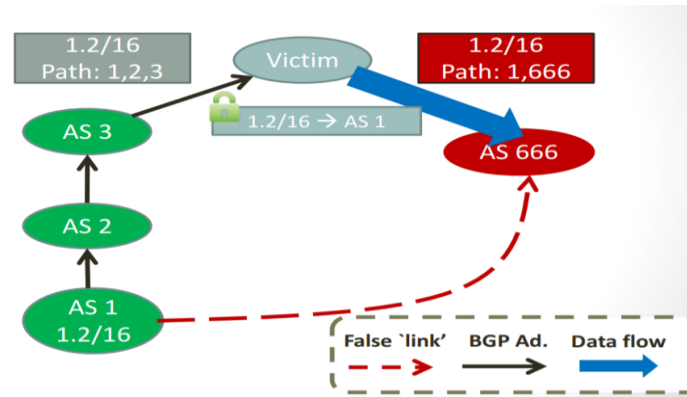
- › Drive traffic by announcing address prefixes that are not their own
- › Google announced Verizon an IP address segment that was originally attributed to NTT. Verizon sent traffic to NTT to Google, causing Japan to disconnect for 1 hour.

<https://www.thedrum.com/news/2017/08/28/google-hijack-made-japan-land-no-internet-more-30-minutes>

Route Leak

"Google was also the victim of a routing leak. In this case Google's prefixes were leaked by Hathway, and accepted by their peer Bharti Airtel. Bharti then advertised routes to dozens of major ASes around the globe. In Figure 5, we can see the leak of an existing prefix 74.125.200/24 from Hathway, with traffic from Bharti (AS9498) transiting via Hathway (AS17488) to Google. This leak lasted for nearly a day, from 10:30 UTC on March 11th to 9:15 UTC on March 12th. "

Path Hijacking



- › Using the characteristics of the AS_PATH attribute being easy to modify, announce incorrect path information to hijack traffic.
- › AS 666 deliberately announced incorrect information, claiming that it was only one hop away from AS1, causing all traffic destined for AS1 to be hijacked to AS666.

*AS (Autonomous System)

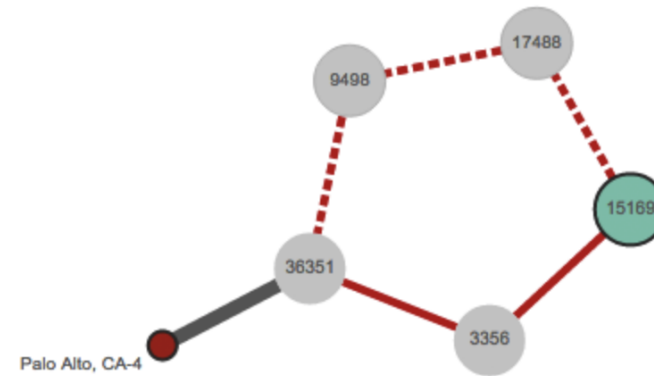


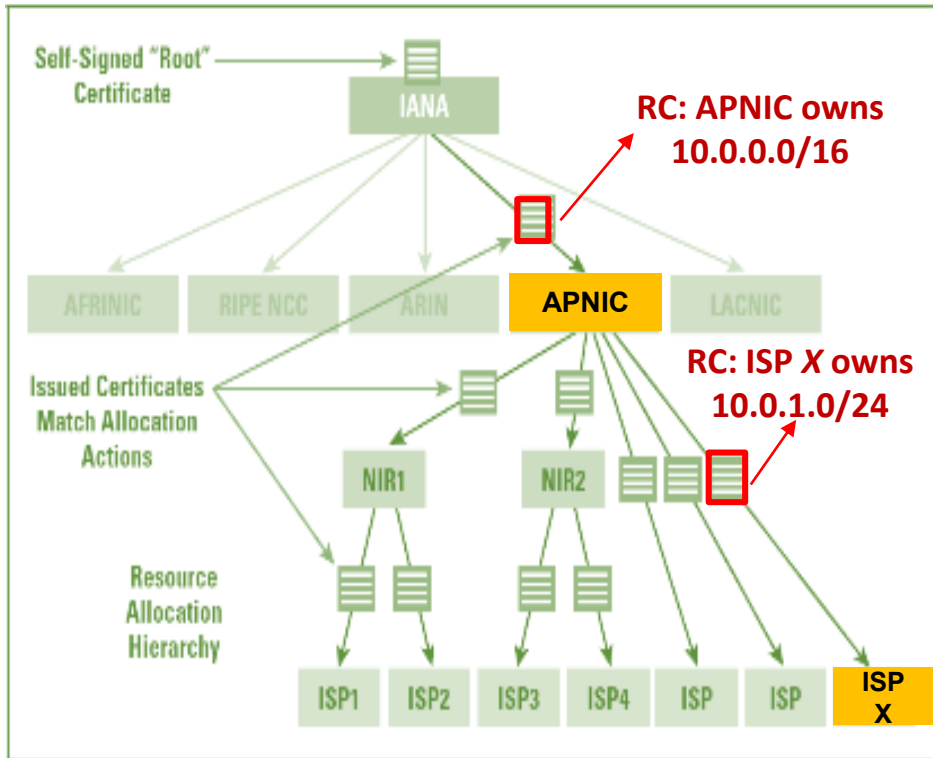
Figure 5: Route leak to Google via Hathway AS17488 that affects Bharti Airtel AS9498.

<https://blog.thousandeyes.com/finding-and-diagnosing-bgp-route-leaks/>

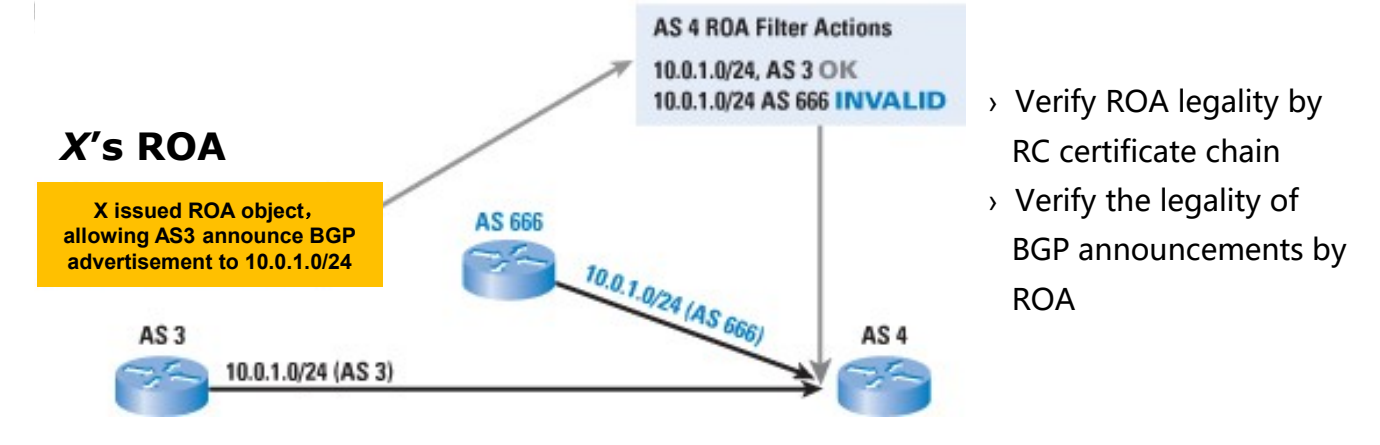
<https://www.internetsociety.org/blog/2018/01/14000-incidents-2017-routing-security-year-review/>

IETF proposed RPKI and BGPSEC

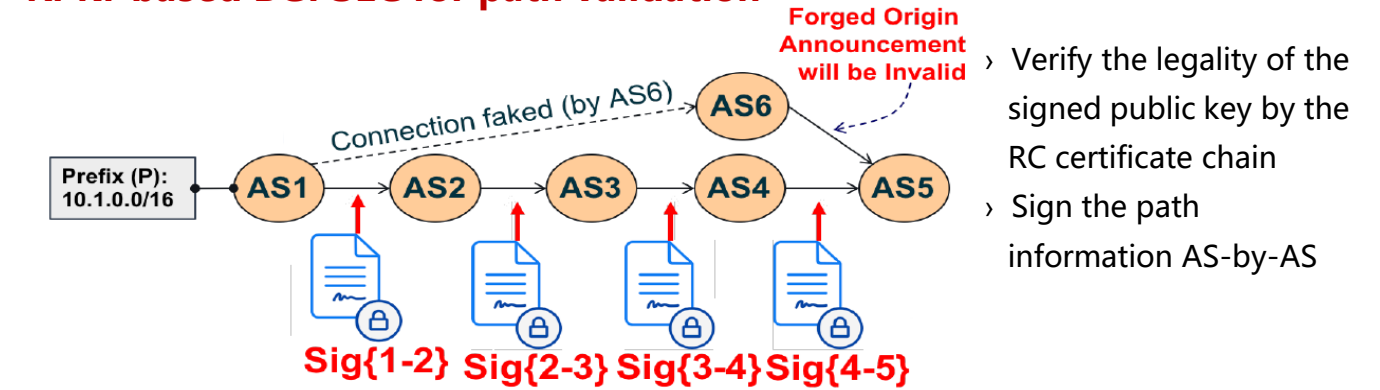
- RPKI provides RC certificate-based verification capabilities
 - › Use Resource Certificate to prove address ownership
 - › The issuance of RC depends on the allocation process of the address



RPKI ROA for BGP source verification

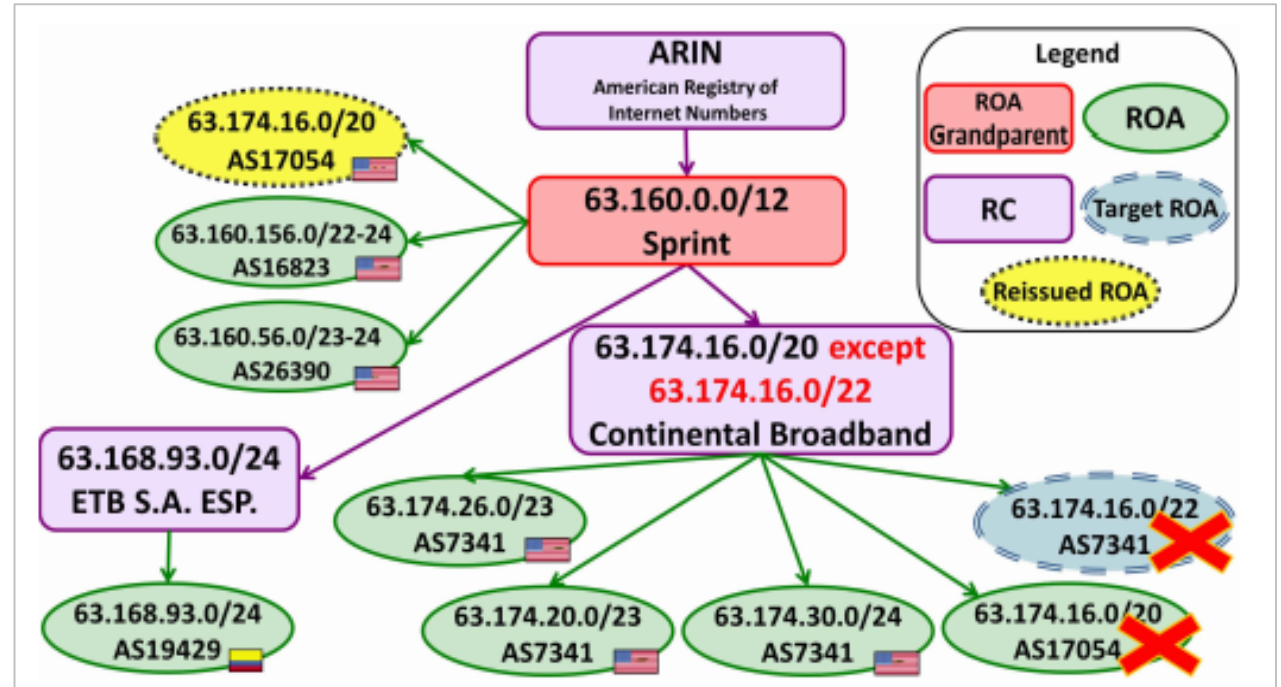


RPKI-based BGPSEC for path validation



BUT...RPKI does not completely solve the problems and introduce centralization issues

- Depending on the centralized trust model, once the Authority node is misconfigured or attacked, it will cause security issues
 - › Certificate revocation/overwrite: Unilaterally cancel the issued RC certificate, causing the BGP announcement of the lower node to be invalid; equivalent to depriving the applicant of the ownership of the IP address.
 - › ROA (Route Origin Authorization) coverage: The superior node issues an ROA that has been distributed to the subordinate institution prefix to attract part of the traffic.
- Path validation requires hop-by-hop signature decryption which affects route convergence speed.



Heilman E, Cooper D, Reyzin L, et al. From the Consent of the Routed: Improving the Transparency of the RPKI[C]//ACM SIGCOMM Computer Communication Review. ACM, 2014, 44(4): 51-62.

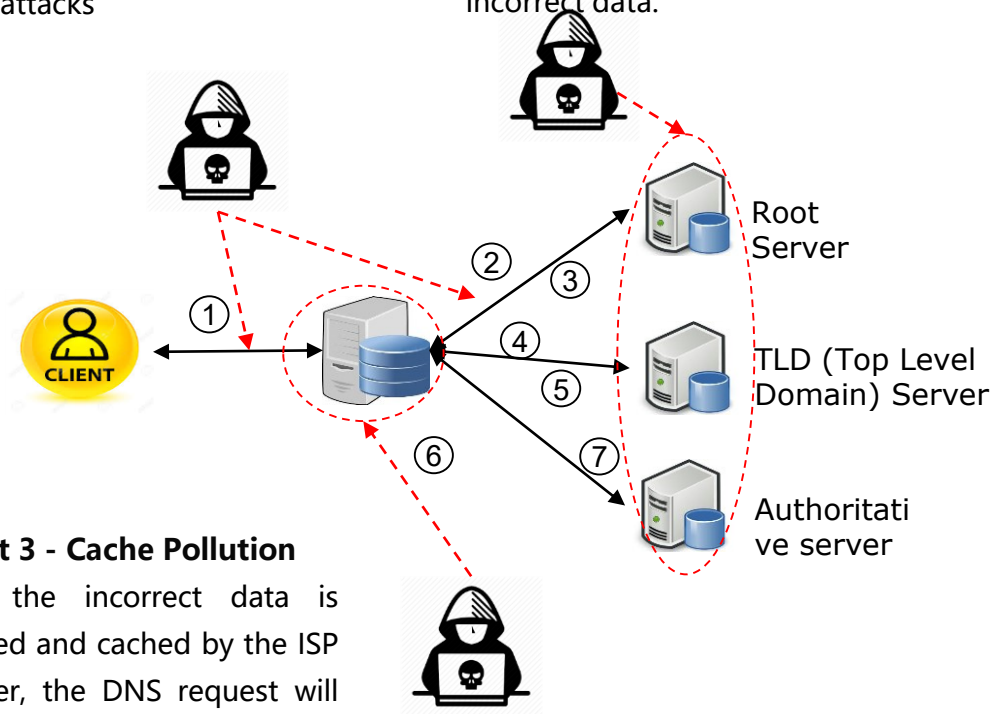
■ Real Case Scenario

- › In Dec,2013, A ROA (79.139.96.0/24, AS 51813) was accidentally deleted, resulting in a certain part of the network prefix in Russia became unreachable.
- › In Jan,2014, the ROA of one of Nigeria' s network was "invalid" , because its parent RC was overwritten.
- › In Dec,2013, ARIN mistakenly issued a ROA, allowing AS6128 to announce the prefix 173.251.0.0/17~24, causing the legal declaration of the prefix to become invalid.

DNSSEC also cannot completely solve the security threats and centralization problems of DNS

■ Threat 1 - DNS Hijacking

- › The data of any link on the DNS resolution path may be subject to MITM attacks



■ Threat 3 - Cache Pollution

- › Once the incorrect data is received and cached by the ISP resolver, the DNS request will receive the incorrect data for a long time.

■ Treat 2 - Chained threat

- › Any device on the DNS resolution tree may be attacked and return incorrect data.

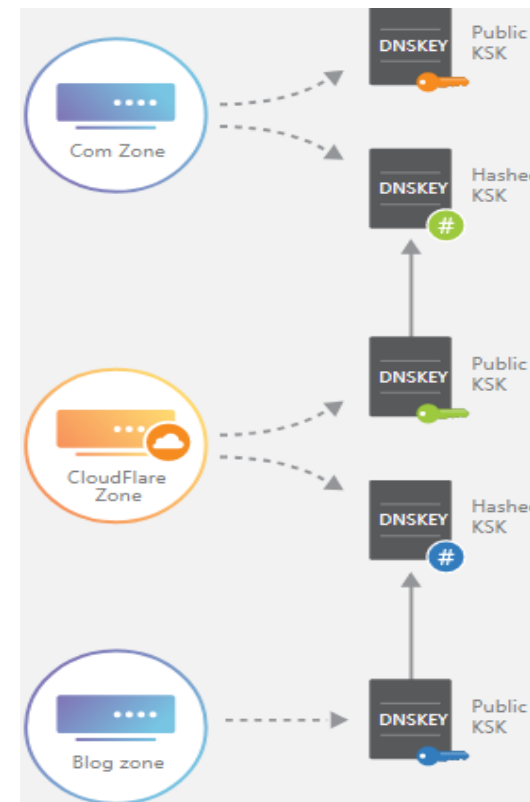
■ DNSSEC only solve DNS hijacking problems

- › Depends on signature information to ensure data integrity
- › Based on the basic principle of PKI, it verify the DNSKEY of the subzone rely on the DNSKEY of the parent zone.

■ Centralization still exists

- › Unilaterally delete the sub-domain DS records in its zone file, so that the subzone's KEY is not trusted.
- › Unilateral fake subzone' s DNSKEY and signed it.

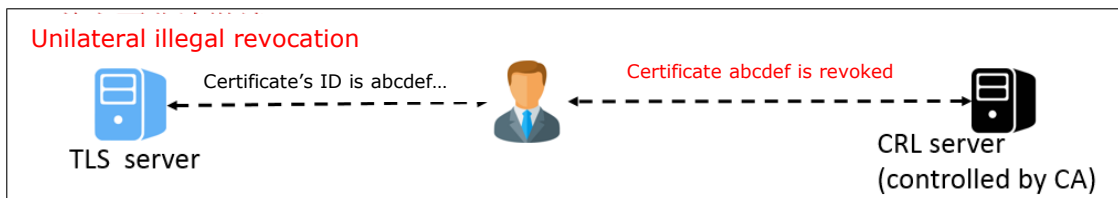
■ Centralization problems still cause cache pollution



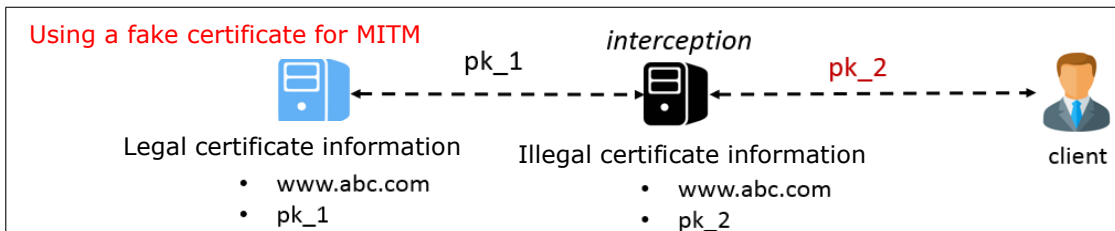
PKI also faces security vulnerability and trust chain failure of central nodes

- All the control of the certificate are owned by the CA, so if the CA is attacked, it will bring the following security threats:

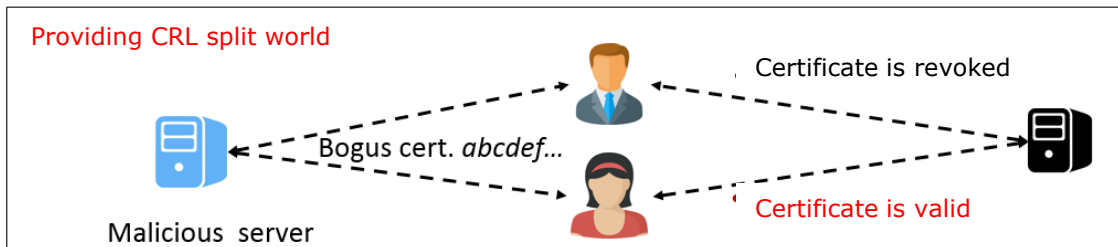
- › Illegal revocation



- › Issuing an illegal identity



- › Issue or revoke CRL



* CRL: Certificate Revocation List

■ Real Case Scenario

- › In Jul, 2011, the Netherlands noted that 8 servers of CA DigiNotar were hacked. At least 531 false certificates were released including Yahoo!, Mozilla, WordPress, The Tor Project, etc.
- › In Jul, 2011, Google service suffered from the above-mentioned illegal certificate attack, affecting the Dutch financial, technology, manufacturing and other industries.

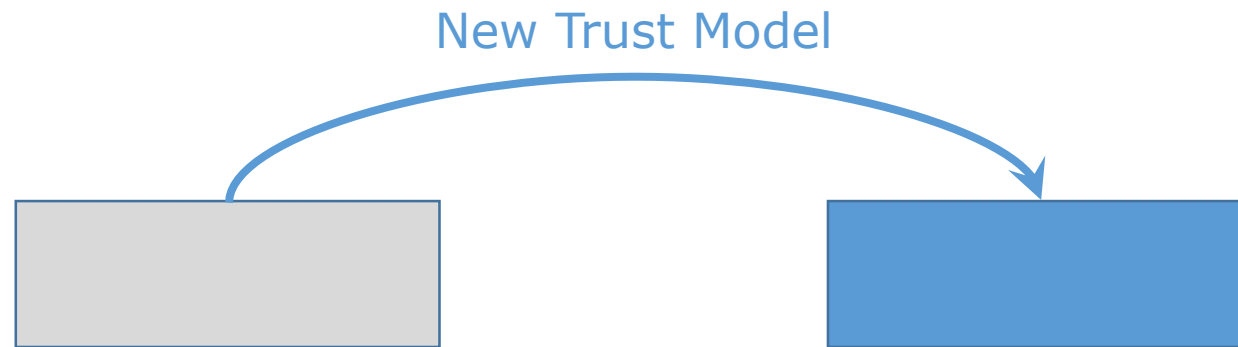
■ Certificate Transparency

- › Use the public certificate Log to a certificate signing example
- › Only be detected afterwards, p evidence of responsibility
- › Unable to fundamentally solve centralization problem of PKI



Privacy Protect and Data Sharing

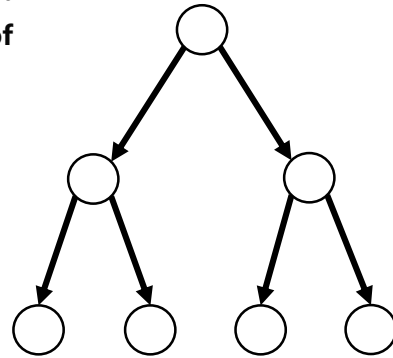
- It is very IMPORTANT now and future
- However, the current trust model can hardly work
- A novel trust model may support more upper layer applications



The root cause is the centralized trust model.

Where is the current problem?

- Reason 1: BGPSEC, DNS (SEC), and PKI all adopt a centralized trust model. There is a single point of security and credibility in the mechanism. Without changing the architecture, it is difficult to solve.
- Reason 2: At a deeper level, the current solution is a patched solution, which does not fundamentally examine where the Internet security credibility is.

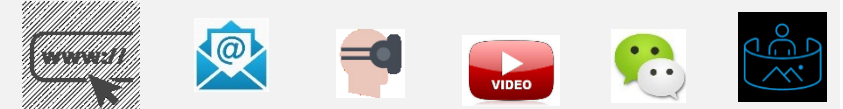


Where is the security and trustworthy foundation?

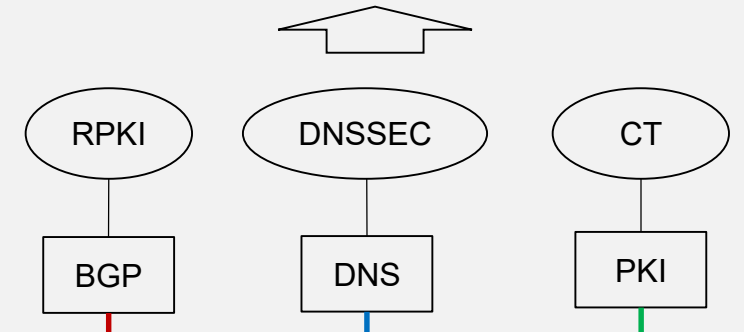
- Decentralization technology to solve problems naturally
 - Additional benefits: increased reliability, increased security, reduced latency...
- Not depending on a single trust anchor is the basis of network security and trustworthiness.
 - What is needed for the Upper Layer: Mapping between resource information
 - What is the dependency of the mapping: the mapping information authorized by the resource owner is trusted

Current network trust model

Application

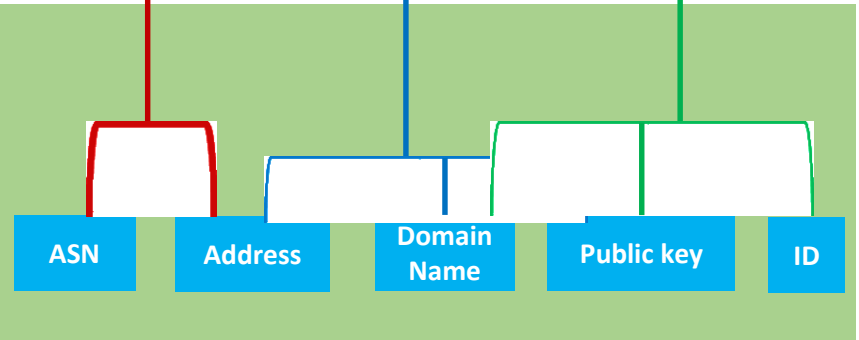


Infrastructure



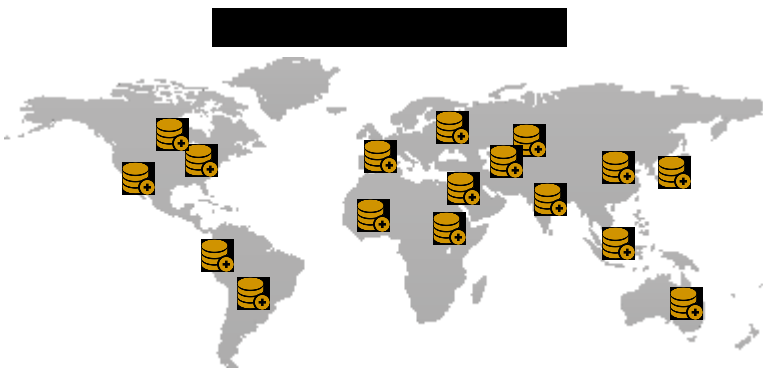
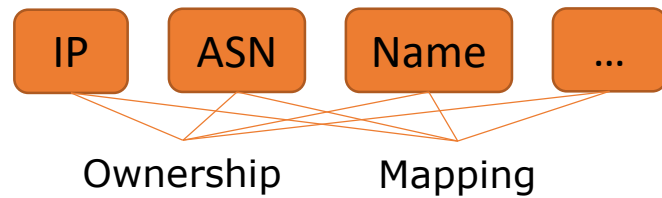
Information mapping

Resources management



Consider to introduce decentralized trust model for resource management and privacy protect

The overview of DNI (Decentralized Network Infrastructure)



Third-party decentralized APP platform

- Decentralized PKI platform
- Pay remote DDoS defense services on demand
- Cross-domain end-to-end QoS capabilities

Trusted name space ownership and mapping system

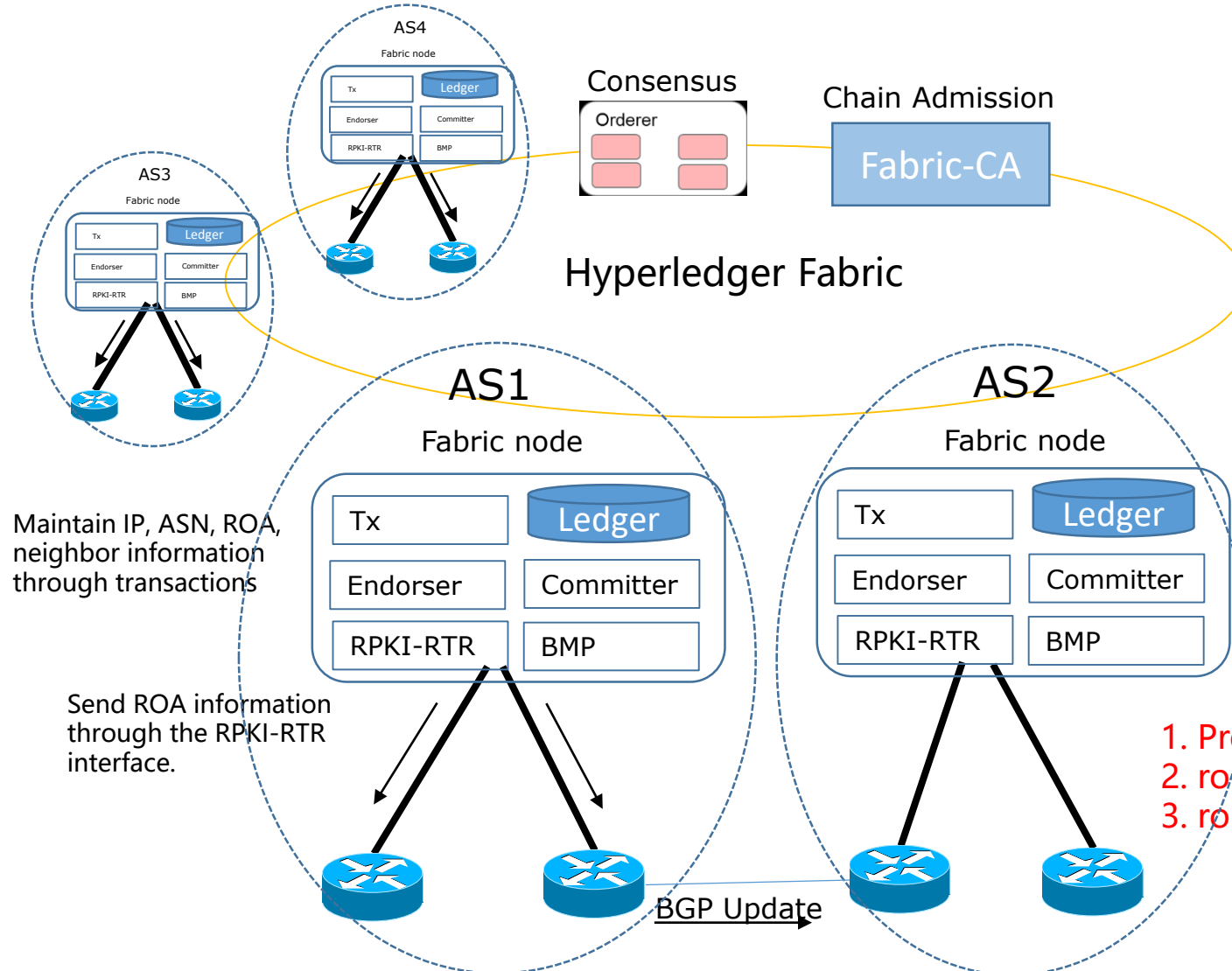
- IP and ASN: Trusted Routing System
- IP and domain name: Trusted DNS resolution system
- Other: trusted host ID, trusted content, trusted IoT ID, etc.

Decentralized network infrastructure based on blockchain

- Decentralized (p2p) network architecture and trusted model
- Consensus mechanism
- Smart contract for computing models
- Monetization trading platform for Internet services

A consortium chain-based DNI Verification System

Blockchain stores Ownership, ROA and neighbor information



Maintain IP, ASN, ROA, neighbor information through transactions

Send ROA information through the RPKI-RTR interface.

1. Prefix origin verification
2. route path validation
3. route leak detection

IP Ownership

IP	Owner	Exp date
1.1.1.0/24	ISP1	19/10

ASN Ownership

ASN	Owner	Exp date
100	ISP1	19/10

ROA (IP->ASN)

IP	Maxlength	ASN
1.1.1.0/24	32	100

ASNeighbor(ASN->ASN)

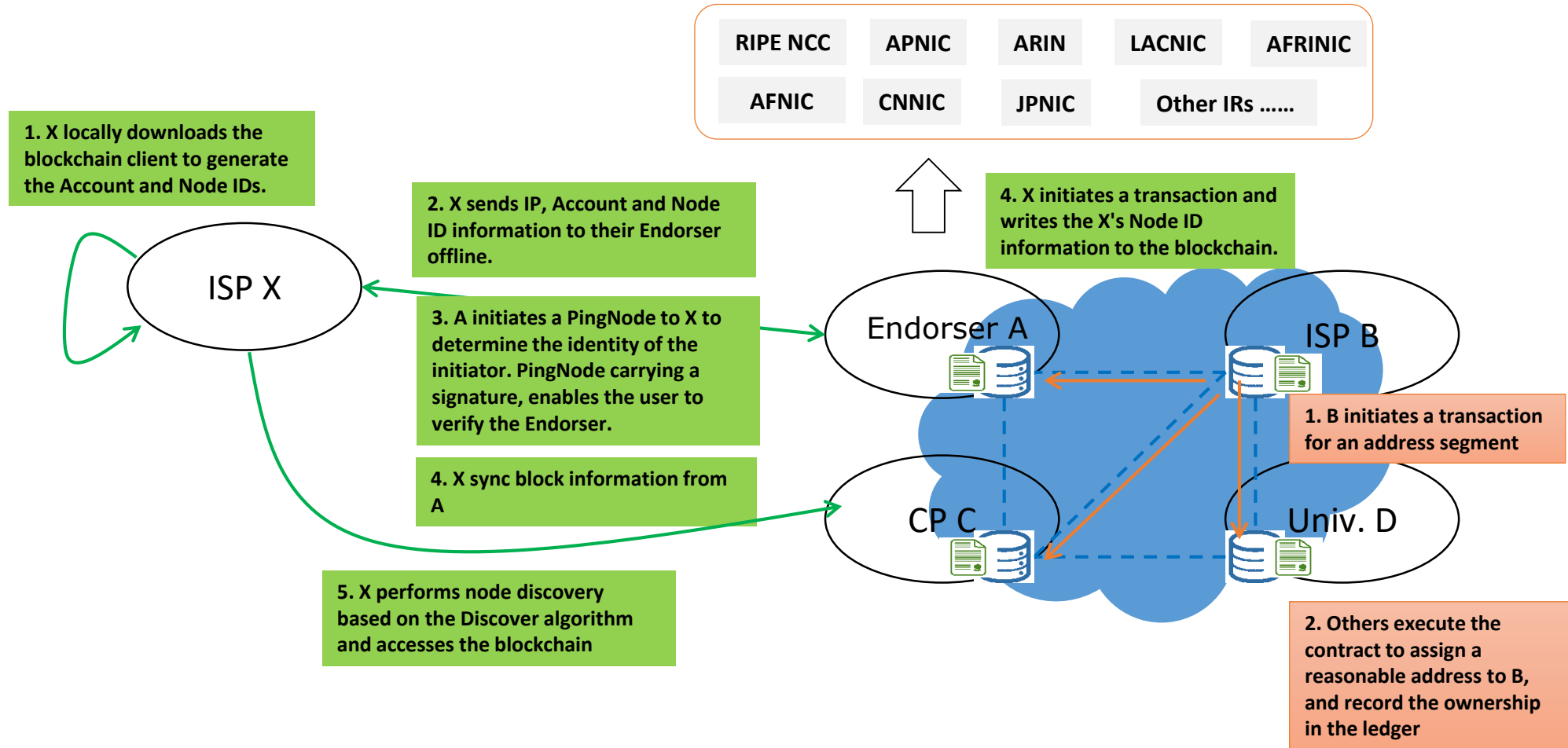
Source	Target	Type
AS1	AS2	P2C
AS2	AS3	P2P

World-state

- RPKI-RTR: RPKI to Router Protocol
- BMP: BGP Monitoring Protocol

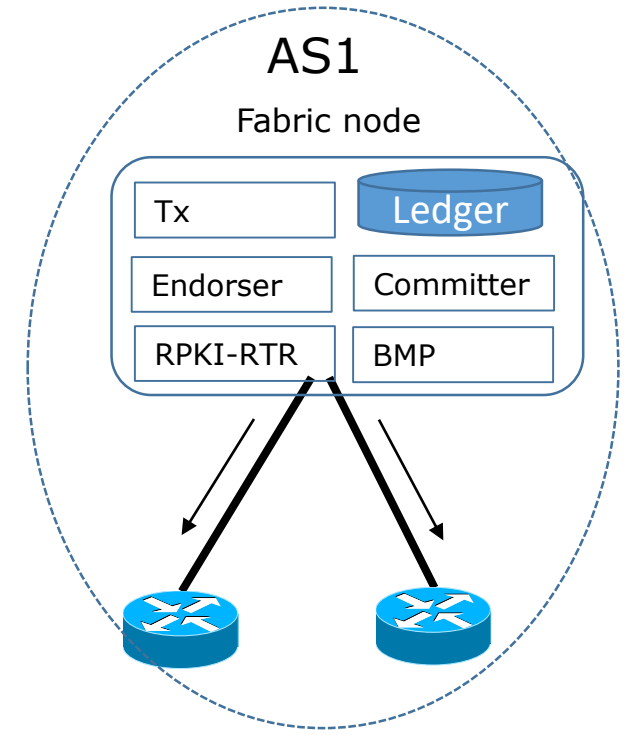
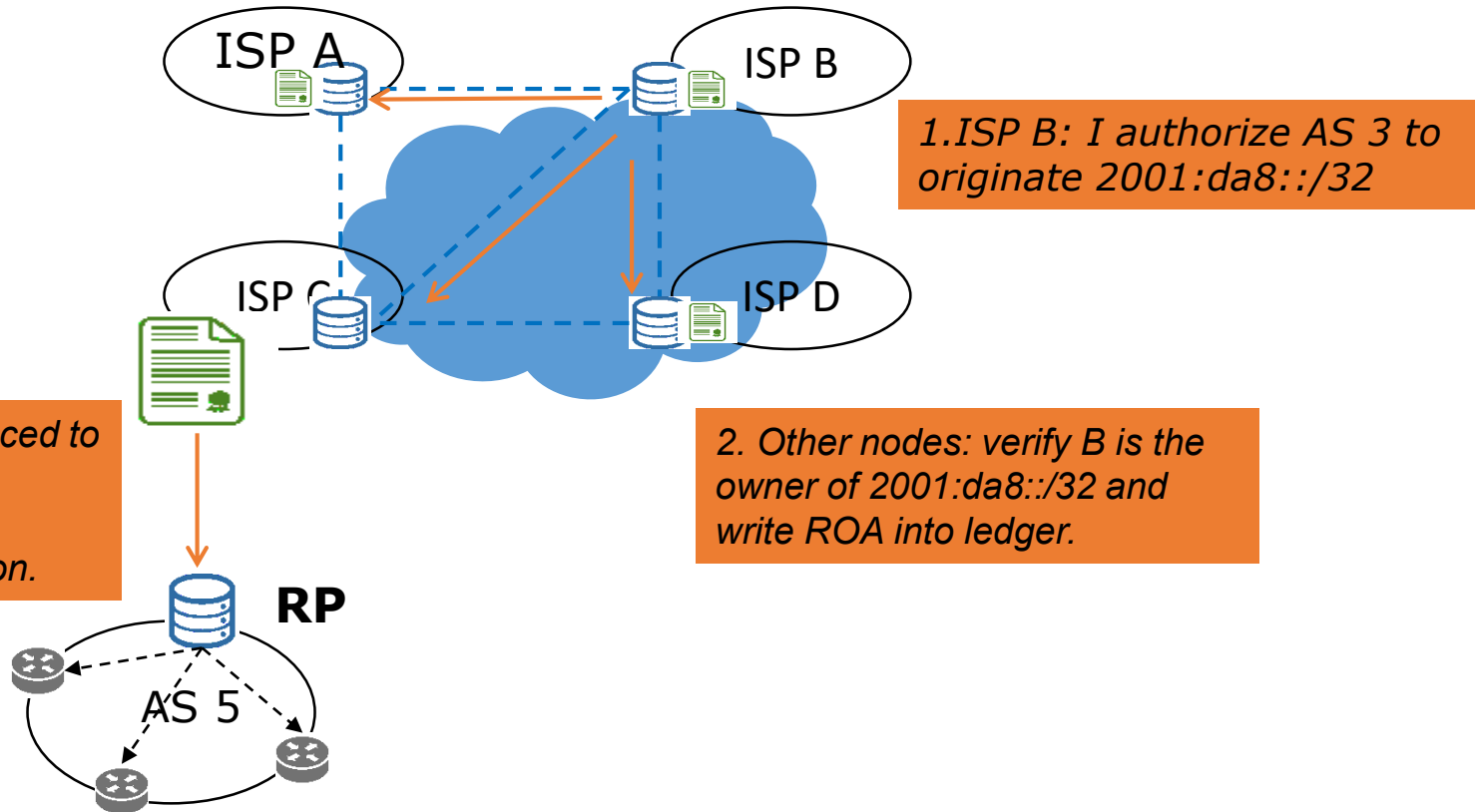
IP address management and access control

- Simultaneous implementation of endorsement access control and dynamic node management
- The blockchain application layer is reversed from the underlying network layer, allowing the network layer to implement a dynamic node admission control strategy based on the consensus result of the application layer.



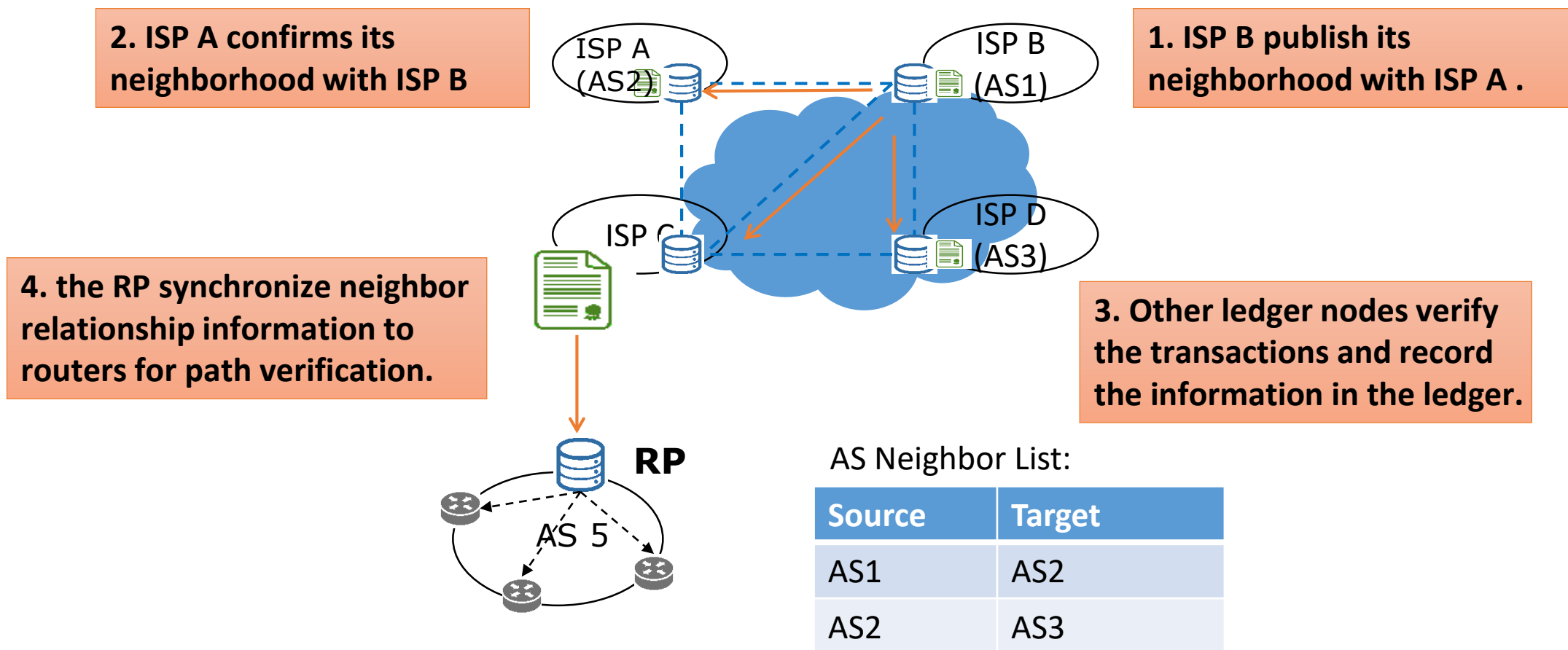
DNI-based BGP Verification - Origin Verification

1. IP address owner initiates an ROA (IP to ASN mapping) as a transaction.
2. Smart contract verifies the address ownership, and writes the ROA into the ledger.
3. Relying parties (RP) get updated ROAs from the ledger, and sync to BGP routers, which then verify BGP routes.



DNI-based BGP Verification - AS Path Verification

1. Each AS publishes its neighbor information in the ledger for AS path verification in BGP advertisement.
2. The Relaying Party (RP) get neighbor information from the ledger and synchronize the information to routers.



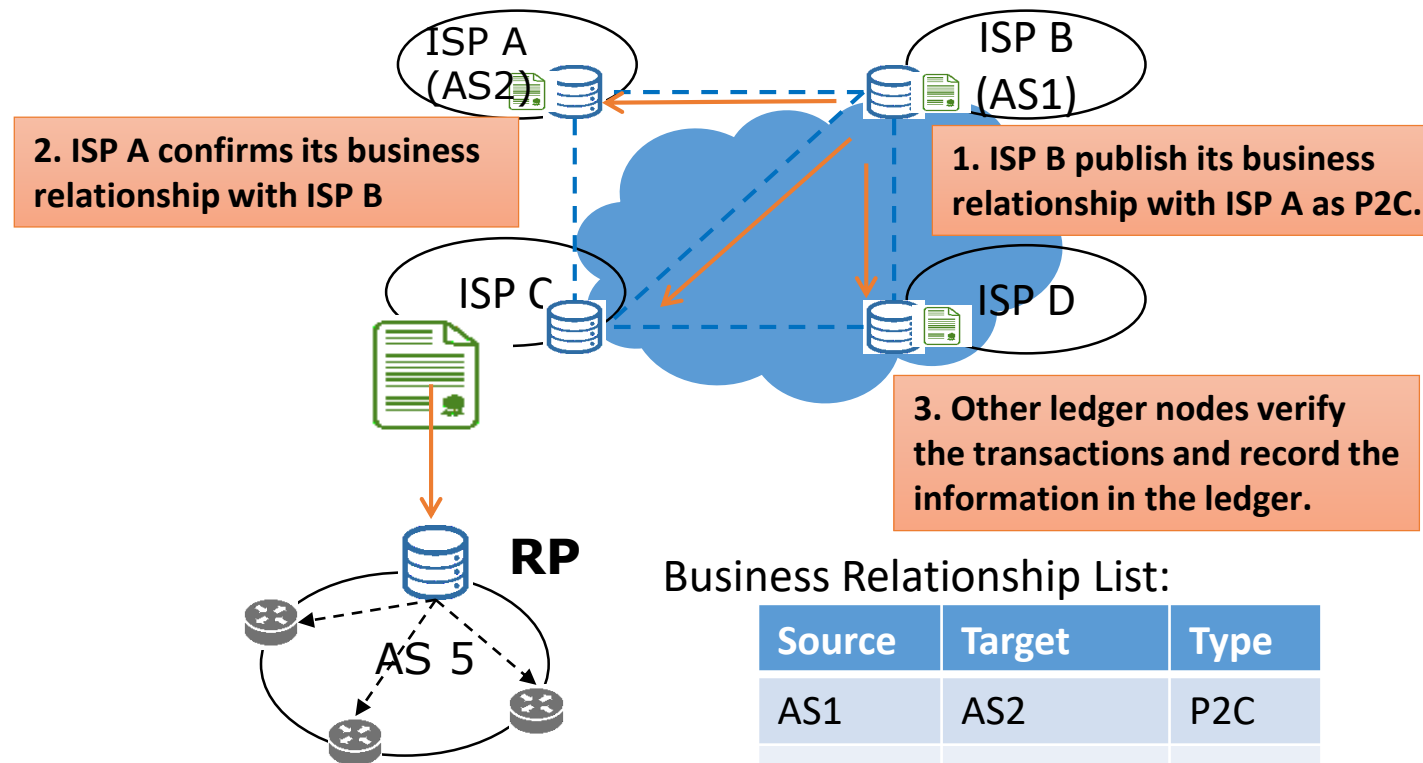
DNI-based BGP Verification - Route Leak Protection

Publish of Business Relationship between ASes

- Each AS registering their business relationship with their neighbors into the ledger.
- The business relationship will be certified by the pair of ASes.

Route leak detection based on ASes' business relationship information

- The Relying Party obtains and analyzes the global neighbor business information from the ledger to generate a route filtering table.
- The Relying Party synchronizes route filtering table to routers.
- Router check each hop of AS Path to decide whether the route leak rule is violated or not.



4. the RP synchronize business relationship information to routers for route leak detection.

D->C->B->A

Business Relationship List:

Source	Target	Type
AS1	AS2	P2C
AS2	AS3	P2P

route leak rules:

Relationship for current hop	Relationship for previous hop	Result
P2P	P2C	Leak
P2P	P2P	Leak
C2P	P2C	Leak
C2P	P2P	Leak

Decentralized domain name management

- The domain name is bound to the public key. As long as the private key is signed, anyone can operate the related domain name.
- Agency needs applicant to provide transaction proof information

Applier X



0. X applies for the domain name example.com to Agency A and provides the relevant public key pk_X

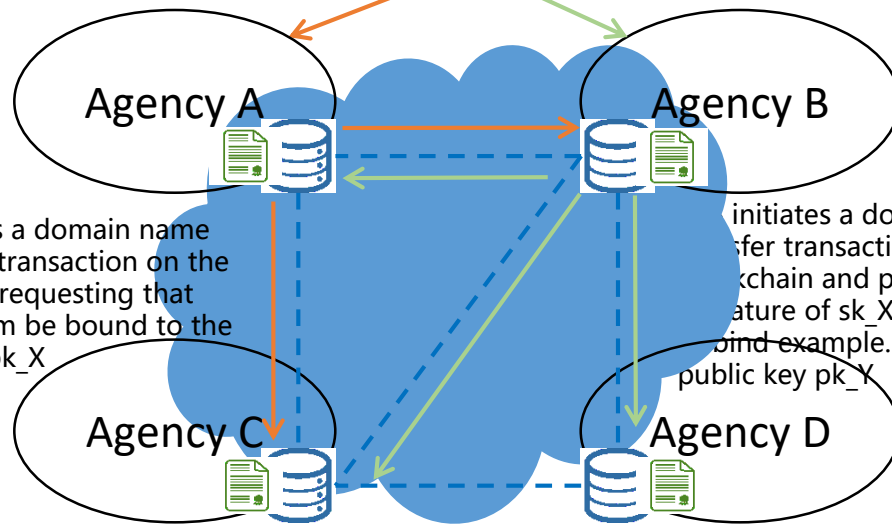
3. X initiates a request to Agency B, which needs to transfer the domain name to pk_Y; and provides the signature of sk_X

1. A initiates a domain name application transaction on the blockchain, requesting that example.com be bound to the public key pk_X

4. B initiates a domain name transfer transaction on the blockchain and provides the signature of sk_X, requesting to bind example.com to the public key pk_Y

2. Others: Verify that the transaction of A is legal; if it is legal, write the information to the blockchain, and the owner public key of the domain name example.com is pk_X

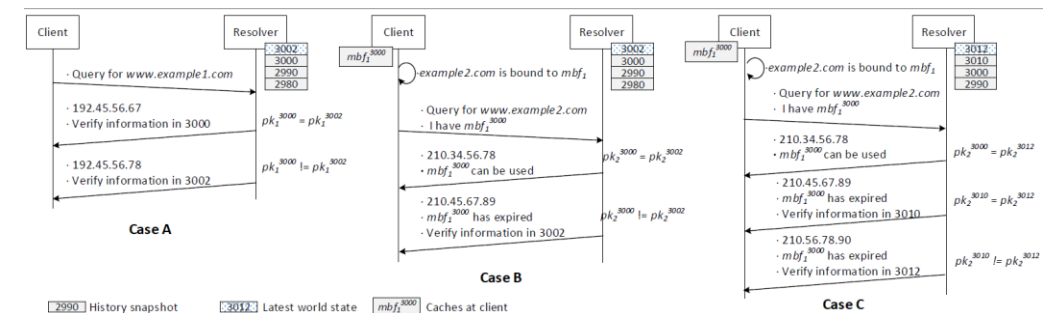
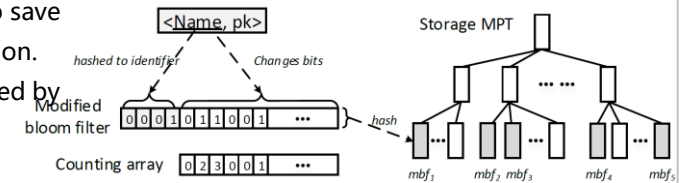
5. Others: Verify that the signature in the transaction initiated by B is legal; if it is legal, the owner of the record example.com is changed to pk_Y



Lightweight data verification mechanism

- In the current DNS system, the client does not have any ability to verify the data authenticity, and can only trust the resolution result unconditionally.
- The Blockchain provides the SPV (Simplified Payment Verification) mode, but it needs to obtain the latest blockchain information to verify each time. The single overhead is at the KB level.
- This mechanism reduces the single verification overhead to the Bytes level.
- A blockchain-based DNS information verification and caching mechanism security enhancement

- › Add a bloom filter to the contract to save the existence of the owner information.
- › Verification information can be reused by the cache
- › A bloom filter can be used to verify multiple domain ownership information



■ Enhance the security capabilities of the DNS protocol instead of DNSSEC

- Data integrity (DNSSEC)
- Cache pollution
- Data authenticity

```
www.example.com A      2.2.2.2
www.example.com RRSIG  xxxxxxx
```

DNS Client



4. Verify ownership;
Verify RRSIG signature



1.1.1.1

0.1 Set my authoritative domain name server information to 1.1.1.1 (also provide sk_X signature)

2. Initiate a DNS request to the authoritative server

1. DNS request www.example.com

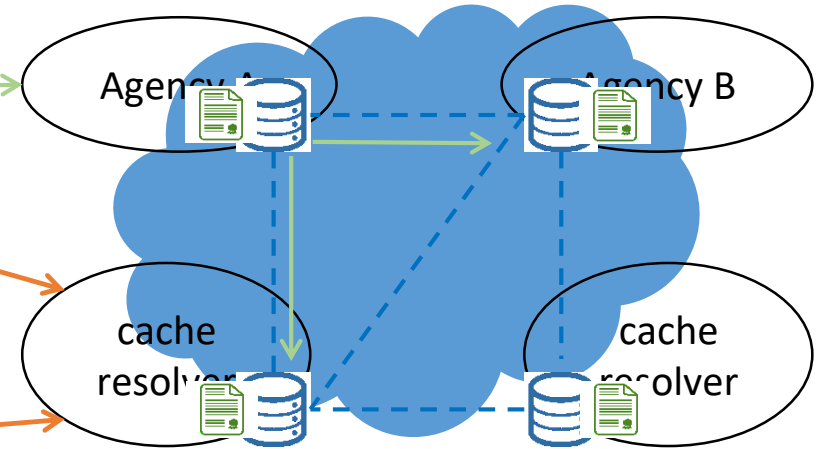
3. DNS response

```
www.example.com A      2.2.2.2
www.example.com RRSIG  xxxxxxx
Other blockchain verification information
```

- The blockchain only stores the ownership information and the authoritative server information because the update frequency of ownership is very low.

0.2 Initiate a transaction on the A blockchain to maintain information

example.com pk_X

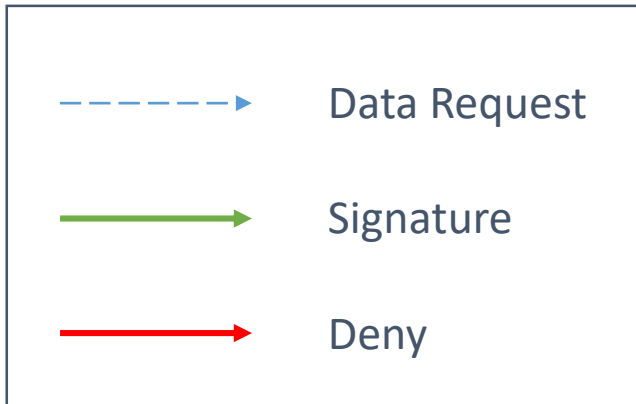


0.3 Other nodes verify that the signature is correct; if correct, write maintenance information to the blockchain.

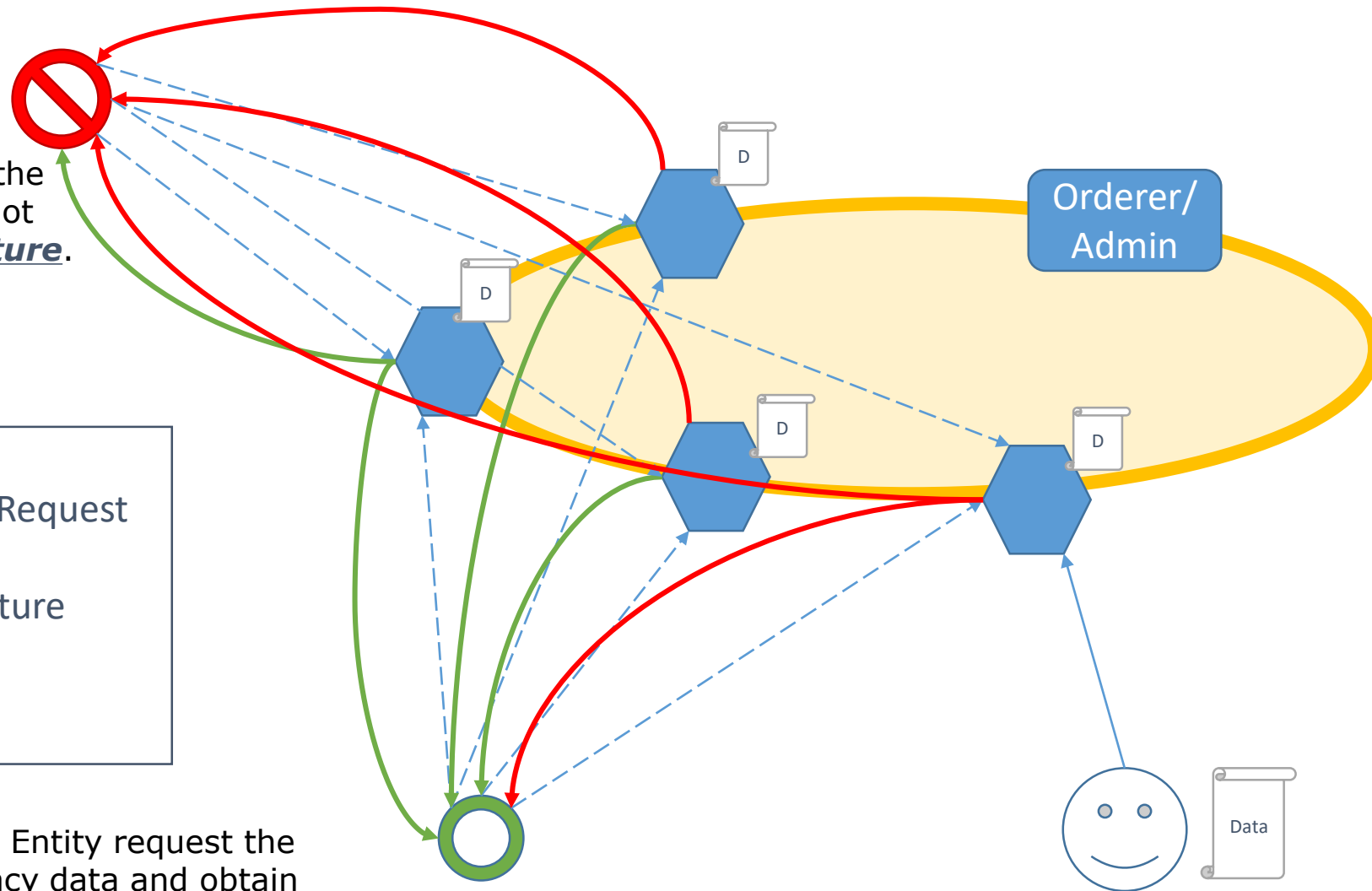
Domain name	Owner	Authoritative server
example.com	pk_X	1.1.1.1

Improve RPKI and Privacy Protect

Illegal Entity request the privacy data and cannot obtain enough **Signature**. FAILED.



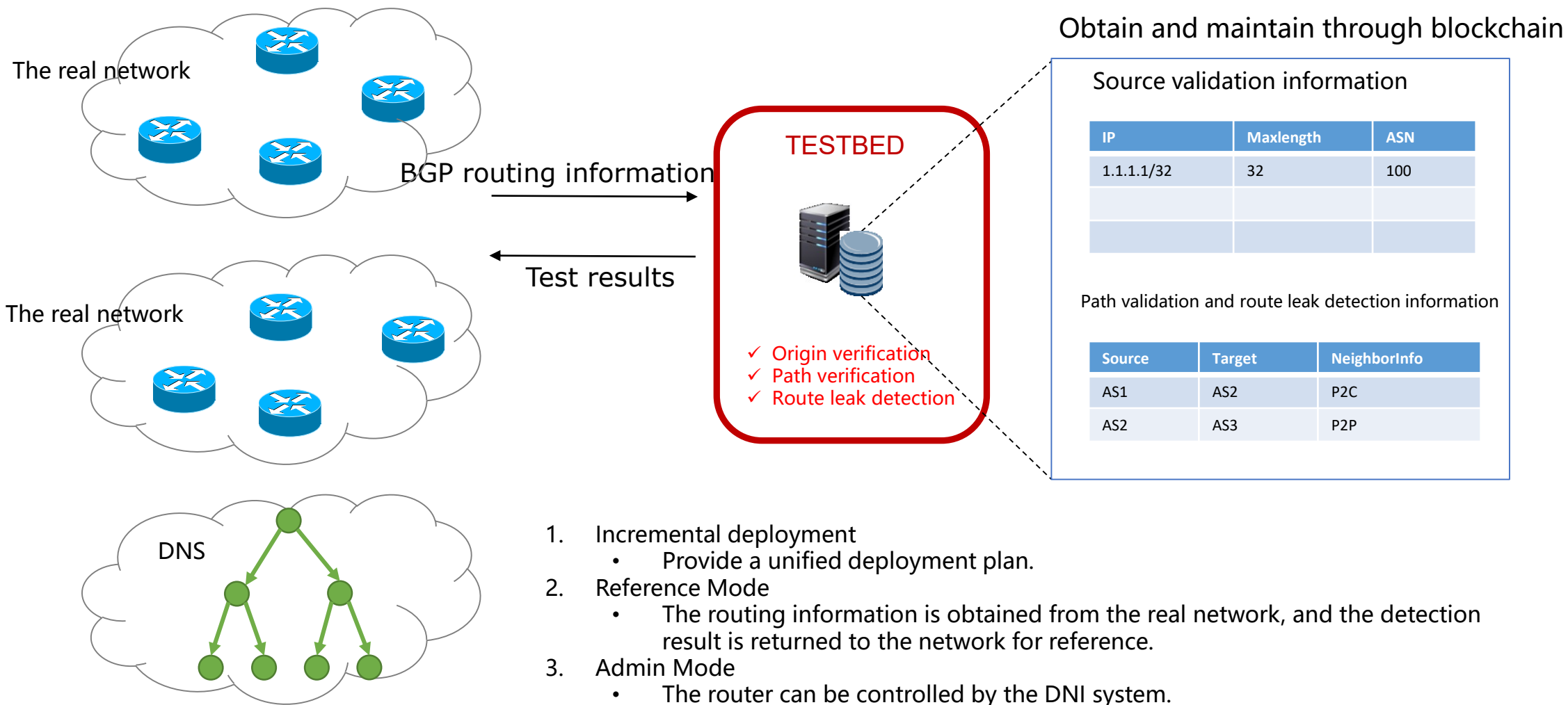
legal Entity request the privacy data and obtain enough **Signature**. SUCCESS.



Share the data with a novel trust model.

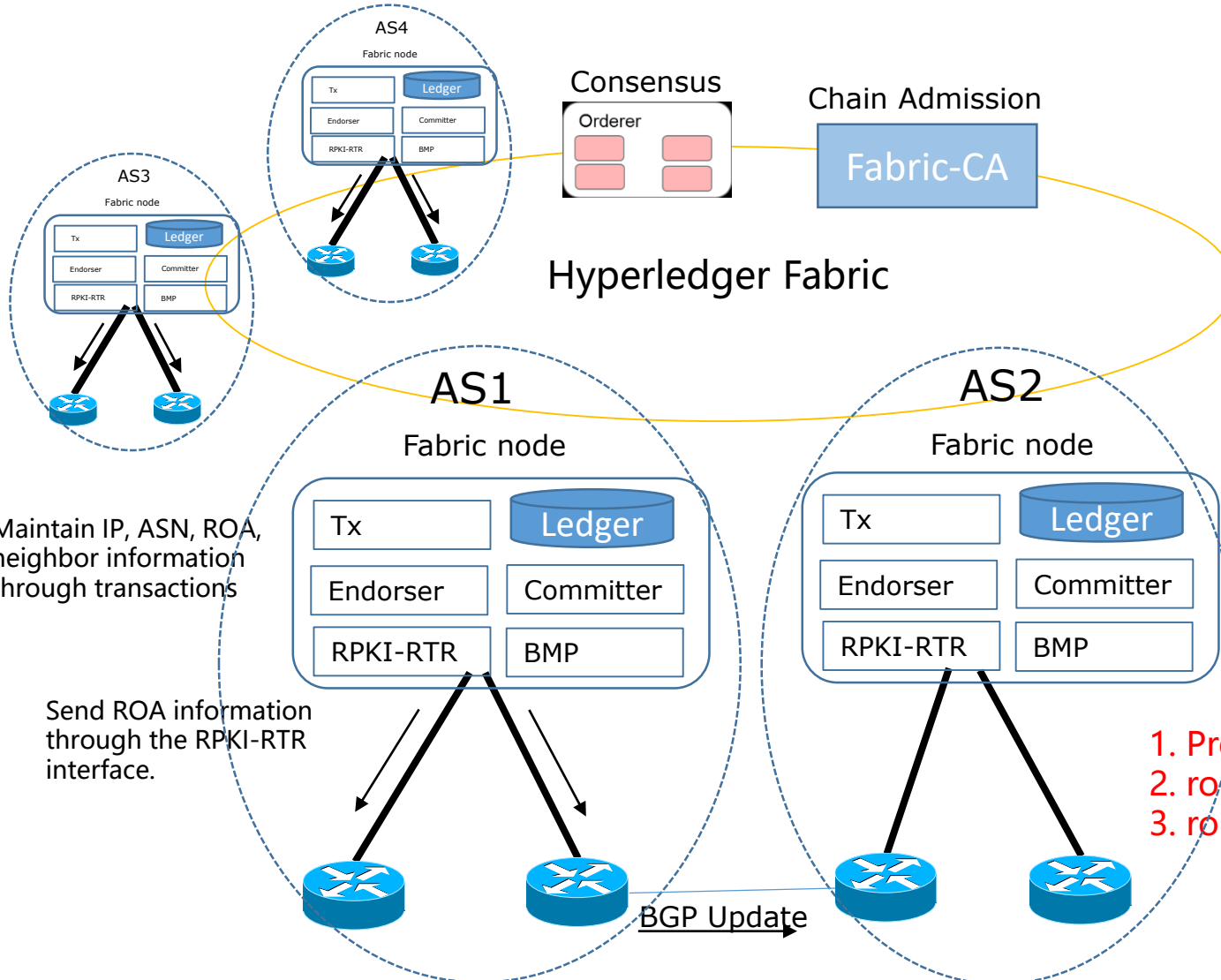
Testbed of BGP security, address management and DNS security based on blockchain

- Solve the single point problem of RPKI.
- Provides a unified solution to support origin validation, path validation, and route leak detection.



DNI System Overview

Blockchain stores Ownership, ROA and neighbor information



Maintain IP, ASN, ROA, neighbor information through transactions

Send ROA information through the RPKI-RTR interface.

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- RPKI-RTR: RPKI to Router Protocol
- BMP: BGP Monitoring Protocol

IP Ownership

IP	Owner	Exp date
1.1.1.0/24	ISP1	19/10

ASN Ownership

ASN	Owner	Exp date
100	ISP1	19/10

ROA (IP->ASN)

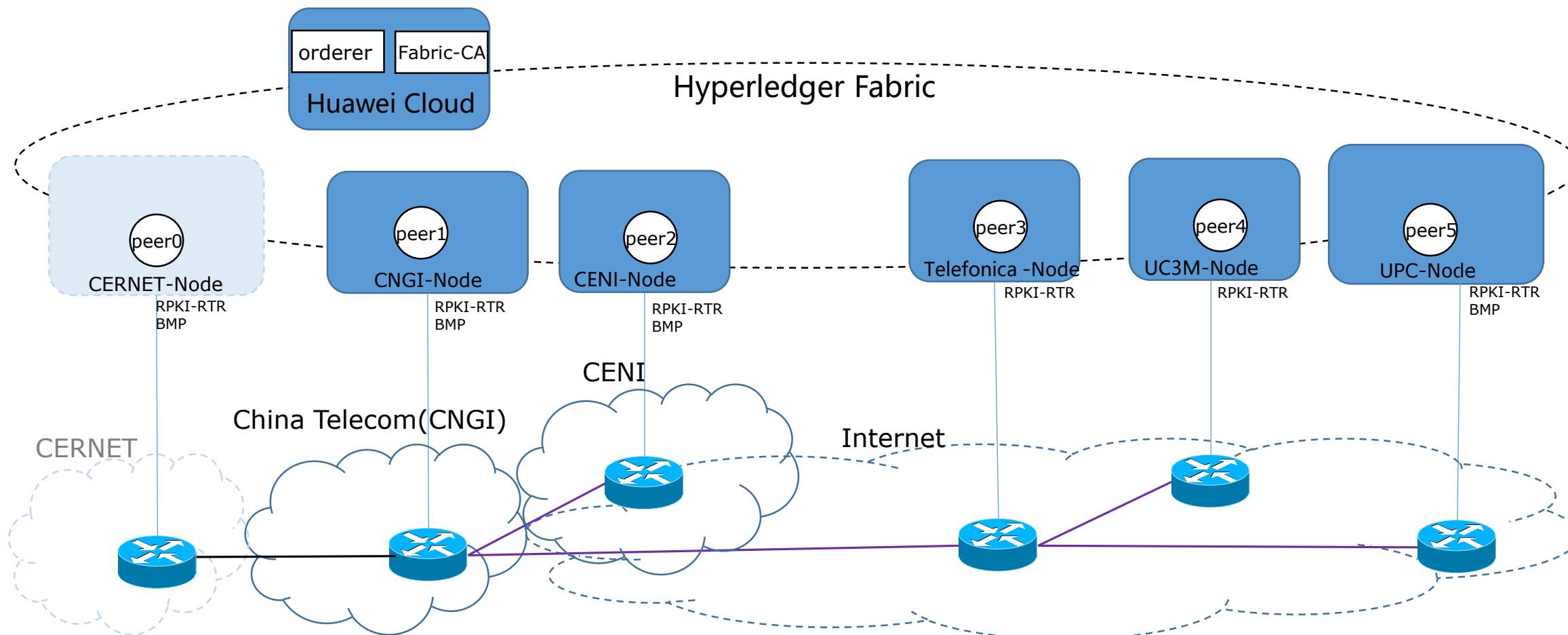
IP	Maxlength	ASN
1.1.1.0/24	32	100

ASNeighbor(ASN->ASN)

Source	Target	Type
AS1	AS2	P2C
AS2	AS3	P2P

World-state

DNI Testbed based on Consortium Chain



- **China Telecom, Telefonica, CENI, CERNET2, UC3M, UPC, BUPT, Tsinghua, ...**

Current Testbed operation

- Based on ethereum / Hyperledger Fabric
- Smart contract development:

by Remix GUI

Endorsement: RIR endorses the ISP User.

IP Allocation: allocate IP to ISP by sparse_allocation

The screenshot shows a transaction hash: `0xc4534293e0e0c02251208180c1c04c4d9c04d1e0190005a5e479`. The decoded output is a JSON object:

```

{
  "0": "address: addr",
  "1": "string name APNIC",
  "2": "bool irrevocable true",
  "3": "bool irrevocable true"
}

```

The screenshot shows a transaction hash: `0x14723a09ac7e02c400c0f7a44af00f00c160c`. The decoded output is a JSON object:

```

{
  "0": "bytes16 ip",
  "1": "address oldOwner",
  "2": "address newOwner",
  "3": "base_prefix",
  "4": "base_prefix"
}

```

ROA: ISP announces the ROA

Askprice: get the realtime ether price

The screenshot shows a transaction hash: `0x2c95215a504135a4d4e0808e0c7912078663b05c3480c5a0122323e`. The decoded output is a JSON object:

```

{
  "0": "bytes16 ip",
  "1": "address oldOwner",
  "2": "address newOwner",
  "3": "base_prefix",
  "4": "base_prefix"
}

```

The screenshot shows a transaction hash: `0x430046827649026e0975732130634e45464702532279e4072a1471`. The decoded output is a JSON object:

```

{
  "0": "uint256 gas",
  "1": "uint256 gas",
  "2": "uint256 gas",
  "3": "uint256 gas"
}

```

- UI & Relying party work is ongoing.

- ITU-T SG13 Q2 WI, Framework and Requirements of Decentralized Trustworthy Network Infrastructure, https://www.itu.int/ITU-T/workprog/wp_item.aspx?isn=15083
- ITU-T DLT FG use case, <https://www.itu.int/en/ITU-T/focusgroups/dlt/Pages/default.aspx>
- IETF dinrg Presentation
 - <https://datatracker.ietf.org/doc/slides-102-dinrg-decentralized-internet-resource-trust-infrastructure-bingyang-liu/>
 - <https://datatracker.ietf.org/meeting/105/materials/slides-105-dinrg-a-blockchainbased-test-bed-for-bgp-verification-00>
- ETSI PDL ISG, <https://portal.etsi.org/TB-SiteMap/PDL/List-of-PDL-Members-and-Participants>

- Decentralized Trust Model can improve the network trust scheme
 - Protect the whole system from single trust anchor failure
 - Improve the privacy and security
 - Co-work with the current trust model
- The BlockChain is not the key but the decentralized idea
- CALL for Joint research and deployment

THANK YOU